

# nRF51 Series Reference Manual

Version 2.0

The nRF51 series offers a range of ultra-low power System on Chip solutions for your 2.4 GHz wireless products. With the nRF51 series you have a diverse selection of devices including those with embedded *Bluetooth*<sup>®</sup> low energy and/or ANT<sup>™</sup> protocol stacks as well as open devices enabling you to develop your own proprietary wireless stack and ecosystem.

The nRF51 series combines Nordic Semiconductor's leading 2.4 GHz transceiver technology with a powerful but low power ARM<sup>®</sup> Cortex<sup>™</sup>-M0 core, a range of peripherals and memory options. The pin and code compatible devices of the nRF51 series offer you the most flexible platform for all your 2.4 GHz wireless applications.



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# **Revision History**

Date	Version	Description
October 2013	2.0	<ul> <li>New sections and chapters added:</li> <li>Added Section 6.1.1 "Override parameters" on page 18.</li> <li>Added Section 6.2.15 "OVERRIDDEN" on page 23.</li> <li>Added Section 6.2.16 "NRF_1MBIT[n] (n=0.4)" on page 23.</li> <li>Added Section 16.1.7 "CRC" on page 76.</li> <li>Added Section 16.2.29 "BCC" on page 94.</li> <li>Added Section 16.2.30 "OVERRIDE[n] (n=03)" on page 94.</li> <li>Added Section 16.2.31 "OVERRIDE[n] (n=03)" on page 94.</li> <li>Added Chapter 26 "SPI Slave (SPIS)".</li> <li>Added Chapter 26 "SPI Slave (SPIS)".</li> <li>Added Chapter 31 "Low Power Comparator (LPCOMP)" on page 181.</li> </ul> Sections updated: <ul> <li>Updated Block diagram, Figure 1 on page 8.</li> <li>Updated Section 6.2 "Registers" on page 19.</li> <li>Updated Section 5.1.1 "Writing to the NVM" on page 14 (added last sentence).</li> <li>Updated Section 7.2 "Registers" on page 24.</li> <li>Updated Section 8.2.1 "PERR0" on page 30.</li> <li>Updated Section 8.2.1 "PERR0" on page 32.</li> <li>Updated Section 8.2.4 "PROTENSETO" on page 33.</li> <li>Updated Section 11.1.2 "System OFF mode" on page 45.</li> <li>Updated Section 11.1.2 "System OFF mode" on page 48.</li> <li>Updated Section 16.1.1 "EasyDMA" on page 73.</li> <li>Updated Section 16.1.5 "Received Signal Strength Indicator (RSSI)" on page 75.</li> <li>Updated Section 16.1.6 "Register" on page 83.</li> <li>Updated Section 16.1.6 "Register" on page 83.</li> <li>Updated Section 16.1.7 "RayDMA" on page 103.</li> <li>Updated Section 16.1.6 "Register" on page 83.</li> <li>Updated Section 16.1.7 "RayDMA" on page 73.</li> <li>Updated Section 16.1.7 "Register" on page 83.</li> <li>Updated Section 16.1.7 "Register" on page 83.</li> <li>Updated Section 16.1.7 "Register" on page 83.</li> <li>Updated Section 16.1.6 "Register" on page 83.</li> <li>Updated Section 16.1.7 "Register" on page 83.</li> <li>Updated Section 18.1.6 "Event Control feature" on page 103.</li> <li>Updated Section 23.1.7 "EasyDMA and ERROR event" on page 129.</li> <li>Updated Section 23.1.7 "EasyD</li></ul>
March 2013	1.1	<ul> <li>Updated DC/DC converter setup description in Section 11.1.1.1 "DC/DC converter setup" on page 40.</li> <li>Updated values in Section 16.1.9 "Maximum consecutive transmission time" on page 77.</li> </ul>
January 2013	1.0	First release



# 1 About this document

This reference manual is a functional description of all the modules and peripherals supported by the nRF51 series and subsequently, is a common document for all nRF51 System on Chip (SoC) devices.

**Note:** nRF51 SoC devices may not support all the modules and peripherals described in this document and some of their implemented modules may have a reduced feature set. Please refer to the individual nRF51 device product specification for details on the supported feature set, electrical and mechanical specifications, and application specific information.

## 1.1 Writing conventions

This Reference Manual follows a set of typographic rules to ensure that the document is consistent and easy to read. The following writing conventions are used:

- Command and event names, and bit state conditions are written in Lucida Console.
- Pin names and pin signal conditions are written in Consolas.
- File names and User Interface components are written in **bold**.
- Internal cross references are italicized and written in *semi-bold*.
- Placeholders for parameters are written in *italic regular font*. For example, a syntax description of SetChannelPeriod will be written as: SetChannelPeriod(*ChannelNumber*, *MessagingPeriod*).
- Fixed parameters are written in regular text font. For example, a syntax description of SetChannelPeriod will be written as: SetChannelPeriod (0, Period).

#### 1.1.1 Peripheral naming and abbreviations

Every peripheral has a unique name or an abbreviation constructed by a single word, e.g. TIMER. This name is indicated in parentheses in the peripheral chapter heading. This name will be used in CMSIS to identify the peripheral.

The peripheral instance name, which is different from the peripheral name, is constructed using the peripheral name followed by a numbered postfix, starting with 0, for example, TIMER0. A postfix is normally only used if a peripheral can be instantiated more than once. The peripheral instance name is also used in the CMSIS to identify the peripheral instance.



### 1.1.2 Register tables

Individual registers are described using register tables. These tables are built up of two sections. The first three rows, which are shaded blue, describe the position and size of the different fields in the register. The following rows, beginning with the row shaded green, describes the fields in more detail.

Bit num	per		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (Field	ID)		0	1	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	A A	A
Reset va	lue		0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	) (	0
ID RW	Field	Value ID	Va	lue			De	scri	ptic	on																							
A RW							Re	egist	ter v	vith	a sir	ngle	fiel	d w	ithc	out e	enur	mera	ated	l val	ues												

Table 1 Example of a register table with a single field

#### 1.1.2.1 Fields and values

The ID (Field ID) row specifies which bits that belong to the different fields in the register.

The **ID** (Field ID) may also specify constants. '1' in this row means that the associated bit is read as '1' and must be written as '1'. Similarly, '0' means that the associated bit is read as '0' and must be written as '0'. A "-" means that the field is reserved and that it is read as undefined and must be written as '0' to secure forward compatibility. If a register is divided into more than one field, a unique field name is specified for each field in the **Field** column.

If a field has enumerated values, then every value will be identified with a unique value ID in the **Value ID** column. Single-bit bit-fields may however omit the "Value ID" when values can be substituted with a Boolean type enumerator range, for example, True, False; Disable, Enable, and On, Off, and so on.

The **Value** column can be populated in the following ways:

- Individual enumerated values, for example, 1, 3, 9.
- Range values, e.g. [0..4], that is, all values from and including 0 and 4.
- Implicit values. If no values are indicated in the **Value** column, all bit combinations are supported, or alternatively the field's translation and limitations are described in the text instead.



If two or more fields are closely related, the value ID, value, and description may be omitted for all but the first field. Subsequent fields will indicate inheritance with "..".

Bit	num	ber		31	30	29	28	3 27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	76	5 5	4	3	2	1	0
ID	(Field	ID)		0	1	-	-	-	-	-	-	-	-	-	-	-	-	-	-	F	Е	D	c	-	-	-	- 1	BE	3 B	В	A	A	A	A
Re	set va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	0	0	0	0	0
ID	RW	Field	Value ID	Va	lue			De	scri	ptic	on																							
А	RW	RANGE		[0	15	]		Ra	ange	e de	scrip	otior	use	es th	nis s	ynta	x																	
В	RW	ENUM						Fi	elds	wit	h en	ume	erate	ed v	alue	es ar	e de	escri	ibed	l like	e thi	s												
			ENUMA	0				Fi	rst e	nun	nera	ted	valu	e																				
			ENUMB	4				Se	ecor	nd er	num	erat	ed v	/alu	e																			
			ENUMC	15	5			Tł	nird	enu	mer	ated	val	ue																				
С	RW	IMP0							-	ter fo						th ir	npli	cit e	enur	nera	ted	val	ues,	, act	ing	as p	oare	ent	for					
				0				Di	isab	le																								
				1				Er	nabl	e																								
D	RW	IMP1																																
Е	RW	IMP2																																
F	RW	IMP3																																

 Table 2
 Example of a register table with multiple fields



# 2 System overview

The nRF51 series of System on Chip (SoC) devices embed a powerful yet low power ARM<sup>®</sup> Cortex<sup>™</sup>-M0 processor with our industry leading 2.4 GHz RF transceivers. In combination with the very flexible orthogonal power management system and a Programmable Peripheral Interconnect (PPI) event system, the nRF51 series enables you to make ultra-low power wireless solutions.

The nRF51 series offers pin compatible device options for *Bluetooth* low energy, proprietary 2.4 GHz, and ANT<sup>™</sup> solutions giving you the freedom to develop your wireless system using the technology that suits your application the best. Our unique memory and hardware resource protection system allows you to develop applications on devices with embedded protocol stacks running on the same processor without any need to link in the stack or strenuous testing to avoid application and stack from interfering with each other.



## 2.1 Block diagram

*Figure 1* illustrates the overall system. Arrows with white heads indicate signals that share physical pins with other signals.

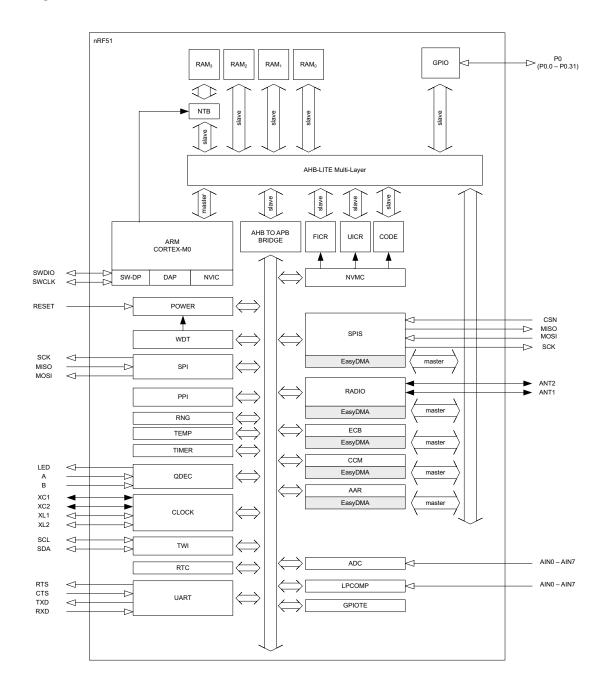


Figure 1 Block diagram



# 2.2 System blocks

#### 2.2.1 ARM<sup>®</sup> Cortex<sup>™</sup>-M0

A low power ARM<sup>®</sup> Cortex<sup>TM</sup>-M0 32 bit CPU is embedded in all nRF51 series devices. The ARM<sup>®</sup> Cortex<sup>TM</sup>-M0 has a 16 bit instruction set with 32 bit extensions (**Thumb-2**<sup>®</sup> **technology**) that delivers high density code with a small memory footprint. By using a single-cycle 32 bit multiplier, a 3-stage pipeline, and a Nested Vector Interrupt Controller (NVIC), the ARM<sup>®</sup> Cortex<sup>TM</sup>-M0 CPU makes program execution simple and highly efficient.

The ARM<sup>®</sup> Cortex Microcontroller Software Interface Standard (CMSIS) hardware abstraction layer for the ARM<sup>®</sup> Cortex-M processor series is implemented and available for M0 CPU. Code is forward compatible with ARM<sup>®</sup> Cortex-M3 based devices.

#### 2.2.2 2.4 GHz radio

The nRF51 series ultra-low power 2.4 GHz GFSK RF transceiver is designed and optimized to operate in the worldwide ISM frequency band at 2.400 GHz to 2.4835 GHz. Configurable radio modulation modes and packet structure makes the transceiver interoperable with *Bluetooth* low energy (BLE), ANT<sup>™</sup>, Gazell, Enhanced Shockburst<sup>™</sup>, and a range of other 2.4 GHz protocol implementations.

The transceiver receives and transmits data directly to and from system memory. It is stored in clear text even when encryption is enabled, so packet data management is flexible and efficient.

#### 2.2.3 Power management

The nRF51 series power management system is orthogonal and highly flexible with only simple ON or OFF modes governing a whole device. In System OFF mode, everything is powered down but sections of the RAM can be retained. The device state can be changed to System ON through reset or wake up from all GPIOs. When in System ON mode, all functional blocks are accessible with each functional block remaining in IDLE mode and only entering RUN mode when required.

#### 2.2.4 PPI system

The Programmable Peripheral Interconnect (PPI) enables different peripherals to interact autonomously with each other using tasks and events without use of the CPU. The PPI provides a mechanism to automatically trigger a task in one peripheral as a result of an event occurring in another. A task is connected to an event through a PPI channel.

#### 2.2.5 Debugger support

The 2 pin Serial Wire Debug interface (provided as a part of the Debug Access Port, DAP) offers in conjunction with the Basic Branch Buffer (BBB) a flexible and powerful mechanism for non-intrusive program code debugging. This includes adding breakpoints in the code, performing single stepping, and capturing instruction trace of parts of the code execution flow.

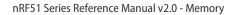


# 3 CPU

A low power ARM<sup>®</sup> Cortex<sup>TM</sup>-M0 32 bit CPU is embedded in all nRF51 series devices. The ARM<sup>®</sup> Cortex<sup>TM</sup>-M0 has a 16 bit instruction set with 32 bit extensions (**Thumb-2**<sup>®</sup> **technology**) that delivers high density code with a small memory footprint. By using a single-cycle 32 bit multiplier, a 3-stage pipeline, and a Nested Vector Interrupt Controller (NVIC), the ARM<sup>®</sup> Cortex<sup>TM</sup>-M0 CPU makes program execution simple and highly efficient.

The ARM<sup>®</sup> Cortex<sup>TM</sup> Microcontroller Software Interface Standard (CMSIS) hardware abstraction layer for the ARM<sup>®</sup> Cortex<sup>TM</sup>-M processor series is implemented and available for M0 CPU. Code is forward compatible with ARM<sup>®</sup> Cortex<sup>TM</sup>-M3 based devices.

For further information on the embedded ARM<sup>®</sup> Cortex<sup>™</sup>-M0 CPU, please refer to www.arm.com/products/ processors/cortex-m/cortex-m0.php.





# 4 Memory

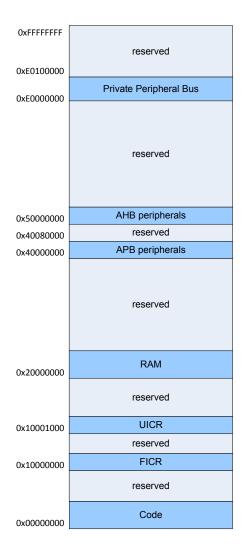


Figure 2 Memory map

## 4.1 Functional description

All memory blocks and registers in the nRF51 series are placed in a common memory map.

#### 4.1.1 Memory categories

There are three main categories of memory:

- Code memory
- Random Access Memory (RAM)
- Peripheral registers (PER)

In addition, there is one information block (FICR) containing read only parameters describing configuration details of the device and another information block (UICR) that can be configured by the user.



#### 4.1.2 Memory types

The various memory categories can have one of the following memory types:

- Volatile memory (VM)
- Non-volatile memory (NVM)

Volatile memory is a type of memory that will lose its contents when the chip loses power. This memory type can be read/written an unlimited number of times by the CPU.

