Computer Laboratory

Exercises in

LINEAR PROGRAMMING AND COMBINATORIAL OPTIMIZATION

LUND INSTITUTE OF TECHNOLOGY

DEPARTMENT OF MATHEMATICS

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Preparation for the Labs

Download the Matlab files needed for the exercises (lab1.zip and lab2.zip) from the course website: http://www.maths.lth.se/education/lth/courses/linkombopt/2013/.

Some Useful MATLAB Functions

First an explanation of the command is given with the MATLAB `help` command and then an example is given:

```
 FIND   Find indices of the non-zero elements.
        I = FIND(X) returns the indices of the vector X that are
        non-zero. For example, I = FIND(A>100), returns the
        indices of A where A is greater than 100. See RELOP.

>> b=[3 5 -2 0 9 -3 -4 8];
>> find(b>0)
ans =
     1     2     5     8

>> help max

 MAX    Largest component.
        For vectors, MAX(X) is the largest element in X. For
        matrices, MAX(X) is a vector containing the maximum element
        from each column. [Y,I] = MAX(X) stores the indices of the
        maximum values in vector I. MAX(X,Y) returns a matrix the
        same size as X and Y with the largest elements taken from X
        or Y. When complex, the magnitude MAX(ABS(X)) is used.

>> max(b)
ans =
     9
>> [maxvalue,index]=max(b)
maxvalue =
     9
index =
     5

>> A=[3 5 4 2 0;7 6 3 5 1;8 6 9 2 0]
A =
     3     5     4     2     0
     7     6     3     5     1
```

```
     8     6     9     2     0
>> [maxv,index]=max(A)
maxv =
     8     6     9     5     1
index =
     3     2     3     2     2

>> help sum

 SUM    Sum of the elements.
        For vectors, SUM(X) is the sum of the elements of X.
        For matrices, SUM(X) is a row vector with the sum over
        each column. SUM(DIAG(X)) is the trace of X.

        See also PROD, CUMPROD, CUMSUM.

>> sum(A)
ans =
    18    17    16     9     1
>> sum(sum(A))
ans =
    61
```

Note also the element-by-element operations, for example, multiplication:

```
>> B=[1 1 1 1 1;0 0 0 0 0;2 2 2 2 2]
B =
     1     1     1     1     1
     0     0     0     0     0
     2     2     2     2     2
>> A.*B
ans =
     3     5     4     2     0
     0     0     0     0     0
    16    12    18     4     0
```

Laboratory Exercise 1

This lab is about linear programming. First you should run an available Matlab program that demonstrates the simplex method. You choose the pivot element on every iteration. Then you have to modify the program so that the pivot elements are chosen automatically. Finally you should test your program on some examples.

You should also make a version that produces the result as fast as possible, and investigate how the execution time increases with the number of constraints and variables.

Cycling can be prevented with *Bland's rule*:

1. If there are more than one variable to enter the basis (that is, columns with negative reduced costs), then choose the one with the lowest index.

2. If there are more than one variable to leave the basis, then choose the one with the lowest index. (Note: the lowest index of a *basis* variable.)

The problem to solve at this session is

$$\begin{array}{ll} \text{maximize} & z = \boldsymbol{c}^T \boldsymbol{x} \\ \text{subject to} & \boldsymbol{A}\boldsymbol{x} = \boldsymbol{b}, \quad \boldsymbol{x} \geq \boldsymbol{0}. \end{array}$$

It is assumed that a feasible choice of basic variables is given. This is easy if the problem has been converted from its standard form with $b \geq 0$. Why?