Non-volatile memory is a type of memory that can retain stored information even when the chip loses power. This memory type can be read an unlimited number of times by the CPU, but have restrictions on the number of times it can be written and also on how it can be written. Writing to non-volatile memory is managed by the Non Volatile Memory Controller (NVMC).

#### 4.1.3 Code memory

The code memory is normally used for storing the program run by the CPU, but can also be used for storing data constants that are retained when the chip loses power.

The code memory is non-volatile.

#### 4.1.4 Random Access Memory (RAM)

RAM is normally used by the CPU program for temporary data storage, but it is also possible to run a CPU program from RAM.

The RAM is divided into multiple RAM blocks that can be further divided into multiple RAM sections, that is, a RAM block may contain one or more RAM sections. More information about how these RAM blocks are divided can be found in the device specific product specifications.

The RAM is located from address 0x20000000 in the address space.

The RAM is volatile and always loses its contents when the chip loses power. The RAM may also lose its contents when entering System OFF power saving mode. Whether the contents of a particular RAM block are lost in System OFF power saving mode is dependent on the settings in the RAMON register in the POWER peripheral.

#### 4.1.5 Peripheral registers

The peripheral registers are registers used for interfacing to peripheral units such as timers, the radio, the ADC, and so on.





## 4.2 Instantiation

The nRF51 series peripheral instantiation is shown in the table below.

ID	Base address	Peripheral	Instance	Description
0	0x40000000	POWER	POWER	Power control
0	0x40000000	CLOCK	CLOCK	Clock control
1	0x40000000 0x40001000	RADIO	RADIO	2.4 GHz Radio
2	0x40002000	UART	UARTO	Universal Asynchronous Receiver/Transmitter 0
2	0x40002000	SPI	SPIO	Serial Peripheral Interface
3	0x40003000	TWI	TWIO	I2C compatible Two-Wire Interface
4	0x40004000	SPI	SPI1	Serial Peripheral Interface
4	0x40004000	TWI	TWI1	·
				I <sup>2</sup> C compatible Two-Wire Interface 1
4	0x40004000	SPIS	SPIS1	SPI slave
6	0x40006000	GPIOTE	GPIOTE	GPIO Tasks and Events
7	0x40007000	ADC	ADC	Analog-to-Digital Converter Timer/counter 0
8	0x40008000	TIMER	TIMER0	
9	0x40009000	TIMER	TIMER1	Timer/counter 1
10	0x4000A000 0x4000B000		TIMER2	Timer/counter 2 Real Time Counter 0
11		RTC	RTCO	
12 13	0x4000C000 0x4000D000	TEMP	TEMP	Temperature sensor
			ECB	Random number generator
14 15	0x4000E000 0x4000F000	ECB CCM	ССМ	Crypto ECB
15	0x4000F000	AAR	AAR	AES CCM mode encryption Accelerated Address Resolver
	0x4000F000	WDT	WDT	Watchdog timer
16 17		RTC	RTC1	Real Time Counter 1
17	0x40011000 0x40012000	QDEC	QDEC	Quadrature Decoder
19	0x40012000	LPCOMP	LPCOMP	Low power comparator
20	0x40013000	LPCOIVIP	LFCOMF	
 29				
30	0x4001E000	NVMC	NVMC	Non-volatile memory controller
31	0x4001F000	PPI	PPI	Programmable Peripheral Interconnect
NA	0x50000000	GPIO	PO	General purpose input and output
NA	0x10000000	FICR	FICR	Factory Information Configuration Registers
NA	0x10001000	UICR	UICR	User Information Configuration Registers
				ripheral instantiation



# 5 Non-Volatile Memory Controller (NVMC)

The Non-volatile Memory Controller (NVMC) is used for writing and erasing Non-volatile Memory (NVM).

Before a write can be performed the NVM must be enabled for writing in CONFIG.WEN. Similarly, before an erase can be performed the NVM must be enabled for erasing in CONFIG.EEN. The user must make sure that writing and erasing is not enabled at the same time, failing to do so may result in unpredictable behavior.

## 5.1 Functional description

#### 5.1.1 Writing to the NVM

When writing is enabled, the NVM is written by writing a word to a word aligned address in the CODE or UICR. The NVMC is only able to write bits in the NVM that are erased, that is, set to '1'.

The time it takes to write a word to the NVM is specified by t<sub>WRITE</sub> in the product specification. The CPU is halted while the NVMC is writing to the NVM.

Only word aligned writes are allowed. Byte or half word aligned writes will result in a hard fault.

#### 5.1.2 Writing to User Information Configuration Registers

UICR registers are written as ordinary non-volatile memory. After the UICR has been written, the new UICR configuration will only take effect after a reset.

#### 5.1.3 Erasing User Information Configuration Registers

There are two registers that can be used for erasing the UICR, the ERASEALL and the ERASEUICR. When readback protection is configured, writing the ERASEUICR register only has an effect when the entire region 1 of code memory is erased (all '1's).

The time it takes to perform an ERASEUICR command is specified by t<sub>PAGEERASE</sub> in the product specification. The CPU is halted while the NVMC performs the erase operation.

#### 5.1.4 Erase all

When erase is enabled, the whole CODE and UICR can be erased in one operation by using the ERASEALL register. ERASEALL will not erase the Factory Information Configuration Registers (FICR).

The time it takes to perform an ERASEALL operation is specified by t<sub>ERASEALL</sub> in the product specification. The CPU is halted while the NVMC performs the erase operation.

#### 5.1.5 Erasing a page in code region 1

When erase is enabled, the NVM can be erased page by page using the ERASEPAGE register or the ERASEPCR1 register. After erasing a NVM page all bits in the page are set to '1'. The time it takes to erase a page is specified by t<sub>PAGEERASE</sub> in the product specification. The CPU is halted while the NVMC performs the erase operation.



#### 5.1.6 Erasing a page in code region 0

ERASEPCR0 is used to erase a page in code region 0. The ERASEPCR0 register can <mark>only be accessed from a</mark> program running in code region 0.

To enable non-volatile storage for program running in code region 0, it is possible for this program to erase and re-write any code page it designates for this purpose within code region 0. The ERASEPCR0 can be used for this purpose. The ERASEPCR0 register has a restriction on its use, enforced by the MPU, where only code running from code region 0 can write to it. It is possible for a program running from code region 0 to erase a page in code region 1 using ERASEPCR1.

The time it takes to erase a page is specified by  $t_{PAGEERASE}$  in the product specification.

# 5.1.7 Availability of erase operations based on presence of pre-programmed factory code

To enable independent readback protection of code region 0 and code region 1, and to ensure that a developer cannot erase factory-programmed code from region 0, the ERASEALL and ERASEUICR registers are not always available. The following table shows when the registers can and cannot be used. When a register is disabled, attempting to write it will have no effect.

FICR.PPFC <sup>1</sup>	Available operations	Unavailable registers
0XFF (Pre-programmed code not present)	ERASEALL	ERASEUICR
0x00 (Pre-programmed code present)	ERASEUICR	ERASEALL

1. See Chapter 6 "Factory Information Configuration Registers (FICR)" on page 18 for details.

 Table 4 Available erase operations dependent on presence of pre-programmed factory code

#### 5.2 Registers

Register	Offset	Description
REGISTERS		
READY	0x400	Ready flag
CONFIG	0x504	Configuration register
ERASEPAGE	0x508	Register for erasing a page in code region 1
ERASEPCR1	0x508	Register for erasing a page in code region 1. Equivalent to ERASEPAGE.
ERASEPCR0	0x510	Register for erasing a page in code region 0
ERASEALL	0x50C	Register for erasing all non-volatile user memory
ERASEUICR	0x514	Register for erasing User Information Configuration Registers

Table 5Register overview



## 5.2.1 READY

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID RW Field	Value ID	Va	lue			De	scri	ptio	on																							
A . D																																
AR																																
AK		0				N١	VMC	is k	ousy	ı (or	ו-go	ing	wri	te o	r era	ise c	oper	atic	on).													

## 5.2.2 CONFIG

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field	d ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		- /	A A
Reset va	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW	Field	Value ID	Va	lue			De	scri	ipti	on																							
A RW	WEN														ode en tl							nen	ded	to	onl	y a	cti	vat	e				
		REN	0				Re	ead	onl	y ac	ces	s.																					
		WEN	1				W	rite	Ena	able	d.																						
		EEN	2				Er	ase	ena	able	d.																						

#### 5.2.3 ERASEALL

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8 3	7 (	6 5	5 4	¥ 3	2	1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-					-	-	-	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	0 0	) (	0 (	0	0	0
ID RW Field Value ID	Val	ue			De	scri	pti	on																							
A RW					en	abl	ed l	oy C	ON	FIG.		bef	-		-	-		regis le m								be					
	0				No	o op	era	tior	ı.																						
	1				St	art o	chip	era	ise.																						

#### 5.2.4 ERASEUICR

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9 8	B 7	6	5	4	3	2	1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-			-	-	-	-	-	-	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	0	0	0	0	0	0
ID RW Field Value ID	Va	lue			De	scri	ptio	on																							
A RW																		Conf pefo									at				
	0				No	о ор	era	tion																							
	1				St	art e	eras	e of	UIC	R.																					



## 5.2.5 ERASEPAGE/ERASEPCR1

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 6	i 5	4	3	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α	A	A	A	A /	A A		A	A	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	) (	0	0	0	0	0
ID RW Field	Value ID	Va	lue			De	scri	ptic	on																							
A RW						Th No	e va ote t	alue that	is t coc	he a le e	addr	ess mu	of tl	•	age	e to	be e	eras	ed (	add		es o fore					•	<i>.</i>				
						yo	u ai	re u	sing	. At	tem	pts	to e	rase	pa	ges	tha	t are	e ou	tsid	e th	code e co y be	de a	area	a m							

#### 5.2.6 ERASEPCR0

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	5 5	4	3	2	1	0
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A	A	A	A	A	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	0	0	0	0	0
ID RW Field Value ID	Val	ue			De	scri	ptic	on																							
A RW					Th Oi ge m W CC Se yc	nly p cess ener emo ritin ONF e pr ou ar	alue bage sed ateo ory i g to IG.E rodu re u	is the ad from d if t regio o ER/ EN h uct s sing	he a drei he i he i ASE has pec	iddr sses orog regi: PRC to b cifica	ess in c iran ster 0 fro e se ation pts	of the ode n run is at om t t to n for to e	ne p reg nnir tten the ena r inf rase	age jion ng ir npte Seri ible orm	e to l 0 ai n co ed a al W era: natic ges	be e re al de n cces /ire l se. on al that	t are	ed (a red. nory fro bug t th	0. addr This reg m a (SWI e tot tside	ion pro D) v tal c	giste 0. A grar vill h code e co	r cai har n in nave size de a	n o rd fa RA no e of area	nly ault M o eff	be t wi or c fect	ll b ode evic	e				



# 6 Factory Information Configuration Registers (FICR)

## 6.1 Functional description

Factory Information Configuration Registers are pre-programmed in factory and cannot be erased by the user. These registers contain chip specific information and configuration.

#### 6.1.1 Override parameters

Factory Information Configuration Registers contain override parameters set during device calibration in production, which need to replace default settings in the RADIO. Override parameters and the RADIO mode they are used for varies between nRF51 devices. Read the OVERRIDDEN register to determine if the FICR contains override parameters for the radio mode you are going to use. If the FICR contains override parameters, they must be copied to the radio OVERRIDE registers before enabling the radio in that mode.



# 6.2 Registers

Register	Offset	Description
CODEPAGESIZE	0x010	Code memory page size
CODESIZE	0x014	Code memory size
CLENRO	0x028	Length of code region 0 in bytes
PPFC	0x02C	Pre-programmed factory code present
NUMRAMBLOCK	0x034	Number of individually controllable RAM blocks
SIZERAMBLOCK[0]	0x038	Size of RAM block 0 in bytes
SIZERAMBLOCK[1]	0x03C	Size of RAM block 1 in bytes
SIZERAMBLOCK[2]	0x040	Size of RAM block 2 in bytes
SIZERAMBLOCK[3]	0x044	Size of RAM block 3 in bytes
CONFIGID	0x05C	Configuration identifier
DEVICEID[0]	0x060	Device identifier, bit 31-0
DEVICEID[1]	0x064	Device identifier, bit 63-32
ER[0]	0x080	Encryption root, bit 31-0
ER[1]	0x084	Encryption root, bit 63-32
ER[2]	0x088	Encryption root, bit 95-64
ER[3]	0x08C	Encryption root, bit 127-96
IR[0]	0x090	Identity root, bit 31-0
IR[1]	0x094	Identity root, bit 63-32
IR[2]	0x098	Identity root, bit 95-64
IR[3]	0x09c	Identity root, bit 127-96
DEVICEADDRTYPE	0x0A0	Device address type
DEVICEADDR[0]	0x0A4	Device address, bit 31-0
DEVICEADDR[1]	0x0A8	Device address, bit 47-32
OVERRIDDEN	0x0AC	Override enable
BLE_1MBIT[0]	0x0EC	RADIO.OVERRIDE[0] values for BLE_1MBIT mode
BLE_1MBIT[1]	0x0F0	RADIO.OVERRIDE[1] values for BLE_1MBIT mode
BLE_1MBIT[2]	0x0F4	RADIO.OVERRIDE[2] values for BLE_1MBIT mode
BLE_1MBIT[3]	0x0F8	RADIO.OVERRIDE[3] values for BLE_1MBIT mode
BLE_1MBIT[4]	0x0FC	RADIO.OVERRIDE[4] values for BLE_1MBIT mode

Table 6 Register overview



#### 6.2.1 CODEPAGESIZE

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6 !	5	4	3	2	10
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	Α	A	A	Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A A
Value after eras	se	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11
ID RW Field	Value ID	Va	lue			De	scri	ptic	on																							
A R						Co	de i	mer	nory	/ pa	ge s	size	in b	ytes	;																	

## 6.2.2 CODESIZE

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	10
ID (Field ID)       A A A A A A A A A A A A A A A A A A A														A A	A A																	
															1	1 1	11															
ID RW Field	/alue after erase 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1																															
A R									nor iESIZ	·				er of	pag	ges.	Tota	al co	ode	spa	ce is	5:										

#### 6.2.3 CLENR0

Bit number	31 3	0 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	55	5 4	4 3	3 3	2	10
ID (Field ID)	A A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	۹ <i>۱</i>	A A	4 <i>i</i>	•	A	A A
Value after erase	1 1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	11	11	1 1	1	1	11
ID RW Field Value ID	Value	e		De	scri	pti	on																							
A R	[0N	1]		by fa N	engt /tes ctor (ma nis re	(FIC 'y co ix va	CR.C ode alue	OD is p e) is	EPA rese (FIC	GES ent o R.CO	IZE) on tl DDE	. Thi ne c PAC	is re hip, iESIZ	gist see ZE*	ter is PP FICI	s on FC. R.CC	ly u DDE	sed SIZE	wh -1,	en p but	ore-	pro	gra	amn	nec	d				
	0xFF	FFFF	FF	Va	alue	if tl	here	e is r	no p	ore-p	orog	ram	me	d co	ode	in tl	he c	hip.	Int	erpr	rete	d a:	s 0.							

## 6.2.4 PPFC

Bit number	31 30 29 2	27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		
Value after erase	1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
ID RW Field Value ID	Value	Description
A R ADDR	0x00 0xFF	Pre-programmed factory code present. Present Not present



#### 6.2.5 NUMRAMBLOCK

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6 !	5	4	3	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Value after erase		1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
ID RW Field	Value ID		Va	lue											D	esc	ript	ion															
A R						Nu	umb	er c	ofin	divi	dua	lly c	ont	rolla	ble	RA	Mb	lock	s.														

# 6.2.6 SIZERAMBLOCK[n] (n=0..3)

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	5 !	5 4	4 3	3 2	2 1	1 0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	4 /	۹ ۱	4 <i>4</i>	4 <i>4</i>	۱, r	A A
Value af	ter erase		1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 '	1 1	1	1 1	1 1	1 1	1
ID RW	Field	Value ID		Val	lue											D	esc	ript	ion														
AR							Siz	ze o	f RA	Мb	locl	k n i	n by	/tes																			

#### 6.2.7 CONFIGID

Bit number	31 30 29	28 27 26	25 24 23	22 21 2	0 19 18	17 16 15	14 13 12 11 10	9876543210
ID (Field ID)	B B B	B B B	B B B	BBB	BB	ВВА	A A A A A	<b>A A A A A A A A A</b>
Value after erase	1 1 1	1 1 1	1 1 1	1 1 1	1 1	1 1 1	1 1 1 1 1	1111111111
ID RW Field Value ID	Value	Descr	iption					
A R HWID		Ident	ification nur	mber for t	he HW.			
B R FWID		Ident	ification nur	mber for t	he firmwa	are that is p	ore-loaded into th	ne chip.

## 6.2.8 DEVICEID[0]

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 0	
ID (Field ID)		Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A	
Value after eras	e	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11	
ID RW Field	Value ID		Va	lue											D	esc	ript	ion	I														
A R						De	evic	e ID	, bit	: 31-	•0, u	iniq	ue l	D fo	or ea	ach	unit	t.															

## 6.2.9 DEVICEID[1]

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Value after eras	e	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
A R						D	evic	e ID	), bit	t 63-	-32,	uni	que	ID f	or e	ach	uni	it.															



## 6.2.10 ER[n] (n=0..3)

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5 4	4	3 2	21	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	4 <i>4</i>	AA
Value after eras	e	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 <sup>.</sup>	11	1
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A R						Er	ncry	ptic	on ro	oot,	woi	rd n.																				

#### 6.2.11 IR[n] (n=0..3)

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Value after eras	e	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
A R						Id	enti	ity r	oot,	wo	rd r	ı.																					

#### 6.2.12 DEVICEADDRTYPE

Bit number	31 30 29 2	3 27 26 25 24 23 22 21 2	0 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)			A
Value after erase	1 1 1 1	1 1 1 1 1 1 1	
ID RW Field Value ID	Value	Description	
A R PUBLIC RANDOM	0	Device address type Public address Random address	

#### 6.2.13 DEVICEADDR[0]

Bit number	31 30 29 28	3 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)	<b>A A A A</b>	
Value after erase	1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
ID RW Field Value ID	Value	Description
A R ADDR		Device address bit 31-0.

## 6.2.14 DEVICEADDR[1]

Bit number	31 30 29 28	3 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		A A A A A A A A A
Value after erase	1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
ID RW Field Value ID	Value	Description
A R ADDR		Device address bit 47-32.



## 6.2.15 OVERRIDDEN

Bit number	31 30 29 28	27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		D A
Value after erase		
ID RW Field	Value ID Value	Description
A R NRF_1MBIT	0 1	Override default values for NRF_1MBIT mode. Use default values for NRF_1MBIT mode.
D R BLE_1MBIT	0 1	Override values for BLE_1MBIT mode is available in FICR.BLE_1MBIT[n] (n=04). No override values for BLE_1MBIT mode available in FICR.

#### 6.2.16 NRF\_1MBIT[n] (n=0..4)

Bit number	3	13	0 29	9 28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5 4	13	2	1	0
ID (Field ID)	Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	4 <i>4</i>	A A	A	A	A
Value after erase	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		• -	-	-	-
ID RW Field	Value ID Va	alu	e		De	scri	ptio	on																							
A R					0	verr	ide	valu	ies f	or N	IRF_	1M	BIT	mo	de.																

### 6.2.17 BLE\_1MBIT[n] (n=0..4)

Bit number	:	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Value after erase		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
ID RW Field	Value ID	Val	ue			De	scri	ptio	on																								
A R						0	verr	ide	valu	ies f	or E	BLE_	1ME	BIT r	noc	le.																	



# 7 User Information Configuration Registers (UICR)

## 7.1 Functional description

The User Information Configuration Registers (UICRs) are NVM registers for configuring user specific settings including code readback protection of the whole code area, or a part of the code area.

The code area can be divided into two regions, code region 0 (CR0) and code region 1 (CR1). Code region 0 starts at address 0x00000000 and stretches into the code area as specified in the CLENR0 register. The area above CLENR0 will then be defined as code region 1. If CLENR0 is not configured, that is, has the value 0xFFFFFFFF, the whole code area will be defined as code region 1 (CR1).

Code running from code region 1 will not be able to write to code region 0. Additionally, the content of code region 0 cannot be read from code running in code region 1 or through the SWD interface if code region 0 is readback protected, see *Section 7.2.2 "RBPCONF"* on page 25.

The main readback protection mechanism that will protect the whole code, that is, both code region 0 and code region 1, is also configured through the UICR, see *Section 7.2.2 "RBPCONF"* on page 25.

The PAGEERASE command in NVMC will only work for code region 1. See *Chapter 5 "Non-Volatile Memory Controller (NVMC)"* on page 14 for more information on how to erase and program the code area and the UICR.

## 7.2 Registers

Register	Offset	Description
REGISTERS		
CLENRO	0x000	Length of code region 0.
RBPCONF	0x004	Read back protection configuration.
XTALFREQ	0x008	Reset value for XTALFREQ in clock, see Chapter 12 "Clock management (CLOCK)" on page 53.
FWID	0x010 0x014 - 0x02C 0x030 - 0x07C 0x080 - 0x0FC	Firmware ID. Reserved. Reserved. Reserved for customer.

 Table 7
 Register overview

#### 7.2.1 CLENR0

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	2 1	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A	۱A	A
Value af	ter eras	e	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11	1	1
ID RW	Field	Value ID	Val	ue			De	scri	ptic	on																							
A RW			[0.	.N]			by N Th Th	rtes (ma lis re lis re	(FIC x va egis egis	cod R. Co Ilue) ter c ter is PPFC	ODE is (F an c s onl	PAC ICR only y us	ESI COI be v sed v	ZE). DEP/ writt	AGE: en i	SIZE f co	E* Fl nter	CR.C nt is	COD 0xF	esiz FFFi	(E-1) FFFF	), bu	t nc	ot lai	rge	r th	nan	25	5.	5			



Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	2 1	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A	A	A
Value af	ter eras	e	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	1	1
ID RW	Field	Value ID	Val	ue			De	scri	ptio	n																							
			0x	FFFF	FFFF	:	Va	lue	afte	r ma	iss e	rase	e of t	flash	n. Int	terp	rete	ed as	i 0.														

## 7.2.2 RBPCONF

Bit	num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 3	32	21	0
ID	Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	В	В	В	В	В	В	В	В	A	A	A /	A /	۹ A	A	A
Va	ue a	fter erase		1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	1	1	1
ID	RW	Field	Value ID	Val	ue			De	escri	pti	on																							
A	RW	PRO		0x				is Di	eadk pres isab nabl	sent led	•				regio	on 0	. Wi	ll be	e igr	nore	d if	pre	pro	grar	nme	ed f	<sup>:</sup> act	ory	/ CC	ode	ž			
В	RW	PALL		0x 0x				D	eadk isab nabl	led	< pro	otec	t all	coc	le in	dev	vice	•																

#### 7.2.3 XTALFREQ

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	32	210
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A A	AAA
Value after erase		1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11	11
ID RW Field	Value ID	Val	ue			De	scri	ptio	on																						
A R	0xFF 0x00					<b>(C</b> 16	5 MH	CK) Iz c	″ OI ryst	n p al is		e 5. ed.	Q in 3.	CLC	DCK	, see	Ch	apt	ter i	12 "	'Clo	ck I	mai	na	gei	me	ent				

## 7.2.4 FWID

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	2 1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	A	A	A	A	A	A	A	A /	۱A	A
Value after erase	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	1	1
ID RW Field Value ID	) Val	lue			De	scri	iptio	on																							
A R												r the FICF						into	the	e chi	ip. T	his l	ID i	s u	sec	w k	hei	n			



# 8 Memory Protection Unit (MPU)

The Memory Protection Unit (MPU) can protect the entire memory against readback and also protect parts of the memory area from accidental access by the CPU.

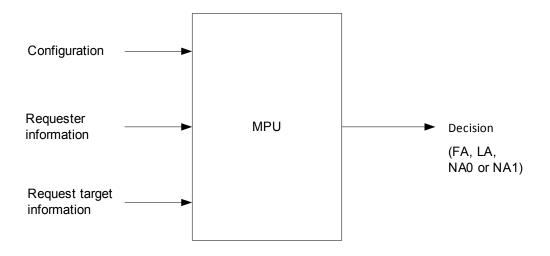


Figure 3 MPU block diagram

## 8.1 Functional description

Protect all (PALL) is configured by writing '0' to UICR.RBPCONF.PALL. When enabled, the debugger (SWD) will no longer have access to code region 0 or code region 1, as well as no longer being able to access RAM or any peripherals except for the NVMC. The debugger will always have access to the NVMC peripheral independent of protection settings.

Code memory, RAM, and peripherals can be divided into two regions: region 0 and region 1. Code memory regions are configured in the CLENRO register in the User Information Configuration Register (UICR), see the Memory isolation and run-time protection section in the *Appendix A: on page 186*. When memory protection is enabled, these regions will be used by the Memory Protection Unit to enforce runtime protection and readback protection of resources classified as region 0.

Independent of protection settings, code region R0 (CR0) will always have full access to the system. The NVMC.ERASEPCR0 register, which is used to erase contents from code region 0, can only be accessed from a program in code region 0.

Only the CPU can do fetches from code memory, and these will always be granted.

Except when generated by the SWD interface, accesses that are not granted by the MPU will result in a hard-fault.

Readback protection of code region 0 is enabled by writing '0' to UICR.RBPCONF.PR0. When enabled, only code running from code region 0 will be able to access the code in code region 0. Accesses generated by code running from code region 1 or from RAM, as well as accesses generated by the debugger (SWD), will not be granted when code region 0 is protected.

The main role for the two region memory protection system is to allow run time protection for SoftDevices installed running on the IC. Please refer to *Appendix A: on page 186* for a description of the nRF51 software architecture and use of the memory protection system with SoftDevices.



#### 8.1.1 Inputs

The MPU has three classes of inputs. These are:

- Configuration
  - Readback protection configuration from UICR and FICR.
- Information about requester
  - Source of memory access request (SWD or CPU program).
  - If the request source is a CPU program; region from which the program is running (region 0 or region 1).
  - Type of access request (read or write).
- Target information
  - Memory category requested access to (code, RAM, or PER).
  - Memory region requested access to (region 0 or region 1).

#### 8.1.2 Output

The MPU outputs the level of memory access that shall be given to a memory access request. The access levels the MPU can give are as follows:

- Full Access (FA)
  - Full read write access to the requested memory.
- Limited Access (LA)
  - Full read access
  - No write access. Write will generate hard fault exception.
- No Access 0 (NA0)
  - No read or write access
  - Read will return 0
  - Write will have no effect
- No Access 1 (NA1)
  - No read or write access
  - Read or write will generate hard fault exception

#### 8.1.3 Output decision table

The output MPU access level based on the MPU inputs is given in the table below.

The given access level is dependent on settings in the Information Configuration Registers (ICRs). See the UICR and FICR chapters for more details.



				R	equest t	arget		
Request source	UICR.RBPCONF.PALL (Readback protect entire code memory)	UICR.RBPCONF.PR0 or FICR.PPFC* (Readback protect code region 0)	Code R0	Code R1	RAM RO	RAM R1	PER R0	PER R1
SWD	0xFF	0xFF	FA	FA	FA	FA	FA	FA
	0xFF	0x00	NA0	FA	FA	FA	FA	FA
	0x00	Х	NA0	NA0	NA0	NA0	NA0	NA0
Code R0	Х	Х	FA	FA	FA	FA	FA	FA
Code R1	Х	0xFF	LA	FA	LA	FA	LA	FA
	Х	0x00	NA1	FA	LA	FA	LA	FA
RAM	0xFF	0xFF	FA	FA	FA	FA	FA	FA
R0 / R1	0xFF	0x00	NA1	FA	FA	FA	FA	FA
	0x00	Х	NA1	NA1	FA	FA	FA	FA

X: Don't care

LA: Limited Access

NA0: No Access 0

NA1: No Access 1

FA: Full Access

Table 8 MPU output decision table based on the MPU inputs and the ICR configuration

#### 8.1.4 Exceptions from table

There are some exceptions from *Table 8*. These exceptions are:

- The NVMC.ERASEALL and NVMC.ERASEUICR registers have conditional write access depending on the readback protection settings in the Information Configuration registers. These exceptions are described in the NVMC chapter, see *Chapter 5 "Non-Volatile Memory Controller (NVMC)"* on page 14.
- The NVMC.ERASEPCR0 register can only be accessed from a program in code region 0.
- The UICR.CLENR0 and the FICR. CLENR0 registers can only be modified when the register value equals the default value (0xFF). This is to avoid that the memory region limits are modified to bypass readback protection.

#### 8.1.5 NVM protection blocks

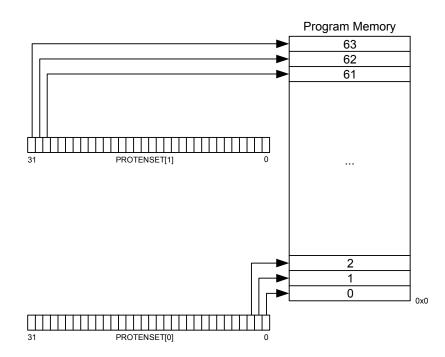
The protection mechanism for NVM can be used to prevent erroneous application code from erasing or writing to protected blocks. Non-volatile memory can be protected from erases/writes depending on settings in the PROTENSET registers. One bit in a PROTENSET register represents one protected block. There are two PROTENSET registers of 32 bits which means there are 64 protectable blocks in total. See the device specific product specifications for more information on how the NVM protection blocks are organized.



**Note:** If an erase or write to a protected block is detected, the CPU will hard fault. If an ERASEALL operation is attempted from the CPU while any block is protected it will be blocked and the CPU will hard fault.

On reset, all the protection bits are cleared. To ensure safe operation, the first task after reset must be to set the protection bits. The only way of clearing protection bits is by resetting the device from any reset source.

The protection mechanism is turned off when in debug mode (a debugger is connected) and the DISABLEINDEBUG register is set to disable.