The simplex method in Matlab. The following procedure requires that the user assigns the correct variables to enter and leave on every iteration. Observe that the tableau is constructed with the function `checkbasic` from hand-in #1.

```
function [tableau,basicvars,steps]=simpmovie(A,b,c,basicvars)
% [y,basicvars,steps] = simpmovie(A,b,c,basicvars)
%
% A     m*n-matrix
% b     m*1-matrix, b>=0
% c     n*1-matrix
% basicvars 1*m - matrix with indices for feasible basic variables.
%
% Shows a movie of how the simplex method works
% on the problem
%                   max(c'x), when Ax=b, x>=0.
%

[m,n]=size(A);

% Create a tableau with slack variables
```

```
[tableau,xb,basic,feasible,optimal]=checkbasic(A,b,c,basicvars);

steps=0;
% Loop until all reduced costs are non-positive
while min(tableau(m+1,1:n)) < -1e-6
        steps = steps+1;

        clc
        disp(tableau)
        basicvars

        % Input variables to enter and leave
        q=input('entering␣variable␣q␣=␣');
        p=input('leaving␣variable␣␣p␣=␣');

        % Update basic vars
        basicvars = union(basicvars, q);
        basicvars = setdiff(basicvars, p);

        % Compute new tableau with the new basic variables
        [tableau,xb,basic,feasible,optimal]=checkbasic(A,b,c,basicvars);

        if ~feasible
                disp(tableau)
                error('You␣have␣chosen␣an␣incorrect␣pivot␣element.␣Restart.');
        end
end

disp('');
disp(tableau)
disp('Congratulation!␣You␣have␣understood␣the␣simplex␣method.')
disp('(At␣least␣for␣this␣example.)')
```

PREPARATORY EXERCISE. Rewrite the program above and name it `simp(A,b,c,basicvars)`, so that the pivot elements are chosen automatically. The Matlab functions `min` and `find` are useful. Unbounded problems should be detected. The problem with cycling can be solved by using Bland's rule, but this is not compulsory. The number of steps of the simplex algorithms is returned from the function.

At the computer

1. Go to the subdirectory `lab1`.
2. Suitable test matrices of size $m \times n$ can be obtained with the procedure `init`:

```
A=9*rand(m,n)+1;
b=ones(m,1);
c=ones(n,1);
A=[A eye(m)];
c=[c;zeros(m,1)];
basicvars=(n+1):(m+n);
```

A trial run with a 3×3-matrix is thus done by:

```
>>m=3
>>n=3
>>init
>>[y,basicvars,steps]=simpmovie(A,b,c,basicvars)
```

3. To study the phenomenon cycling, write `chvatal` to create the matrices

```
% http://people.orie.cornell.edu/~miketodd/or630/SimplexCyclingExample.pdf
A=[0.5 -5.5 -2.5 9 1 0 0;
   0.5 -1.5 -0.5 1 0 1 0;
    1    0    0  0   0 0 1];
b=[0 0 1]';
c=[10 -57 -9 -24 0 0 0]';
basicvars=[5 6 7];
```

Choose the pivot elements according to “the most negative reduced cost should enter the basis” and “if more than one variable can leave the basis, then choose the one with the lowest index”. (The cycling indices are $(q, p) = (1, 5), (2, 6), (3, 1), (4, 2), (5, 3), (6, 4)$.)

4. Make the copy

```
>>!cp simpmovie.m simp.m        (Unix)
>>!copy simpmovie.m simp.m      (Windows)
```

and adjust `simp.m` according to the preparatory exercise. Test your program on Example 1 in section 2.1 of the book (tableau 2.1, 2.3 and 2.4). Test random matrices (created by `init`) of sizes from 10×10 to 100×100 and make a table of the elapsed time. How does the execution time seem to depend on m and n? How does the number of steps grow with m and n? Exponential or polynomial? If you have solved the problem with cycling, then try the Chvatal example, too.

5. For every rule on how to choose the pivot element it is possible to construct examples that make the simplex method very slow. Take a look at `simpb` (write `edit simpb` from matlab) which is a simplex method using a ‘Bland’-like rule for selecting the incoming

basic variable (q). The script `badexample` creates a difficult example for this simplex algorithm. The variable d sets the size of the matrices A of the problem.

```
>>d=3
>>badexample
>>[y,basicvars,steps]=simpb(A,b,c,basicvars)
```

How does the number of steps grow with d? Is simplex a polynomial time solution for the linear programming problem?

Laboratory Exercise 2

This lab deals with the transportation problem, the maximal flow problem, the local search method and the branch and bound method.

You should write a Matlab program that solves the transportation problem and then test on some problems.

By means of an available Matlab program you should solve a maximal flow problem for a network with 23 nodes. You should also verify that "max-flow"="min-cut."