*Figure 4 Protected regions of program memory* 

#### 8.2 Registers

Register	Offset	Description
PERRO	0x528	Definition of peripherals in memory region 0.
RLENRO	0x52C	Length of RAM region 0.
<b>PROTENSET0</b>	0x600	Protection bit enable set register
PROTENSET1	0x604	Protection bit enable set register
DISABLEINDEBUG	0x608	Disable protection mechanism in debug mode

Table 9 Register overview



#### 8.2.1 PERR0

Bit	num	ber		31	30	29	28	27	26	25	5 24	2	3 22	21	20	19	18	3 17	16	5 1	5 1·	41	3	12	11	10	9	8	7	65	4	3 2	210
ID (	Field	IID)		z	Y	-	-	-	-	-	-	-	-	-	-	т	S	R	Q	P	, c		N	м	L	K	J	I	н	G -	E	DO	В
Res	et va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	) C	)	0	0	0	0	0	0	0	0 0	0	0 0	00
ID I	RW	Field	Value ID	Va	lue			De	scri	pti	ion																						
A	RW	POWER_CLOCK	REGION0 REGION1	0 1				pe Re	assi eripł egio egio	ner n 0	ral. )	/ER	and	CLC	CK	and	all	l oth	er p	oeri	phe	eral	s w	vith	ID=	0, i	as i	egi	ion	0 0	r reg	gioi	n 1
В	RW	RADIO	REGION0 REGION1	0 1				Re	assi egio egio	n 0	)	10	as re	gion	0 c	or re	gio	on 1 p	per	iph	eral	l.											
С	RW	UARTO	REGION0 REGION1	0 1				Re	assi egio egio	n 0	)	то	as re	gion	0 c	or re	gio	on 1 j	per	iph	era	I.											
D	RW	SPI0_TWI0	REGION0 REGION1	0 1				Cl	assi	fy S	SPI0	an	d TW	′10 a	s re	gior	۱O	or re	egio	on 1	pe	rip	her	al.									
E	RW	SPI1_TWI1	REGION0 REGION1	0 1				Re	assi egio egio	n 0	)	, SF	PIS1, a	and	тw	l1 as	s re	egior	0 ר	or r	egi	on	1 p	erip	hei	al.							
G	RW	GPIOTE	REGION0 REGION1	0 1				Cl	assi	fy (	GPIC	DTE	as re	gio	n 0 i	or re	egi	on 1	pe	ripł	hera	al.											
Η	RW	ADC	REGION0 REGION1	0 1				Re	assi egio egio	n 0	)	as	regio	on 0	or r	egio	on	1 pe	rip	her	al.												
I	RW	TIMERO	REGION0 REGION1	0 1				Re	assi egio egio	n 0	)	ERO	as re	gio	n 0 (	or re	egi	on 1	pe	ripł	nera	al.											
J	RW	TIMER1	REGION0 REGION1	0 1				Re	assi egio egio	n 0	)	ER1	as re	gio	ר 0 י	or re	egi	on 1	pe	ripł	nera	al.											
К	RW	TIMER2	REGION0 REGION1	0 1				Cl	assi	fy ٦	TIME	R2	as re	gio	ר 0 י	or re	egi	on 1	pe	ripł	hera	al.											
L	RW	RTCO	REGION0 REGION1	0 1				Cl	assi	fy F	RTC	) as	s regi	on (	or	regi	on	1 pe	erip	he	ral.												
Μ	RW	TEMP	REGION0 REGION1					Cl	assi	fy 1	TEM	Ра	s reg	ion	) or	reg	ior	n1p	eri	ohe	ral.												
N	RW	RNG	REGION0 REGION1	0 1				Cl	assi	fy F	RNG	as	regio	on 0	or r	egio	on	1 pe	rip	her	al.												
0	RW	ECB	REGION0 REGION1	0 1				Cl	assi	fy E	ECB	as	regio	n 0 (	or re	egio	n 1	l per	riph	nera	ıl.												
Ρ	RW	CCM_AAR	REGION0 REGION1	0 1				Cl	assi	fy (	ССМ	l ar	nd AA	R as	reg	gion	0 0	or re	gio	n 1	pei	ripł	nera	al.									
Q	RW	WDT	REGION0 REGION1					Cl	assi	fy ۱	WDT	as	regi	on 0	or	regi	on	1 pe	erip	her	al.												



Bi	num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	87	6	5	4 3	2	10
ID	(Field	IID)		z	Y	-	-	-	-	-	-	-	-	-	-	т	S	R	Q	Ρ	ο	Ν	м	L	K	J	IH	G	-	EC	) C	ΒA
Re	set va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0 0	0	0 0
ID	RW	Field	Value ID	Va	lue			Des	scri	ptic	on																					
R	RW	RTC1	REGION0 REGION1	0 1				Cla	assif	fy R⁻	TC1	as r	egi	on 0	or	regi	on ′	1 pe	riph	nera	ıl.											
S	RW	QDEC	REGION0 REGION1	0 1				Cla	assif	fy Q	DEC	Cas	reg	ion	0 or	reg	ion	1 p	erip	her	al.											
Т	RW	COMP_LPCOMP	REGION0 REGION1	0 1				Cla	assif	fy C	OM	P_L	PCC	MP	as r	egio	on C	) or	regi	on	1 pe	eripł	nera	I.								
Y	RW	NVMC	REGION0 REGION1	0 1				Cla	assif	fy N	VM	C as	reg	ion	0 0	r reg	gion	1 p	erip	her	ral.											
Z	RW	PPI	REGION0 REGION1	0 1				Cla	assif	fy Pl	PI as	s reg	gion	00	r re	gior	n 1 p	beri	phei	ral.												

#### 8.2.2 RLENRO

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	87	6	5	4	3	2	10
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A	A	A	A	A	A	AA
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0 0
ID RW Field Value ID																rip	tion	I												
ID     RW     Field     Value ID     Value       A     RW     This register specifies the size of RAM region 0.       Given a base address for the RAM called RAMBA, RAM addresses < RAMBA +																														
					RL	EN	R0 a	re c	lass	ifie	d as	r the reg 1 RA	ion														20			
					Tł	ne a	ddr	ess	(RA	MBA	4 +	rlei	NR0	) ha	s to	o be	wo	rd-a	lign	ied.										
						AMB ie cł						vaila	ble	RAN	A is	def	ineo	d in	the	pro	duc	t sp	ecif	ficat	ion	of				



## 8.2.3 PROTENSET0

Bit	numb	per		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	76	5	4	3	2	10
ID (	Field	ID)		FF	EE	DD	сс	BB	AA	z	Y	Х	w	v	U	т	S	R	Q	Ρ	0	Ν	М	L	к	J	H	ΗG	F	E	D	C	ΒA
Res	et va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0 0
ID	RW	Field	Value ID		Va	lue											De	escr	ripti	ion													
A	W R	PROTREGO	ENABLE		ROTE	EN.0		Ena	tect ables adba	s pr	otec	tio	n of	blo	ck 0																		
В	W R	PROTREG1	ENABLE		Ena	tect ables adba	s pro	otec	tio	n of	blo	ck 1	•																				
FF	W R	PROTREG31	ENABLE		ROTE	EN.3	1	Ena	tect ables adba	s pr	otec	tio	n of	blo	ck 3	1.																	



## 8.2.4 PROTENSET1

Bit	numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	87	6	5	4	3	2	10
ID (	Field	ID)		FF	EE	DD	сс	BB	AA	z	Y	Х	w	v	U	т	S	R	Q	Ρ	0	Ν	м	L	к	J	I F	I G	F	E	D	C	BA
Res	et va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0 0
ID	RW	Field	Value ID		Va	lue											De	esc	ript	ion													
A	W R	PROTREG32	ENABLE		Rote	EN.32	2	Ena	otect ables adba	s pro	otec	tio	n of	blo	ck 3	2.																	
В	W R	PROTREG33	3	Ena	otect ables adba	s pro	otec	tio	n of	blo	ck 3	3.																					
FF	W R	PROTREG63	ENABLE	1 PF	Rote	EN.6	3	Ena	otect ables adba	s pro	otec	tio	n of	blo	ck 6	53.																	

#### 8.2.5 DISABLEINDEBUG

Bit number	31	30 29	28	27	26	25	24	23	22	21 2	20	19	18	17	16	15	14	13	12	11	10	9 8	B 7	6	5 4	43	3 2	2 1	0
ID (Field ID)	-		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	- 1		• -	-	A
Reset value	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0 (	0 0	0	0	1
ID RW Field Value ID Value Description																													
A RW DISABLEINDEBUG DISABLE ENABLE	•			reg ma Dis	sable giste ode. sable able	r wil e in c	I on debi	ly d ug.																					



# 9 Peripheral interface

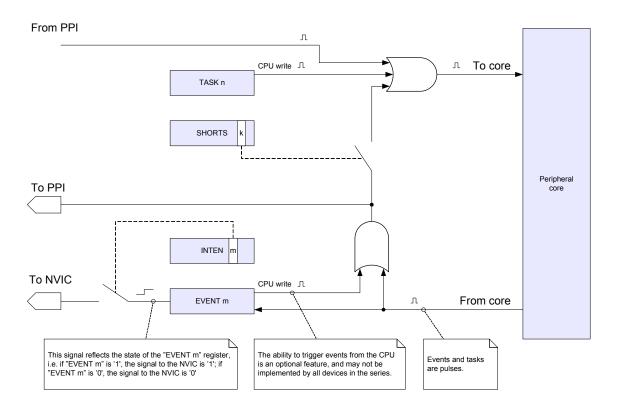


Figure 5 Tasks, events, shortcuts, and interrupts

## 9.1 Functional description

All peripherals on nRF51 series devices can be accessed through the standard ARM<sup>®</sup> Cortex Advanced Peripheral Bus (APB) and AMBA High-performance Bus (AHB) registers as well as through task, event, and interrupt registers.

#### 9.1.1 Peripheral ID

For peripherals on the APB bus there is a direct relationship between its ID and its base address.

Every peripheral is assigned a fixed block of 0x1000 bytes, that is, a total of 1024 registers of 4 bytes on the APB bus. The peripheral with base address 0x40000000 is therefore assigned ID=0, and a peripheral with the base address 0x40001000 is assigned ID=1. The peripheral with the base address 0x4001F000 is assigned ID=31.



Peripherals may share the same ID, which may impose one or more of the following limitations:

- Peripherals do not share any registers or common resources, but the total number of registers available for each peripheral is reduced compared to a peripheral that has a dedicated ID.
- Peripherals share some registers or other common resources.
- Only one of the peripherals can be used at a time.
- Both peripherals are optional in the series, and only one of them is instantiated in any given chip.

#### 9.1.2 Bit set and clear

Registers with multiple single-bit bit-fields may implement the "set and clear" pattern. This pattern enables firmware to set and clear individual bits in a register without having to perform a read-modify-write operation on the main register. This pattern is implemented using three consecutive addresses in the register map where the main register is followed by a dedicated SET and CLR register in that order. The SET register is used to set individual bits in the main register while the CLR register is used to clear individual bits in the SET or CLR register will set or clear the same bit in the main register respectively. Reading the SET or CLR registers returns the value of the main register.

#### 9.1.3 Tasks

Tasks are used to trigger actions in a peripheral, for example, to start a particular behavior. A peripheral can implement multiple tasks with each task having a separate register in that peripheral's task register group.

A task is triggered when firmware writes a '1' to the task register or when the peripheral itself, or another peripheral, toggles the corresponding task signal, see *Figure 5* on page 34. All tasks follow the register layout in *Table 10* on page 37.

#### 9.1.4 Events

Events are used to notify peripherals and the CPU about events that have happened, for example, a state change in a peripheral. A peripheral may generate multiple events with each event having a separate register in that peripheral's event register group.

An event is generated when the peripheral itself toggles the corresponding event signal, whereupon the event register is updated to reflect that the event has been generated, see *Figure 5* on page 34. An event register is only cleared when firmware writes a '0' to it.

Events can be generated by the peripheral even when the event register is set to '1'. All events follow the register layout described in *Table 10* on page 37.

#### 9.1.5 Shortcuts

A shortcut is a direct connection between an event and a task within the same peripheral. If a shortcut is enabled, its associated task is automatically triggered when its associated event is generated.

Using a shortcut is the equivalent to making the same connection outside the peripheral and through the PPI. However, the propagation delay through the shortcut is usually shorter than the propagation delay through the PPI.

Shortcuts are predefined, which means their connections cannot be configured by firmware. Each shortcut can be individually enabled or disabled through the shortcut register, one bit per shortcut, giving a



maximum of 32 shortcuts for each peripheral. All shortcut registers follow the register layout described in *Table 10* on page 37, (SHORTS).

#### 9.1.6 Interrupts

An interrupt is an exception that is generated by an event and can interrupt the program flow of the CPU. All peripherals on the APB bus support interrupts. A peripheral only occupies one interrupt, and the interrupt number follows the peripheral ID, for example, the peripheral with ID=4 is connected to interrupt number 4 in the Nested Vector Interrupt Controller (NVIC).

Using the INTEN register, you can configure every event in a peripheral to generate that peripheral's interrupt. You can enable multiple events to generate interrupts simultaneously. To resolve the correct interrupt's source, firmware can query the event registers found in the event group in the peripherals register map.

Each event implemented in the peripheral is associated with a specific bit position in the INTEN register. The correct bit position can be derived from the event's address. The event on address 0x100 is associated with bit 0 in the INTEN register, the event at address 0x104 is associated with bit 1, and so on. The event at address 0x17C is identified with bit 31 in the INTEN register. This pattern effectively limits the maximum number of events in a peripheral to 32.

The INTEN register implements the "set and clear" pattern, which is illustrated in *Table 10* on page 37, that is, INTEN, INTENSET, and INTENCLR. The relationship between tasks, events, shortcuts, and interrupts is shown in *Figure 5* on page 34.



## 9.2 Register overview tables

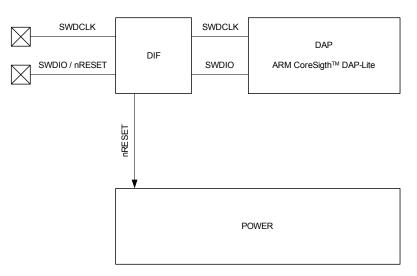
All peripherals follow the register group pattern in *Table 10*; tasks are grouped together, events are grouped together, and all other register types are grouped together. In addition, SHORTS and INTEN registers have a fixed location in the register map.

Register	Offset	Description
TASKS		
{TASK0}	0x000	Description of the first task
{TASK1}	0x004	Description of the second task
<>		
{TASK31}	0x07C	Description of the 32nd task (last task)
EVENTS		
{EVENT0}	0x100	Description of the first event
{EVENT1}	0x104	Description of the second event
<>		
{EVENT31}	0x17C	Description of the 32nd event (last event)
REGISTERS		
SHORTS	0x200	Shortcut register
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
{REG0}	0x400	First generic register
<>	<>	$\Leftrightarrow$
{REGN}	0x7FC	Last generic register

Table 10 Example of register overview table



# 10 Debugger Interface (DIF)



*Figure 6 Debugger interface* 

## 10.1 Functional description

nRF51 devices support the Serial wire Debug (SWD) interface from ARM. The interface has two lines; SWDCLK and SWDIO. SWDIO and nRESET share the same physical pin. The Debugger Interface (DIF) module is responsible for handling the resource sharing between SWD traffic and reset functionality. The **SWDCLK** pin has an internal pull down resistor and the **SWDIO/nRESET** pin has an internal pull up resistor.

#### 10.1.1 Normal mode

The DIF module will be in normal mode after power on reset. In this mode the SWDIO/nRESET pin acts as a normal active low reset pin.

To guarantee that the device remains in normal mode, the SWDCLK line must be held low, that is, '0', at all times. Failing to do so may result in the DIF entering into an unknown state and may lead to undesirable behavior and power consumption.

#### 10.1.2 Debug interface mode

Debug interface mode is initiated by clocking one clock cycle on SWDCLK with SWDIO=1. Due to delays caused by starting up the DAP's power domain, a minimum of 150 clock cycles must be clocked at a speed of minimum 125 kHz on SWDCLK with SWDIO=1 to guaranty that the DAP is able to capture a minimum of 50 clock cycles.

If the device is in System OFF mode, see the POWER chapter, *Chapter 11 "Power management (POWER)"* on page 40, for more information about System OFF mode, entering into debug interface mode will generate a wake-up.

In debug interface mode, the SWDIO/nRESET pin will be used as SWDIO. The pin reset mechanism will therefore be disabled as long as the device is in debug interface mode.

In debug interface mode, System OFF will be emulated to facilitate debugging of the device while in System OFF. Power numbers will naturally be higher in emulated System OFF compared to normal System OFF. See emulated System OFF in *Chapter 11 "Power management (POWER)"* on page 40 for more information.



## 10.1.3 Resuming normal mode

Normal mode can always be resumed by performing a "hard-reset" through the SWD interface:

- 1. Enter debug interface mode.
- 2. Enable reset through the RESET register in the POWER peripheral.
- 3. Hold the SWDCLK and SWDIO/nRESET line low for a minimum of 100 μs.

You can also generate a "hard-reset" by performing a power on reset, or a brown-out reset.



# 11 Power management (POWER)

The power management on the nRF51 series gives you unique flexibility through the orthogonal power control of all system blocks on the devices.

## 11.1 Functional description

#### 11.1.1 Power supply

The nRF51 supports three different power supply alternatives: internal DC/DC converter setup, internal LDO setup, and Low Voltage mode setup.

#### 11.1.1.1 DC/DC converter setup

Selected nRF51 series devices have a Buck type DC/DC converter that steps down the supply voltage VDD. The resulting voltage is then used by an internal LDO that supplies the system with power, see *Figure 7*.

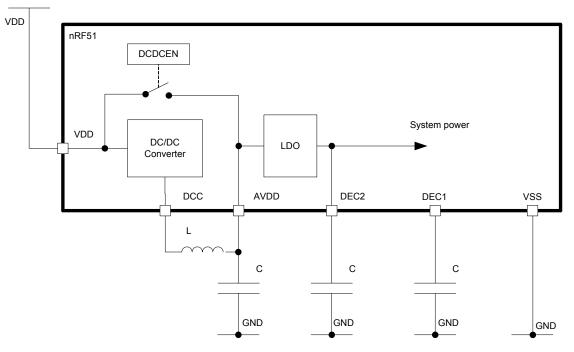


Figure 7 DC/DC converter

The DC/DC converter requires an external LC filter and is enabled through the DCDCEN register as illustrated in *Figure 7*. See the product specification for more information about component values.

The DC/DC converter only reduces the system's net power consumption if VDD is above the minimum voltage (stated in the product specification) and the internal current consumption ( $I_{DD}$ ) gives a DC/DC conversion factor of  $F_{DCDC} < 1$ . Therefore, to save power it is recommended to disable the DC/DC converter if VDD is below the minimum voltage or  $F_{DCDC} > 1$ .

AVDD is connected to VDD internally if the DC/DC converter is not enabled. This internal connection introduces a small series resistance between VDD and AVDD, see the product specification for more information.



#### Calculating current when the DC/DC converter is enabled

The device current consumption when the DC/DC converter is enabled (I<sub>DD,DCDC</sub>) can be calculated using *Equation 1*, the parameters in *Table 11* and the conversion factor chart in *Figure 7* on page 40.

Parameter	Description	Value
I <sub>DD</sub>	Internal current consumption (current drawn from device power regulators) under Normal Test Conditions (NTC)	Calculated by adding current values from Electrical Specification tables in the device product specification.
I <sub>DD,DCDC</sub>	Current drawn from the external power supply (VDD) when the DC/DC converter is enabled	Calculated using <i>Equation 1</i> .
F <sub>DCDC</sub>	DC/DC current conversion factor based on DC/DC converter efficiency	Interpolated from <i>Figure 7</i> .
VDD	Voltage at VDD pin	
T <sub>START,DCDC</sub>	DC/DC converter startup time.	Calculated using <i>Equation 2</i> .

 Table 11 DC/DC current calculation parameters

$$I_{DD, DCDC} = F_{DCDC} \cdot I_{DD}$$

#### **Equation 1** DC/DC current calculation

The internal current consumption ( $_{IDD}$ ), calculated using electrical specification data from the product specification, is used with the supply voltage (VDD) to find the current factor ( $F_{DCDC}$ ) using *Figure 8*.

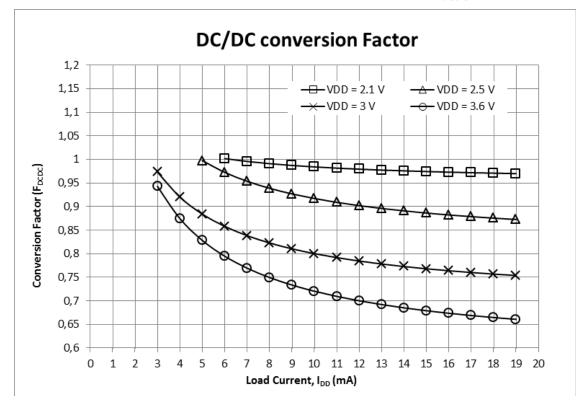


Figure 8 DC/DC conversion factor



If  $F_{DCDC} < 1$ , then  $I_{DD,DCDC} < I_{DD}$  when the DC/DC converter is enabled, resulting in a decrease in power consumption.

If  $F_{DCDC} > 1$ , then  $I_{DD} < I_{DD,DCDC}$  when the DC/DC converter is enabled, resulting in an increase in power consumption. This is due to the base run current of the DC/DC converter ( $I_{DCDC}$ ) being the dominant factor.

#### Continuous use of the DC/DC converter

Using the DC/DC converter continuously can save power when average load current is expected to be larger than 4 to 6 mA. When average load current is less than 4 mA, the conversion factor  $F_{DCDC}$  approaches a value > 1 and continuous use of the DC/DC converter will increase current consumption.

For example, if a battery voltage VDD = 3V and an average internal current consumption, including the DC/DC converter run current ( $I_{DCDC}$ ), is  $I_{DD}$  = 10 mA:  $F_{DCDC}$  would be 0.8 using *Figure 7* on page 40. The current drawn from VDD, when the DC/DC converter is enabled, would be  $I_{DD,DCDC}$  = 0.8 x 10 mA = 8 mA.



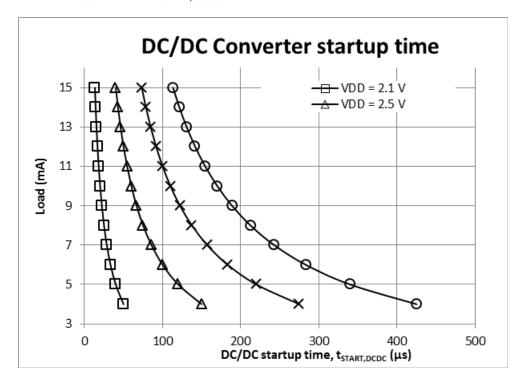
#### Non-continuous use of the DC/DC converter

The DC/DC converter has a startup time of  $t_{START,DCDC}$ . This is the time it takes for the DC/DC converter to pull down AVDD to the internal supply voltage of 1.9 V. This is a function of the decoupling capacitance on the AVDD pin (C<sub>AVDD</sub>), supply voltage (V<sub>DD</sub>), and load current (I<sub>DD</sub>):

$$t_{START, DCDC} = \frac{(VDD - 1.9)(C_{AVDD})}{I_{DD}}$$

#### **Equation 2** DC/DC converter startup time

For example, given a decoupling capacitance ( $C_{AVDD}$ ) of 1  $\mu$ F, supply voltage (VDD) of 3.6 V, and internal current consumption ( $I_{DD}$ ) of 4 mA, t<sub>START,DCDC</sub> is 425  $\mu$ s.



**Figure 9** DC/DC converter startup time when  $C_{AVDD} = 1 \, \mu f$ 

t<sub>START,DCDC</sub> must be considered when enabling the DC/DC converter during periods of peak current consumption because the DC/DC converter will draw base current of I<sub>DCDC</sub> when enabled, including the startup period.



For example, if the DC/DC converter was enabled and VDD = 3V and peak current  $I_{DD}$  = 10 mA,  $F_{DCDC}$  would be 0.8 using *Figure 8* on page 41 and  $t_{START,DCDC}$  would be 110 µs using *Equation 2* on page 43. The current drawn from the battery in the first 110 µs would be  $I_{DD}$  = 10 mA + ( $I_{DCDC}$ ) = 10.3mA. The current drawn from the battery after the first 110 µs would be  $I_{DD,DCDC}$  = 8.24 mA. Peak current, in this case, would have to be drawn for more than 198 µs to save power:

 $\frac{8.24mA(110\mu s)}{10.3mA} + 110\mu s = 198\mu s$ 

- **Note:** If the DC/DC converter is to be used when the Radio is enabled, it must start (AVDD = 1.9V) before the TXENABLE or RXENABLE tasks can be set. The DC/DC converter cannot be enabled within t<sub>START,DCDC</sub> of the Radio starting. The software managing the non-continuous use of the DC/DC converter must ensure this is true.
  - The DC/DC converter does not operate over the whole device voltage range and must be disabled (switching the device to use the internal LDO) when the voltage drops to the lower threshold of the supply voltage range.

#### 11.1.1.2 Internal LDO setup

The internal DC/DC converter can be bypassed if it is not going to be used. When the DC/DC converter is bypassed, only the internal LDO is active as illustrated in *Figure 10*. The internal LDO will then generate the system power directly from the supply voltage VDD. It is recommended that the DC/DC converter is disabled in this setup.

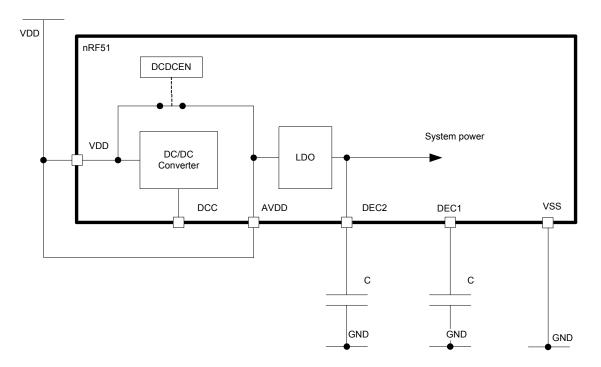
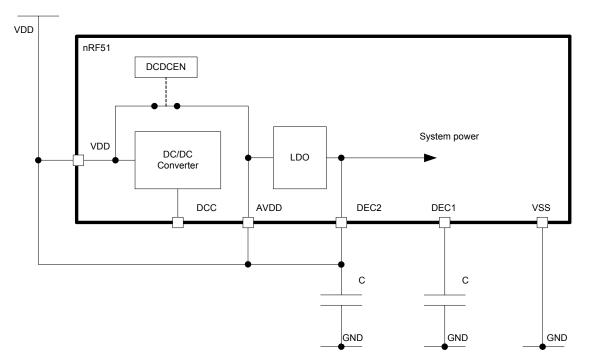


Figure 10 LDO regulator only



#### 11.1.1.3 Low voltage mode setup

If you have a stable, low voltage available for the nRF51 device, it is possible to configure the device in low voltage mode as illustrated in *Figure 11*. In this mode the internal LDO is bypassed and the system is powered directly from the supply voltage VDD. See the product specification for more information about which voltage levels are supported in low voltage mode. In low voltage mode, the DC/DC converter must be disabled. Additional requirements may apply to the accuracy and stability of the supply voltage in low voltage mode. See the product specification.



*Figure 11* Low voltage mode

#### 11.1.2 System OFF mode

System OFF is the deepest power saving mode the system can enter. In this mode, the system's core functionality is powered down and all ongoing tasks are terminated. The only mechanism that is functional and responsive in this mode is the reset and the wake-up mechanism.

One or more blocks of RAM can be retained in System OFF mode depending on the settings in the RAMON register.

The system can be woken up from System OFF mode either from the DETECT signal generated by the GPIO peripheral, by the ANADETECT signal generated by the LPCOMP module, or from a reset. When the system wakes up from OFF mode, a system reset is performed.

Before entering System OFF mode you must make sure that all on-going EasyDMA transactions have been completed. This is usually accomplished by making sure that the EasyDMA enabled peripheral is not active when entering System OFF.

#### 11.1.2.1 Emulated System OFF mode

If the device is in debug interface mode, System OFF will be emulated to secure that all required resources needed for debugging are available during System OFF, see *Chapter 10 "Debugger Interface (DIF)"* on



page 38, for more information. This includes the following key components: DAP, DIF, CLOCK, POWER, NVMC, MPU, CPU, CODE, and RAM. Since the CPU is kept on in emulated System OFF mode, it is recommended to add an infinite loop directly after entering System OFF, to prevent the CPU from executing code that normally should not be executed.

### 11.1.3 System ON mode

System ON mode is a fully operational mode, where the CPU and selected peripherals can be brought into a state where they are functional and more or less responsive depending on the sub power mode selected.

In System ON mode the CPU can either be active or sleeping. The CPU enters sleep by executing the WFI or WFE instruction found in the CPU's instruction set. In WFI sleep the CPU will wake up as a result of an interrupt request if the associated interrupt is enabled in the NVIC. In WFE sleep the CPU will wake up as a result of an result of an interrupt request regardless of the associated interrupt being enabled in the NVIC or not.

The system implements mechanisms to automatically switch on and off the appropriate power sources depending on how many peripherals are active, and how much power is needed at any given time. The power requirement of a peripheral is directly related to its activity level. The activity level is usually raised and lowered when specific tasks are triggered or events generated, see individual chapters describing the different peripherals for more information on how to optimize power consumption in System ON mode.

#### 11.1.3.1 Sub power modes

During CPU sleep, in System ON mode, the system can reside in one of the following two sub power modes:

- Constant Latency
- Low Power

In Constant Latency mode (for more information, see the device specific product specification) the CPU wakeup latency and the PPI task response will be constant and kept at a minimum. This is secured by forcing a set of base resources while in sleep, see the device specific product specification for more information about which resources are forced on. The advantage of having a constant and predictable latency will be at the cost of having increased power consumption. The Constant Latency mode is selected by triggering the CONSTLAT task.

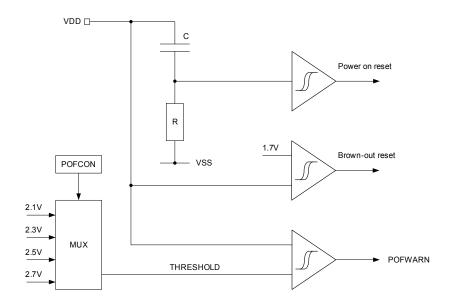
In Low Power mode the automatic power management system, described in *Section 11.1.3 "System ON mode"*, will be most effective and save most power. The advantage of having low power will be at the cost of having varying CPU wakeup latency and PPI task response. The Low Power mode is selected by triggering the LOWPWR task.

When the system enters ON mode, it will, by default, reside in the Low Power sub power mode.



### 11.1.4 Power supply supervisor

The power supply supervisor initializes the system at power-on and provides an early warning of impending power failure. In addition the power supply supervisor puts the system in a reset state if the supply voltage is too low for safe operation (brown-out). The power supply supervisor is illustrated in *Figure 12*.

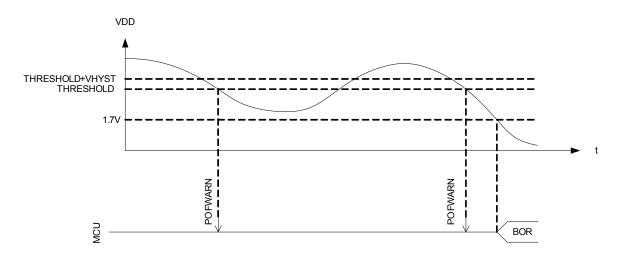


*Figure 12 Power supply supervisor* 

#### 11.1.4.1 Power-fail comparator

The power failure comparator provides the CPU with an early warning of impending power failure. It will not reset the system, but give the CPU time to prepare for an orderly power-down. It also provides hardware protection of data stored in program memory by preventing write and erase instructions from being executed.