Finally, you should test the local search algorithm and the branch and bound algorithm on two test problems: the travelling salesman problem and the VigenÃÍre crypto analysis.

The transportation problem

PREPARATORY EXERCISE. Write a Matlab program with the feature

```
function [x,cost]=transport(s,d,c);
% [x,cost]=transport(s,d,c)
% Input:
%   s = supplies         (m*1)
%   d = demands          (n*1)
%   c = costs            (m*n)
% Output
%   x = optimal solution (m*n)
%   cost = minimal transport cost
...
```

Use the following subroutines, which you can find in your subdirectory `lab2`.

```
function [x,b]=northwest(s,d)
% [x,b]=northwest(s,d)
% x: shipments using nw-rule (m*n)
% b: 1 for each basic variables 0 for nonbasic (m*n)
% s: supplies (m*1)
% d: demands (n*1)
if (sum(s)~=sum(d)),
  disp('ERROR:␣The␣total␣supply␣is␣not␣equal␣to␣the␣total␣demand.');
  return;
end
m=length(s);
n=length(d);
i=1;
j=1;
x=zeros(m,n);
b=zeros(m,n);
```

```
while ((i<=m) & (j<=n))
   if s(i)<d(j)
      x(i,j)=s(i);
      b(i,j)=1;
      d(j)=d(j)-s(i);
      i=i+1;
    else
      x(i,j)=d(j);
      b(i,j)=1;
      s(i)=s(i)-d(j);
      j=j+1;
    end
end
```

```
function [u,v]=multipliers(x,c,b)
% [u,v]=multipliers(x,c,b)
% x: current solution  (m*n)
% b: 1 for each basic variables 0 for nonbasic (m*n)
% c: costs (m*n)
% u: lagrange multipliers for rows (m*1)
% v: lagrange multipliers for columns (n*1)
[m,n]=size(x);
if sum(sum(b))< m+n-1
  disp('Error_in_multipliers')
  return
else
  u=Inf*ones(m,1);
  v=Inf*ones(n,1);
  u(1)=0;   % choose an arbitrary multiplier = 0
  nr=1;
  while nr<m+n  % until all multipliers are assigned
    for row=1:m
      for col=1:n
        if b(row,col)>0
          if (u(row)~=Inf) & (v(col)==Inf)
            v(col)=c(row,col)-u(row);
            nr=nr+1;
          elseif (u(row)==Inf) & (v(col)~=Inf)
            u(row)=c(row,col)-v(col);
            nr=nr+1;
          end
        end
      end
```

```
    end
  end
end
```

```
function [y,bout]=cycle(x,row,col,b)
% [y,bout]=cycle(x,row,col)
% x: current solution (m*n)
% b: entering basic variables (m*n)
% row,col: index for element entering basis
% y: solution after cycle of change (m*n)
% bout: new basic variables after cycle of change (m*n)
bout=b;
y=x;
[m,n]=size(x);
loop=[row col]; % describes the cycle of change
x(row,col)=Inf; % do not include (row,col) in the search
b(row,col)=Inf;
rowsearch=1;    % start searching in the same row
while (loop(1,1)~=row | loop(1,2)~=col | length(loop)==2),
  if rowsearch, % search in row
    j=1;
    while rowsearch
       if (b(loop(1,1),j)~=0) & (j~=loop(1,2))
         loop=[loop(1,1) j ;loop]; % add indices of found element to loop
         rowsearch=0;  % start searching in columns
       elseif j==n,    % no interesting element in this row
         b(loop(1,1),loop(1,2))=0;
         loop=loop(2:length(loop),:); % backtrack
         rowsearch=0;
       else
         j=j+1;
       end
    end
  else  % column search
    i=1;
    while ~rowsearch
       if (b(i,loop(1,2))~=0) & (i~=loop(1,1))
         loop=[i loop(1,2) ; loop];
         rowsearch=1;
       elseif i==m
         b(loop(1,1),loop(1,2))=0;
         loop=loop(2:length(loop),:);
         rowsearch=1;
```

```
      else
        i=i+1;
      end
    end
  end
end
% compute maximal loop shipment
l=length(loop);
theta=Inf;
minindex=Inf;
for i=2:2:l
 if x(loop(i,1),loop(i,2))<theta,
  theta=x(loop(i,1),loop(i,2));
  minindex=i;
 end;
end
% compute new transport matrix
y(row,col)=theta;
for i=2:l-1
  y(loop(i,1),loop(i,2))=y(loop(i,1),loop(i,2))+(-1)^(i-1)*theta;
end
bout(row,col)=1;
bout(loop(minindex,1),loop(minindex,2))=0;
```