The comparator has approximately 0.1 V of hysteresis (VHYST), as illustrated in *Figure 13*.



*Figure 13 Power failure comparator (BOR = Brown-out reset)* 



The power failure comparator is only intended to be used as long as the CPU is running. If the power failure comparator is enabled when the CPU is not running, this may result in unpredictable behavior in the power failure comparator.

### 11.1.5 RAM blocks

Each of the available RAM blocks, which each may contain multiple RAM sections, can power up and down independently in both System ON and System OFF mode. See *Chapter 4 "Memory"* on page 11 for more information about RAM blocks and their sections.

#### 11.1.6 Reset

The nRF51 series implements various reset sources. After a reset the CPU can query the RESETREAS (reset reason register) to find out which source generated the reset.

#### 11.1.6.1 Power-on reset

The power-on reset generator initializes the system at power-on. The system is held in reset state until the supply has reached the minimum operating voltage, see the device specific product specification for more information.

#### 11.1.6.2 Pin reset

A pin reset is generated when the physical reset pin on the device is asserted. Since the debugger interface uses the same pin as the pin reset mechanism, a pin reset will not be available when the device is in debug interface mode unless explicitly enabled in the RESET register.

#### 11.1.6.3 Wakeup from OFF mode reset

The device is reset when it wakes up from OFF mode.

The DAP is not reset following a wake up from OFF mode if the device is in debug interface mode, see *Chapter 10 "Debugger Interface (DIF)"* on page 38 for more information.

#### 11.1.6.4 Soft reset

A soft reset is generated when the SYSRESETREQ bit of the Application Interrupt and Reset Control Register (AIRCR register) in the ARM<sup>®</sup> core is set.

#### 11.1.6.5 Watchdog reset

A Watchdog reset is generated when the watchdog times out.

#### 11.1.6.6 Brown-out reset

The brown-out reset generator puts the system in reset state if the supply voltage drops below the brownout reset threshold.

#### 11.1.6.7 Retained registers

A retained register is a register that will retain its value in System OFF mode, and through a reset depending on reset source. See individual peripheral chapters for information of which registers are retained for the different peripherals.



#### 11.1.6.8 Reset behavior

				Rese	t target			
Reset source	CPU	Peripherals	GPIO	DAP	RAM <sup>1</sup>	WDT	Retained registers	RESETREAS
CPU lockup <sup>2</sup>	√	$\checkmark$	$\checkmark$					
Soft reset	$\checkmark$	$\checkmark$	$\checkmark$					
Wakeup from System OFF mode reset	√	✓		√3	√4			
Watchdog reset <sup>5</sup>	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	
Pin reset <sup>6</sup>	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	
Brownout reset	✓	✓	~	$\checkmark$	✓	✓	✓	$\checkmark$
Power on reset	$\checkmark$	✓	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$	✓

1. The RAM is never reset, but depending on reset source, RAM content may be corrupted.

2. Reset from CPU lockup is disabled if the device is in debug interface mode. CPU lockup is not possible in System OFF.

3. The DAP will not be reset if the device is in debug interface mode.

4. RAM is not reset on wake-up from OFF mode, but depending on settings in the RAMON register parts, or the whole RAM, may not be retained after the device has entered System OFF mode.

5. Watchdog reset is not available in System OFF.

6. Not available when device is in debug interface mode.

## 11.2 Registers

Register	Offset	Description
TASKS		
CONSTLAT	0x078	Enable constant latency mode
LOWPWR	0x07C	Enable low power mode (variable latency)
EVENTS		
POFWARN	0x108	Power failure warning
REGISTERS		
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
RESETREAS	0x400	Reset reason
SYSTEMOFF	0x500	System OFF register
POFCON	0x510	Power failure configuration
GPREGRET	0x51C	General purpose retention register
RAMON	0x524	RAM on/off
RESET	0x544	Configure reset functionality
DCDCEN	0x578	DCDC enable register

Table 12 Register overview



### 11.2.1 RESETREAS

Unless cleared, this register is cumulative. A field is cleared by writing '1' to it. If none of the reset sources are flagged, it indicates that the chip was reset from the on-chip reset generator.

Bit r	umbe	r		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	87	6	5	4	3	2 <sup>.</sup>	10
ID (F	ield ID	))		-	-	-	-	-	-	-	-	-	-	-	-	-	G	F	E	-	-	-	-	-	-	-		-	-	-	D	C I	ΒA
Rese	et valu	e		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0 0	0 0
ID	RW	Field	Value ID	Va	lue			De	scri	ptie	on																						
А	RW	RESETPIN		1				Re	eset	froi	n p	in de	eteo	cted	I																		
В	RW	DOG		1				Re	eset	froi	n w	atch	ndo	g de	etec	ted																	
С	RW	SREQ		1				Re	eset	froi	n A	IRCF	R.SY	SRE	SET	REC	) det	tect	ted														
D	RW	LOCKUP		1				Re	eset	fror	n C	PU l	ock	-up	dete	ecte	ed																
E	RW	OFF		1								wak al fro		•	om O	OFF	mc	de	whe	en w	/ake	eupi	is tr	igge	ered	fro	m t	he					
F	RW	LPCOMP		1										•	om LPC			de	whe	en w	/ake	eupi	is tr	igge	ered	fro	m						
G	RW	DIF		1								wak nterf		•	om ode	OFF	<sup>:</sup> mo	de	whe	en w	/ake	eupi	is tr	igge	ered	fro	m e	nte	rin	g			

### 11.2.2 SYSTEMOFF

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	2 1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-			-	Α
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A W		1				Er	nter	Sys	tem	OF	Fm	ode	è																			

## 11.2.3 GPREGRET

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	32	: 1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	Α.	A	A /	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	D	0	0	o (	0	0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A RW						Ge	ene	ral p	ourp	ose	ret	enti	on r	egis	ster.	Thi	s re	gist	er is	a re	etai	ned	regi	ste	er.							



### 11.2.4 **POFCON**

Bit	num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6 :	5	4	3	2	10
ID	(Field	I ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		•		-	-	-	B	ΒA
Re	set va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	D	0 (	0	0	0	0	0 0
ID	RW	Field	Value ID	Va	lue			De	scri	pti	on																							
Α	RW	POF		0				D	isab	le p	owe	er fa	ilur	e co	mp	arat	or																	
				1				Er	nabl	e po	owe	r fai	lure	e cor	npa	rato	or																	
В	RW	THRESHOLD	V21	0				Se	et th	res	holc	to i	2.1	V																				
			V23	1				Se	et th	res	holc	l to :	2.3	V																				
			V25	2				Se	et th	res	holc	to i	2.5	V																				
			V27	3				Se	et th	res	hold	l to :	2.7	V																				

### 11.2.5 RAMON

The RAM is divided into separate blocks for power management purposes. These blocks are not related in any way to the region concept described in the memory chapter.

			•	0
ID (Field ID) H G F E	D	c	В	A
Reset value 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	1	1	1	1
ID RW Field Value Description				
A RW ONRAMO 0 Keep RAM block 0 off in ON mode				
1 Keep RAM block 0 on in ON mode				
B RW ONRAM1 0 Keep RAM block 1 off in ON mode				
1 Keep RAM block 1 on in ON mode				
C RW ONRAM2 0 Keep RAM block 2 off in ON mode				
1 Keep RAM block 2 on in ON mode				
D RW ONRAM3 0 Keep RAM block 3 off in ON mode				
1 Keep RAM block 3 on in ON mode				
E RW OFFRAMO 0 Keep RAM block 0 off in OFF mode				
1 Keep RAM block 0 on in OFF mode				
F RW OFFRAM1 0 Keep RAM block 1 off in OFF mode				
1 Keep RAM block 1 on in OFF mode				
G RW OFFRAM2 0 Keep RAM block 2 off in OFF mode				
1 Keep RAM block 2 on in OFF mode				
H RW OFFRAM3 0 Keep RAM block 3 off in OFF mode				
1 Keep RAM block 3 on in OFF mode				



## 11.2.6 RESET

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	2 1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-			-	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	1
ID RW Field Value ID	Va	lue			De	scri	ptie	on																							
A RW							•				lebu reg	-		face	mc	ode,	see	the	DIF	per	riph	eral	reg	gist	er.	Thi	is				
	0				Di	sab	le																								
	1				Er	abl	e																								

### 11.2.7 DCDCEN

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	21	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		· -	Α
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	) (	0
ID RW Field Value ID	Va	lue			De	scri	pti	on																							
A RW					Er	nabl	e D	C/D	C cc	onve	ertei	r.																			
	0				Di	sab	le																								
	1				Er	nabl	e																								



# 12 Clock management (CLOCK)

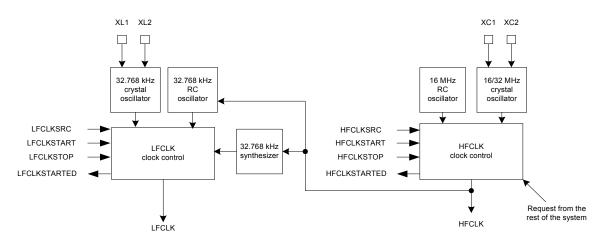


Figure 14 Clock control

## 12.1 Functional description

The system depends on, and generates, two different clocks: a high frequency clock (HFCLK) and a low frequency clock (LFCLK). These clocks are only available when the system is in ON mode.

The HFCLK is fixed to 16 MHz and the LFCLK is fixed to 32.768 kHz.

### 12.1.1 Low frequency clock (LFCLK)

The system supports three LFCLK clock sources: the 32.768 kHz crystal oscillator, the 32.768 kHz RC oscillator, and the 32.768 kHz synthesized clock, see *Figure 14*. The 32.768 kHz crystal oscillator requires an external AT-cut quartz crystal to be connected to the XL1 and XL2 pins in parallel resonant mode. The XL1 and XL2 share pins with the GPIO.

**Note:** GPIOs that share pins with XL1 and XL2 differ from device to device. For more information, see the device specific product specification.

The LFCLK clock and all of the available LFCLK sources are switched off by default when the system is propagated from OFF to ON mode.

The LFCLK clock is started by first selecting the preferred clock source in the LFCLKSRC register and then triggering the LFCLKSTART task. If the selected clock source cannot be started immediately the 32.768 kHz RC oscillator will start automatically and generate the LFCLK until the selected clock source is available.

The LFCLK is stopped by triggering the LFCLKSTOP task. The LFCLKSRC register should only be modified when the LFCLK is not running.

The 32.768 kHz crystal oscillator utilizes an amplitude regulated architecture to achieve low current consumption and fast start-up. The 32.768 kHz crystal oscillator is also designed to work with one of the following alternative external sources:

- A rail-to-rail clock signal applied to the XL1 pin. The XL2 pin shall then be left unconnected.
- A low swing clock signal applied to the XL1 pin. The XL2 pin shall then be left unconnected.



The synthesized 32.768 kHz clock depends on the HFCLK to run. If 250 ppm accuracy is required for the LFCLK running off the synthesized 32.768 kHz clock, the HFCLK must be generated from the 16/32 MHz crystal oscillator.

### 12.1.2 High frequency clock (HFCLK)

The system supports two high frequency clock sources: the 16/32 MHz crystal oscillator and the 16 MHz RC oscillator, see *Figure 14* on page 53. The HFCLK (16/32 MHz) crystal oscillators require an external AT-cut quartz crystal to be connected to the **XC1** and **XC2** pins in parallel resonant mode. If a 32 MHz crystal is used the XTALFREQ register must be configured accordingly.

When the system enters ON mode, the 16 MHz RC oscillator will start up automatically to provide the HFCLK to the CPU and other active parts of the system.

The HFCLK crystal oscillator is started by triggering the HFCLKSTART task and stopped using the HFCLKSTOP task. A HFCLKSTARTED event will be generated when the selected HFCLK crystal oscillator has started. The start-up times of the 16 MHz and 32 MHz crystal oscillators are described in the device specific product specification.

The 16 MHz RC oscillator is automatically switched off when one of the HFCLK crystal oscillators is running; it will be switched back on automatically when the HFCLK crystal oscillator is stopped.

If the system does not require a 16 MHz clock, the 16 MHz RC oscillator may be switched off automatically to save power. This occurs if all peripherals that require the HFCLK are appropriately stopped or disabled, and the CPU is sleeping. When this condition is no longer met the 16 MHz RC oscillator is automatically restarted. These optimization steps are only performed when the HFCLK is generated from the 16 MHz RC oscillator.

To use the RADIO and the calibration mechanism associated with the 32.768 kHz RC oscillator<mark>, the HFCLK must be generated from a HFCLK crystal</mark> oscillator.

The HFCLK crystal oscillators utilize amplitude regulated architecture to achieve low current consumption and fast start-up. The HFCLK crystal oscillators are also designed to work with one of the following alternative external sources:

- A 16 MHz rail-to-rail clock signal applied to the **XC1** pin. The **XC2** pin shall then be left unconnected.
- A 16MHz low swing clock signal applied to the **XC1** pin. The **XC2** pin shall then be left unconnected.

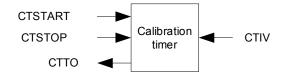
### 12.1.3 Calibrating the 32.768 kHz RC oscillator

After the 32.768 kHz RC oscillator is started and running, it can be calibrated triggering the CAL task. The 32.768 kHz RC oscillator will then temporarily request the HFCLK to calibrate itself against. A DONE event will be generated when calibration has finished. The calibration mechanism will only work as long as HFCLK is generated from the 16/32MHz crystal oscillator. Therefore, it is necessary to explicitly start this crystal oscillator before calibration can be started. See the device product specification for recommendations on calibration intervals and crystal accuracy.

#### 12.1.3.1 Calibration timer

The calibration timer can be used to time the calibration interval of the 32.768 kHz RC oscillator. The calibration timer is started by triggering the CTSTART task and stopped by triggering the CTSTOP task. The calibration timer will always start counting down from the value specified in CTIV and generate a CTTO timeout event when it reaches 0. The Calibration timer will stop by itself when it reaches 0.





*Figure 15 Calibration timer* 

Due to limitations in the calibration timer, only one task related to calibration, that is, CAL, CTSTART and CTSTOP, can be triggered for every period of LFCLOCK.

## 12.2 Registers

Register	Offset	Description
TASKS		
HFCLKSTART	0x000	Start HFCLK crystal oscillator
HFCLKSTOP	0x004	Stop HFCLK crystal oscillator
LFCLKSTART	0x008	Start LFCLK source
LFCLKSTOP	0x00C	Stop LFCLK source
CAL	0x010	Start calibration of LFCLK RC oscillator
CTSTART	0x014	Start calibration timer
CTSTOP	0x018	Stop calibration timer
EVENTS		
HFCLKSTARTED	0x100	16 MHz oscillator started
LFCLKSTARTED	0x104	32 kHz oscillator started
DONE	0x10C	Calibration of LFCLK RC oscillator complete event
СТТО	0x110	Calibration timer timeout
REGISTERS		
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
HFCLKSTAT	0x40C	Which HFCLK source is running
LFCLKSTAT	0x418	Which LFCLK source is running
LFCLKSRC	0x518	Clock source for the 32 kHz clock
CTIV	0x538	Calibration timer interval
XTALFREQ	0x550	Crystal frequency

Table 13Register overview



## 12.2.1 HFCLKSTAT

Bit n	uml	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	21	0
ID (F	ield	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	В	-	-	-	-	-	-	-	-	-	-	-	-			A
Rese	t va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0 0
ID R	wı	Field	Value ID	Va	lue			De	scri	pti	on																							
A F	2	SRC						Ad	tive	e Clo	ock :	soui	rce																					
			RC	0				16	5 MF	lz R	Cos	scilla	ator	run	nin	g ar	nd g	ene	rati	ng t	he F	IFCI	_K											
			XTAL	1				16	6/32	Мŀ	lz cr	ysta	al os	scilla	ator	run	ning	g ar	nd g	enei	ratir	ng tl	ne F	HFCL	LΚ									
ΒF	2	STATE						HI		۲nc	ate ot ru nnir		ng																					
			NOTRUNNING	0																														
			RUNNING	1																														

## 12.2.2 LFCLKSRC

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
A RW SRC						CI	ock	sou	irce																								
	RC	0				32	2.76	8 kŀ	lz R	Co	scill	ator	•																				
	XTAL	1				32	2.76	8 kŀ	Ηz c	ryst	al o	scill	atoı																				
	SYNTH	2				32	2.76	8 kŀ	Hz s	yntł	hesi	zer	synt	hes	izin	g 32	2.76	8 k⊦	lz fr	om	16	MH	z sys	ter	n c	loc	k						

## 12.2.3 LFCLKSTAT

Bit number	31 30 29 2	3 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2	210
ID (Field ID)			AA
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	00
ID RW Field Value ID	Value	Description	
A R SRC		Active Clock source	
RC	0	32.768 kHz RC oscillator running and generating the LFCLK	
XTAL	1	32.768 kHz crystal oscillator running and generating the LFCLK	
SYNTH	2	32.768 kHz synthesizer synthesizing 32.768 kHz from 16 MHz system clock	
B R STATE		LFCLK state	
NOTRUN	IING 0	LFCLK not running	
RUNNING	1	LFCLK running	



## 12.2.4 CTIV

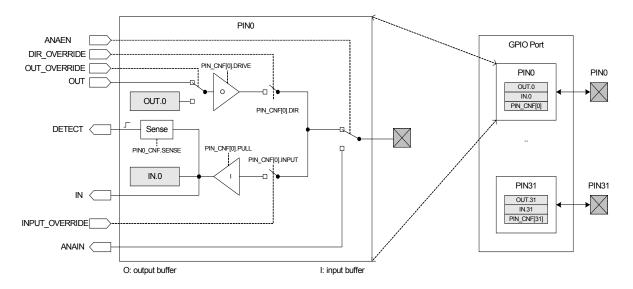
Bit number	31 30 29 28	27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field Value ID	Value	Description
A RW		Calibration timer interval in multiples of 0.25 seconds. Range: 0.25 seconds to 31.75 seconds.

## 12.2.5 XTALFREQ

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	A A	A A
Reset value	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11	1
ID RW Field Value ID	Va	lue			De	scri	ptio	on																							
A RW										•							al fo havi		CLK	K. Th	nis re	egis	ter	has	s to	ma	atcł	n th	ıe a	ctu	ıal
	0	ĸFF			16	5 MF	Ηz c	ryst	al is	use	ed.																				
	0	<00			32	2 MF	Ηz c	ryst	al is	use	ed.																				



# 13 GPIO



*Figure 16* GPIO Port and the GPIO pin details

*Figure 16* illustrates the GPIO port containing 32 individual pins, where **PIN0** is illustrated in more detail as a reference. All the signals on the left side of the illustration are used by other peripherals in the system, and therefore, are not directly available to the CPU.

## 13.1 Functional description

The GPIO Port peripheral implements up to 32 pins, **PIN0** through **PIN31**. Each of these pins can be individually configured in the PIN\_CNF[n] registers (n=0..31). The following parameters can be configured through these registers:

- Direction
- Drive strength
- Enabling of pull-up and pull-down resistors
- Pin sensing
- Input buffer disconnect
- Analog input (for selected pins)

The PIN[n].CNF registers are retained registers. See *Chapter 11 "Power management (POWER)"* on page 40 for more information about retained registers.

Pins can be individually configured, through the pin sense mechanism, to detect either a high level or a low level on their input. When the correct level is detected, the sense mechanism raises the DETECT signal line, which can then be read by other peripherals in the system, see *Figure 16*. This mechanism is functional in both ON and OFF mode.

The input buffer of a GPIO pin can be disconnected from the pin to enable power savings when the pin is not used as an input, see *Figure 16.* Inputs must be connected in order to get a valid input value in the IN register and for the sense mechanism to get access to the pin.

Other peripherals in the system can attach themselves to GPIO pins and override their output value and configuration, or read their analog or digital input value, see *Figure 16*.



Selected PINs also support analog input signals, see ANAIN in *Figure 16* on page 58. Pins that support analog input signals vary between devices, see the product specification for your device for more details.

Pin direction can be configured both in the DIR register as well as through the individual PIN\_CNF[n] registers. A change in one register will automatically be reflected in the other register.

## 13.2 Registers

Register	Offset	Description
REGISTERS		
OUT	0x504	Write GPIO port
OUTSET	0x508	Set individual bits in GPIO port
OUTCLR	0x50C	Clear individual bits in GPIO port
IN	0x510	Read GPIO port
DIR	0x514	Direction of GPIO pins
DIRSET	0x518	Setting DIR register
DIRCLR	0x51C	Clearing DIR register
PIN_CNF[0]	0x700	Configuration of pin 0
PIN_CNF[31]	0x77C	Configuration of pin 31

 Table 14
 Register overview

### 13.2.1 OUT

Bit numb	per		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW	Field	Value ID	Va	lue			De	scri	ipti	on																								
A RW								-						wha g. bi			•		•				-		rela	tes	s to	pir	n					

### 13.2.2 IN

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW	Field	Value ID	Va	lue			De	scri	ptio	on																								
A R								5						ole G elate		•		•				egis	ter	rela	tes	to	pin	nu	mt	ber				



## 13.2.3 OUTSET

Bit number	3	81 3	0 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)	ŀ	A A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset value	C	) ()	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field	Value ID \	/alue	•		De	scri	pti	on																								
A RW						-			•				PIO it 0 i	•								-					•					

### 13.2.4 OUTCLR

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW	Field	Value ID	Va	alue			De	scri	ipti	on																								
A RW							nı	umb	ber i		PIO	por	t, e.	the g. bi	t 0 r	relat	tes t	o G	PIO	pin	nur	nbe	er 0.	Set	ting	g a	<b>'1'</b> i	in c						

### 13.2.5 DIR

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α	A	A	A	A	A	A	A	A	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
A RW						ու th	imb e bi	oer i ts ir	n Gl n the	PIO e re	por gist	t, e. er w	g. bi rill co	t 0 i onfi	relat gur	tes t	to G e cc	PIO prres	pin spoi	nuı ndir	mbe ng G	er 0. GPIC	egist Sett pin t.	ting	g a	'1' i	in o	ne	of				

## 13.2.6 DIRSET

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)		A	A	A	A	A	A	A	Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	<b>\</b>
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	) (
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																								
A RW								-							ts in outs			_	-					•					erv	will				



## 13.2.7 DIRCLR

Bit number	3	1 30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field Value ID	V	alue			De	scri	ptio	on																								
A RW													bits ut. Se														wil					

## 13.2.8 PIN\_CNF[n] (n=0..31)

Bit	numl	ber		31	30	29	28	27	7 26	52	25 2	4 2	23 2	22 2	12	20 1	9 1	8 17	7 10	6 1	5 14	13	12	11	10	9	8	7 (	6 !	54	3	2	1	0
ID (	Field	ID)		-	-	-	-	-	-	-	-	-			-	-	-	E	E	-	-	-	-	-	D	D	D				C	c	В	A
Res	et va	lue		0	0	0	0	0	0	0	) 0	0	) (	) 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	) (	0	0	1	0
ID	RW	Field	Value ID	Va	lue			D	escr	rip	tion																							
А	RW	DIR						Р	'in d	lire	ectio	n																						
			INPUT	0				C	Confi	Ϊgι	ure p	oin a	as a	n in	out	pin																		
			OUTUT	1				C	Confi	Ϊgι	ure p	oin a	as a	n ou	tpι	ıt pir	۱																	
В	RW	INPUT						C	Conn	nec	t or	dis	con	inect	: inp	out k	ouff	er																
			CONNECT	0				C	Conn	nec	t in	out	but	ffer																				
			DISCONNECT	1				D	Disco	onr	nect	inp	out	buffe	er																			
с	RW	PULL																																
			DISABLED	0				Ν	lo p	ull																								
			PULLDOWN	1				Ρ	ull c	dov	wn c	on p	oin																					
			PULLUP	3				Ρ	ull u	up	on p	oin																						
D	RW	DRIVE																																
			S0S1	0				S	tand	dar	rd 0,	sta	anda	ard 1																				
			H0S1	1				F	ligh	dr	ive (	), st	tand	dard	1																			
			S0H1	2				S	tand	dar	rd 0,	hig	gh c	lrive	1																			
			H0H1	3				F	ligh	dr	ive (	), h	igh	driv	e 1																			
			D0S1	4				C	Disco	onr	nect	0, s	star	dard	11																			
			D0H1	5				C	Disco	onr	nect	0, ł	higł	n driv	/e 1																			
			S0D1	6				S	tand	dar	rd 0,	dis	scor	nnec	t 1																			
			H0D1	7										onne																				
E	RW	SENSE						P	in se	ens	sing	me	echa	anisr	n																			
			DISABLED	0				C	Disak	ble	d																							
			HIGH	2				S	ense	e fo	or hi	igh	lev	el																				
			LOW	3				S	ense	e fo	or lo	wl	eve																					



# 14 GPIO tasks and events (GPIOTE)

## 14.1 Functional description

The GPIO Tasks and Events (GPIOTE) module provides functionality for accessing GPIO pins using tasks and events.

A task can be used for performing the following write operations to a pin:

- Set
- Clear
- Toggle

An event can be generated from any of the following input pins using the GPIO DETECT signal:

- Rising edge
- Falling edge
- Any change

#### 14.1.1 Pin events and tasks

The GPIOTE module has a number of tasks and events that can be configured to operate on individual GPIO pins; the OUT[n] tasks and the IN[n] events. The tasks can be used for writing to individual pins, and the events can be generated from changes occurring at the inputs of individual pins.

The tasks and events are configured using the CONFIG[n] registers. Every pair of OUT[n] tasks and IN[n] events has one CONFIG[n] register associated with it.

When an OUT[n] task or an IN[n] event has been configured to operate on a pin, the pin can only be written from the GPIOTE module. Attempting to write a pin as a normal GPIO pin will have no effect.

As long as an OUT[n] task or an IN[n] event is configured to control a pin **n**, the pin's output value will only be updated by the GPIOTE module. The pin's output value as specified in the GPIO will therefore be ignored as long as the pin is controlled by GPIOTE. When the GPIOTE is disconnected from a pin, see MODE field in CONFIG[n] register, the associated pin will get the output and configuration values specified in the GPIO module.

When a GPIOTE channel is configured to operate on a pin as a task, the initial value of that pin is configured in the OUTINIT field of CONFIG[n].

#### 14.1.2 Port event

PORT is an event that can be generated from multiple input pins using the GPIO DETECT signal. The event will be generated on the rising edge of the DETECT signal. See *Section 13.1 "Functional description"* on page 58 for more information about the DETECT signal.

This feature is always enabled although the peripheral itself appears to be IDLE, that is, no clocks or other power intensive infrastructure have to be requested to keep this feature enabled. This feature can therefore be used to wake-up the CPU from a WFI or WFE type sleep in System ON with all peripherals and the CPU idle, that is, lowest power consumption in System ON mode.



# 14.2 Registers

Registers	Offset	Description
TASKS		
OUT[0]	0x000	Task for writing to pin specified by PSEL in CONFIG[0].
OUT[1]	0x004	Task for writing to pin specified by PSEL in CONFIG[1].
OUT[2]	0x008	Task for writing to pin specified by PSEL in CONFIG[2].
OUT[3]	0x00C	Task for writing to pin specified by PSEL in CONFIG[3].
EVENTS		
IN[0]	0x100	Event generated from pin specified by PSEL in CONFIG[0].
IN[1]	0x104	Event generated from pin specified by PSEL in CONFIG[1].
IN[2]	0x108	Event generated from pin specified by PSEL in CONFIG[2].
IN[3]	0x10C	Event generated from pin specified by PSEL in CONFIG[3].
PORT	0x17C	Event generate from multiple input pins.
REGISTERS		
INTENSET	0x304	Interrupt enable set register.
INTENCLR	0x308	Interrupt enable clear register.
CONFIG[0]	0x510	Configuration for OUT[0] task and IN[0] event.
CONFIG[1]	0x514	Configuration for OUT[1] task and IN[1] event.
CONFIG[2]	0x518	Configuration for OUT[2] task and IN[2] event.
CONFIG[3]	0x51C	Configuration for OUT[3] task and IN[3] event.

 Table 15
 Register overview

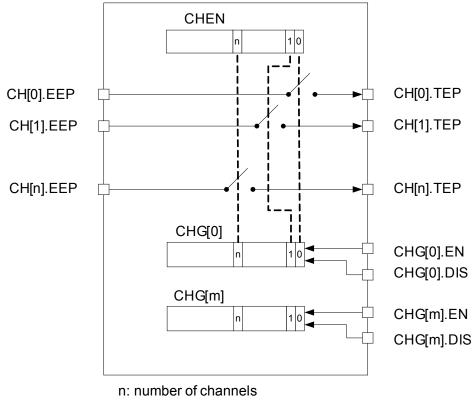


# 14.2.1 CONFIG[n] (n=0..3)

Bi	t num	ber		31	1 30	) 29	<b>9</b> 2	8 2	27 2	62	25 2	4 2	3 2	2 21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	65	5 4	13	2	10
ID	(Field	ID)		-	-	-	-	-	-	-	-	-	-	-	D	-	-	С	С	-	-	-	В	В	В	В	В	-			· -	- 1	A A
Re	eset va	lue		0	0	0	0	(	) 0	0	) ()	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0 0	0 0	0 (	0 0
ID	RW	Field	Value ID	Va	alue	2		1	Desc	rip	tion	ı																					
A	RW	MODE							Mod	e																							
			DISABLED	0	)				Disa	ble	d. Pi	in sp	peci	fied	by P	SEL	will	not	be	acq	luire	ed b	y th	e G	PIO	TE n	noc	dule	<b>.</b>				
			EVENT	1					Even IN[n] pin.				•	•			-					-				•							
			TASK	3					Task trigg the p pin c	jeri oin.	ing t . Wh	he (	OUT enal	[n] ta pled	ask v as a	vill p task	oerf « the	orm e GP	the NOT	e op E m	era Iodu	tion ule v	spe vill a	ecifi acqu	ed k uire	y P the	OL pii	ARI n an	d tł				
В	RW	PSEL		[(	031	]]			Pin r	nun	nbei	rass	socia	ted	witł	n OU	T[n]	] tas	k ar	nd I	N[n	eve	ent.										
C	RW	POLARITY							Whe trigg Whe	jere	ed.			•				•					•				-	-	ask i	is			
			LOTOHI	1					Task Ever				•						n ri	sing	) ed	ge o	on p	in.									
			HITOLO	2					Task Ever				•							Illin	g eo	lge	on J	oin.									
			TOGGLE	3					Task Ever			-		•					cha	nge	on	pin											
C	RW	OUTINIT							Whe conf				node	: Init	ial v	alue	e of	the	out	put	wh	en t	he (	GPIC	OTE	cha	nn	el is	5				
									Whe	n iı	n ev	ent	mod	de: N	lo ef	fect	•																
			LOW	0	)				Task	ma	ode:	Init	ial v	alue	ofp	oin b	efo	re ta	ask t	trig	geri	ng i	s lo	N.									
			HIGH	1					Task	ma	ode:	Init	ial v	alue	ofp	oin b	efo	re ta	ask 1	rig	geri	ng i	s hig	gh.									