You can test your code using the files example511, example512 and example513. The solutions are

```
>> example511; [x cost] = transport(s,d,c)
x =

   100     0    20     0
     0    60    60    20
     0     0     0   100
cost =
        1900

>> example512; [x cost] = transport(s,d,c)
x =
     0     0     0    30    70
    20    60    80     0     0
    70     0     0    70     0
cost =
        1930

```

```
>> example513; [x cost] = transport(s,d,c)
x =
     0     0     0    30    70
    20    60    80     0     0
    30     0     0    70     0
cost =
        1730
```

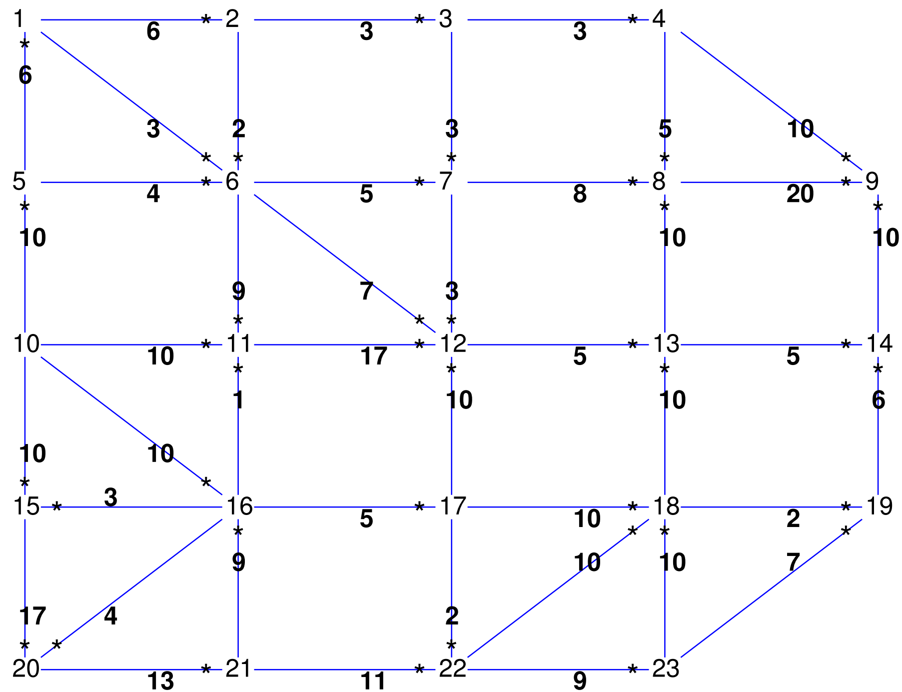

Figure 1: Matlab's graph of the network. The capacities of the arcs at the start are shown with bold figures and the asterisks symbolizes arrows, e.g. 6 units can be shipped from node 1 to node 2.

The maximal flow problem

In this session you will find the maximal flow in the network in Figure 1 *from node 10 to node 9*. When you have found a new breakthrough for the flow and input it, the graph will be updated automatically and the asterisks (=arrows) will show in which direction of the arcs more flow can be shipped. When you think you have found the optimal solution you should divide the nodes according to the max flow-min cut theorem.

Local search and Branch-and-bound

Study the local search (Steepest descent) method and the branch and bound method on two problems. In the directory `lab2` are two subdirectories `@vigcrypto` and `@tsp`. These contain methods for two new objects: vigcrypot and tsp objects. One can create new objects of these types using

```
>> problem = demoproblem(tsp);
>> problem = demoproblem(vigcrypto);
```

One can also construct other instances of the problem classes above using the constructors

```
>> problem = tsp(relevantdata);
>> problem = vigcrypto(relevandata);
```

To each of these object a number of methods are given. For example

```
x=randomindomain(problem);
```

generates a representative x for a point in the domain of the combinatorial optimization problem. In general the points x are represented as row matrices.

```
f=evaluate(problem,x);
```

evaluates the goal function f at the point x in the domain of the combinatorial optimization problem.

```
xlist = getneighbours(problem,x);
```

generates a list xlist of all neighbours to the point x in the domain of the optimization problem. Each row of the matrix xlist is a representative of a point (a neighbour) close to x.

```
D = getdomain(problem);
```

generates a representative D of the whole domain of the problem. Subsets are also represented as a row matrix.

```
[listofsubsets,sizes]=branch(problem,S)
```

generates a list of representatives of subsets to the set S.

```
[fl,f,fu]=bound(problem,subset);
```

calculates upper fu and lower fl bounds on the optimal value of the function f in the subset. More information can (hopefully) be found in `Contents.m` and in the comments in each file. Try for example

```
help lab2
help tsp
methods tsp
help tsp/evaluate
help branchandbound
```

At the computer

1. Download the Matlab files needed for the session from the course home page. Copy this file to your home directory, decompress and unpack it. Now you have a subdirectory `lab2`. Start Matlab, go to this directory, and make sure the Matlab search path includes the directory `lab2`. This can be done in matlab using the path command.

```
>> path(path,pwd);
```

   Information about the routines in the directory can be found in the text file `Contents.m`

2. Run the program `transportmovie.m` on Example **1** of section 5.1 in the book with the following commands

```
>> example511
>> transportmovie(s,d,c)
```

   Compare the result with your own calculations by hand! Get the matrices $\boldsymbol{s}$, $\boldsymbol{d}$ and $\boldsymbol{c}$ from the m-files `example511.m`, `example512.m` and `example513.m`.

3. Test your program `transport.m` with data from the examples.

4. Solve the maximal flow problem by running `maxflow`. Mark the solution in Figure 1. Divide the nodes by a minimal cut into two groups M och M' such that the source belongs to M and the sink to M'. (Make this division by hand with help of the result from the run).

5. Copy the code in `steepdescstep.pre.m` to `steepdescstep.m`. Modify the code so that the function returns the neighbour with the lowest value on the goal function and a boolean `lokmin` which indicates whether `xin` is a local minimum to the problem.

```
function [xout,lokmin]=steepdescstep(problem,xin)
% One step in local search
% Input:
%   problem - The optimization problem.
%   xin     - a point in the domain of problem.
% Output:
%   xout    - the neighbour to xin with the lowest
%             goal function.
%   lokmin  - 1 if xin is a local minimum, 0 otherwise

% Generate all neighbours to xin.
        neighbours=getneighbours(problem,xin);
        lokmin=1;
        f=evaluate(problem,xin);
```

```
        xout=xin;
        for ii=1:size(neighbours,1);
                y=neighbours(ii,:);
                fnew=evaluate(problem,y);
                ...
                        add apropriate code here.
                ...
        end
end
```

Now try the routine `steepdesc` (which uses your steepdescstep routine) on both a travelling salesman problem and the vigcrypto problem. Generate problem objects with

```
>> problem = demoproblem(tsp);
>> problem = demoproblem(vigcrypto);
```

Generate starting points for steepest descent with

```
xin=randomindomain(problem);
```

Try with different random starting points. Do the routine end up in different local minima or most often in the same?

```
function [xout,steps,locmin]=steepdesc(problem,xin)
        if nargin < 2,
                xin = randomindomain(problem);
        end;

        steps=0;
        locmin=0;
        xout=xin;
        while ~locmin,
                [xout,locmin]=steepdescstep(problem,xout);
                steps=steps+1;
%               f = evaluate(problem,xout)
        end
end
```

6. Type

```
>> edit branchandbound
```

to see the code for the branch and bound algorithm. Try the algorithm

```
>> [dmin,fumin,res]=branchandbound(problem);
```

on both a travelling salesman problem and on the vigcrypto problem. Did the local search find the global minimum?