# 15 Programmable Peripheral Interconnect (PPI)



m: number of channel groups

*Figure 17 PPI block diagram* 

## 15.1 Functional description

The Programmable Peripheral Interconnect (PPI) enables peripherals to interact autonomously with each other using tasks and events independent of the CPU.

The PPI provides a mechanism to automatically trigger a task in one peripheral as a result of an event occurring in another peripheral. A task is connected to an event through a PPI channel. The PPI channel is composed of two end-point registers, the Event End-Point (EEP) and the Task End-Point (TEP). A peripheral task is connected to a Task End-Point using the address of the task register associated with the task. Similarly, a peripheral event is connected to an Event End-Point using the address of the event register associated with the event.

There are two ways of enabling and disabling PPI channels:

- Enable or disable PPI channels individually using the CHEN, CHENSET, and CHENCLR registers.
- Enable or disable PPI channels in PPI channel Groups through the Groups' ENABLE and DISABLE tasks. Prior to these tasks being triggered, the PPI channel Group must be configured to define which PPI channels belongs to which groups.



PPI tasks (for example, CHG0EN) can be triggered through the PPI like any other task, which means they can be hooked up to a PPI channel as a TEP. One event can trigger multiple tasks by using multiple channels and one task can be triggered by multiple events in the same way.

### 15.1.1 Pre-programmed channels

The PPI system has in addition to the fully programmable peripheral interconnections, a set of channels where the event (EEP) and task (TEP) endpoints are set in hardware. These fixed channels can be individually enabled, disabled, or added to PPI channel groups in the same way as ordinary PPI channels.

Channel	EEP	ТЕР
20	TIMER0->EVENTS_COMPARE[0]	RADIO->TASKS_TXEN
21	TIMER0->EVENTS_COMPARE[0]	RADIO->TASKS_RXEN
22	TIMER0->EVENTS_COMPARE[1]	RADIO->TASKS_DISABLE
23	RADIO->EVENTS_BCMATCH	AAR->TASKS_START
24	RADIO->EVENTS_READY	CCM->TASKS_KSGEN
25	RADIO->EVENTS_ADDRESS	CCM->TASKS_CRYPT
26	RADIO->EVENTS_ADDRESS	TIMER0->TASKS_CAPTURE[1]
27	RADIO->EVENTS_END	TIMER0->TASKS_CAPTURE[2]
28	RTC0->COMPARE[0]	RADIO-> TASKS_TXEN
29	RTC0->COMPARE[0]	RADIO-> TASKS_RXEN
30	RTC0->COMPARE[0]	TIMER0->TASKS_CLEAR
31	RTC0->COMPARE[0]	TIMER0->TASKS_START

 Table 16
 Pre-programmed channels

## 15.2 Registers

Register	Offset	Description	
TASKS			
CHG[0].EN	0x000	Enable channel group 0	
CHG[0].DIS	0x004	Disable channel group 0	
CHG[1].EN	0x008	Enable channel group 1	
CHG[1].DIS	0x00C	Disable channel group 1	
CHG[2].EN	0x010	Enable channel group 2	
CHG[2].DIS	0x014	Disable channel group 2	
CHG[3].EN	0x018	Enable channel group 3	
CHG[3].DIS	0x01C	Disable channel group 3	
REGISTERS			
CHEN	0x500	Channel enable	
CHENSET	0x504	Channel enable set	
CHENCLR	0x508	Channel enable clear	
CH[0].EEP	0x510	Channel 0 event endpoint	
CH[0].TEP	0x514	Channel 0 task endpoint	



Register		Description
CH[1].EEP	0x518	Channel 1 event endpoint
CH[1].TEP	0x51C	Channel 1 task endpoint
CH[2].EEP	0x520	Channel 2 event endpoint
CH[2].TEP	0x524	Channel 2 task endpoint
CH[3].EEP	0x528	Channel 3 event endpoint
CH[3].TEP	0x52C	Channel 3 task endpoint
CH[4].EEP	0x530	Channel 4 event endpoint
CH[4].TEP	0x534	Channel 4 task endpoint
CH[5].EEP	0x538	Channel 5 event endpoint
CH[5].TEP	0x53C	Channel 5 task endpoint
CH[6].EEP	0x540	Channel 6 event endpoint
CH[6].TEP	0x544	Channel 6 task endpoint
CH[7].EEP	0x548	Channe 7 event endpoint
CH[7].TEP	0x54C	Channel 7 task endpoint
CH[8].EEP	0x550	Channel 8 event endpoint
CH[8].TEP	0x554	Channel 8 task endpoint
CH[9].EEP	0x558	Channel 9 event endpoint
CH[9].TEP	0x55C	Channel 9 task endpoint
CH[10].EEP	0x560	Channel 10 event endpoint
CH[10].TEP	0x564	Channel 10 task endpoint
CH[11].EEP	0x568	Channel 11 event endpoint
CH[11].TEP	0x56C	Channel 11 task endpoint
CH[12].EEP	0x570	Channel 12 event endpoint
CH[12].TEP	0x574	Channel 12 task endpoint
CH[13].EEP	0x578	Channel 13 event endpoint
CH[13].TEP	0x57C	Channel 13 task endpoint
CH[14].EEP	0x580	Channel 14 event endpoint
CH[14].TEP	0x584	Channel 14 task endpoint
CH[15].EEP	0x588	Channel 15 event endpoint
CH[15].TEP	0x58C	Channel 15 task endpoint
CHG[0]	0x800	Channel group 0
CHG[1]	0x804	Channel group 1
CHG[2]	0x808	Channel group 2
CHG[3]	0x80C	Channel group 3



### 15.2.1 CHEN

Bit n	umb	er		31	30	29	28	27	26	25	24	23	3 22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (F	ield	ID)		FF	EE	DD	сс	BB	AA	z	Y	Х	W	v	U	-	-	-	-	Ρ	0	Ν	м	L	К	J	I	н	G	F	Е	D	с	В	A
Rese	et val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID	RW	Field	Value ID		Va	alue											C	)esc	ript	tion	1														
A	R W	CH0						Ena	able	or	disal	ble	cha	nne	10																				
				0				Dis	able	•																									
				1				Ena	able																										
В		CH1																																	
С		CH2																																	
D		CH3																																	
Е		CH4																																	
F		CH5		••																															
G		CH6																																	
Н		CH7		••																															
I		CH8																																	
J	•	CH9		••				••																											
К		CH10																																	
L		CH11		••																															
М		CH12		••																															
Ν	••	CH13		••																															
0		CH14		••																															
P		CH15		••				••																											
U		CH20																																	
V	••	CH21		••																															
W		CH22																																	
X	•	CH23		••				•																											
Y Z		CH24 CH25																																	
AA		CH25		••				•																											
BB		CH20																																	
CC		CH28						••																											
DD		CH28																																	
EE		CH29 CH30		•				••																											
FF																																			
FF		CH31		••				••																											



## 15.2.2 CHENSET

Bit n	umb	er		31	30	29	28	27	26	25	5 24	23	3 22	2 21	20	19	18	3 17	16	5 15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (F	ield I	D)		FF	EE	DD	сс	BB	AA	z	Y	Х	w	v	U	-	-	-	-	Ρ	ο	Ν	м	L	к	J	I	н	G	F	E	D	с	В	A
Rese	t val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID	RW	Field	Value ID	Val	ue			De	scrip	otio	n																								
А		CH0						En	able	e ch	ann	el 0																							
	W			1				En	able	2																									
	R			CH	IEN	.CHC	)	Re	ad v	alu	e of	СН	0 fie	eld i	n CH	IEN	reg	giste	r																
В		CH1																																	
С		CH2																																	
D		CH3																																	
E		CH4																																	
F		CH5																																	
G		CH6																																	
Н		CH7						••																											
I		CH8																																	
J		CH9						•																											
К		CH10																																	
L	••	CH11		•				•																											
M		CH12																																	
N O		CH13 CH14		•				•																											
P	•	CH15																																	
U	•	CH20		••																															
v	•	CH21																																	
W		CH22																																	
x		CH23																																	
Y		CH24																																	
z	•	CH25																																	
AA		CH26																																	
BB		CH27																																	
CC		CH28						•																											
DD		CH29																																	
EE		CH30																																	
FF	•	CH31																																	
		Chor		••				••																											



## 15.2.3 CHENCLR

Bit r	umb	er		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (F	ield	ID)		FF	EE	DD	сс	BB	AA	z	Y	х	w	v	U	-	-	-	-	Ρ	0	Ν	м	L	к	J	I	н	G	F	E	D	C I	3 A
Rese	et val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0 0
ID	RW	Field	Value ID	Va	lue			De	scrip	tior	ı																							
А		CH0						En	able	cha	nnel	0																						
	W			1				En	able																									
	R			Cł	HEN.	СНО	)	Re	ad v	alue	of C	H0	fiel	d in	CHE	EN r	egist	ter																
В		CH1																																
С		CH2																																
D		CH3						••																										
Е		CH4																																
F		CH5		••				•																										
G		CH6																																
Н		CH7																																
1		CH8																																
J		CH9		••				••																										
ĸ		CH10																																
L		CH11		••				••																										
M N		CH12 CH13																																
0	 	CH14		••																														
P		CH15																																
U		CH20																																
V		CH21																																
w		CH22																																
х		CH23																																
Y		CH24																																
Z		CH25																																
AA		CH26																																
BB		CH27																																
сс		CH28																																
DD		CH29																																
EE		CH30																																
FF		CH31																																



## 15.2.4 CH[n].EEP (n=0..15)

Event endpoints are only able to recognize addresses from the EVENT group.

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field	Value ID	Va	lue			De	scri	ipti	on																								
A RW							oint roup		o ev	ent	reg	iste	r. Ao	cep	ots c	only	ado	dres	ses	to r	egi	ster	s frc	m	the	e Ev	ver	nt					

### 15.2.5 CH[n].TEP (n=0..15)

Task endpoints are only able to recognize addresses from the TASK group.

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 :	3 3	21	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A i	A I	A A	AA
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	0 0	0 0
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																							
A RW							Pc	ointe	er to	o tas	sk re	egist	ter.	Acce	epts	on	ly a	ddre	esse	s to	reg	iste	rs fr	om	the	e Ta	ask	gro	ou	р.			

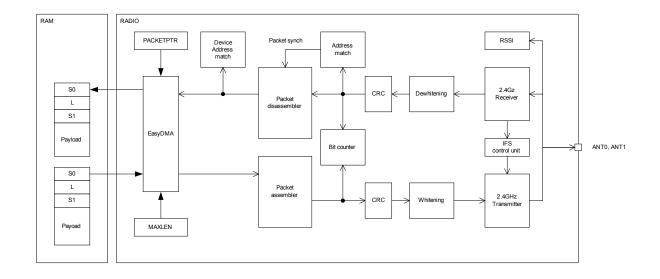


# 15.2.6 CHG[n] (n=0..3)

Reset																																
	valu		ID (Field ID)			DD	cc	BB	AA	z	Y	х	w v	U V	-	-	-	-	Ρ	0	N	м	L	к	Γ	н	G	F	E	D	С В	5 A
ID I		Reset value			0 (	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0 (	) 0	0	0	0	0 0	) (	0
	RW	Field	Value ID	Valu	ue			Des	crip	tior	ı																					
А	RW	CH0						Inc	lude	e or e	excl	ude	char	nnel	0 in	gro	up	n														
				0				Exc	lude	e																						
				1				Inc	lude	2																						
В	RW	CH1																														
С	RW	CH2																														
D	RW	CH3						••																								
	RW	CH4																														
	RW	CH5						••																								
	RW	CH6																														
	RW	CH7						••																								
	RW	CH8						••																								
	RW	CH9						••																								
	RW RW	CH10 CH11		•																												
	RW	CH12						••																								
	RW	CH13																														
	RW	CH14																														
	RW	CH15																														
U	RW	CH20																														
V	RW	CH21																														
W	RW	CH22																														
Х	RW	CH23																														
Y	RW	CH24																														
Z	RW	CH25																														
AA	RW	CH26																														
BB	RW	CH27																														
CC	RW	CH28																														
DD	RW	CH29						••																								
EE	RW	CH30																														
FF	RW	CH31						••																								



## 16 2.4 GHz radio (RADIO)



#### Figure 18 Radio block diagram

The RADIO contains a 2.4 GHz radio receiver and a 2.4 GHz radio transmitter that is compatible with Nordic's proprietary 2 Mbps, 1 Mbps, and 250 kbps radio modes in addition to 1 Mbps *Bluetooth* Low Energy mode.

The RADIO implements EasyDMA. EasyDMA in combination with an automated packet assembler and packet disassembler, and an automated CRC generator and CRC checker, makes it very easy to configure and use the RADIO. See *Figure 18* for more information.

The RADIO includes a Device Address Match unit and an interframe spacing control unit that can be utilized to simplify address white listing and interframe spacing respectively, in *Bluetooth* Low Energy and similar applications.

The RADIO also includes a Received Signal Strength Indicator (RSSI) and a bit counter. The bit counter generates events when a preconfigured number of bits have been sent or received by the RADIO.

## 16.1 Functional description

#### 16.1.1 EasyDMA

The RADIO implements EasyDMA for reading and writing of data packets from and to the RAM without CPU involvement.

As illustrated in *Figure 18*, the RADIO's EasyDMA utilizes the same PACKETPTR pointer for receiving packets and transmitting packets. The CPU should therefore reconfigure this pointer every time the radio is switched between transmit and receive mode. The MAXLEN register configures the maximum number of bytes that can be transmitted or received by the RADIO within the same packet. This feature can be used to ensure that the RADIO does not overwrite, or read beyond, the RAM assigned to the packet.

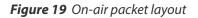
The EasyDMA is guaranteed to have finished accessing the RAM when the DISABLED event is generated.



## 16.1.2 Packet configuration

A radio packet contains the following fields: PREAMBLE, ADDRESS, LENGTH, SO, S1, PAYLOAD, and CRC. The radio sends the different fields in the packet in the order they are shown in *Figure 19*, from left to right. The preamble will be sent least significant bit first on-air.

0x55 = 10101010 0xAA= 01010101										
PREAMBLE	BASE	B PREFIX	S0	LENGTH	S1		PAYLOAD		~ Cl	RC
	LSByte	1	1	I	I	LSByte	I	I	MSByte	I
	ADDRESS		J							



The PREAMBLE is always one byte long taking the value 0xAA or 0x55 depending on the first ADDRESS bit that is sent on air. If the first bit of the ADDRESS is 0 the preamble is set to 0xAA. Otherwise, the PREAMBLE is set to 0x55.

Radio packets are stored in memory inside instances of a radio packet data structure as shown in *Figure 20*. The PREAMBLE, ADDRESS and CRC fields are omitted in this data structure since they are added dynamically when the packet is sent.



*Figure 20* In-RAM representation of radio packet, S0, LENGTH, and S1 are optional

The byte ordering on air is always least significant byte first for the ADDRESS and PAYLOAD fields and most significant byte first for the CRC field. The ADDRESS fields are always transmitted and received with least significant bit first on-air. The CRC field is always transmitted and received with the most significant bit first. The bit-endianness (which means the order the bits are sent and received in) of the S0, LENGTH, S1, and PAYLOAD fields can be configured through the ENDIAN field in PCNF1.

The sizes of the S0, LENGTH, and S1 fields can be individually configured through S0S, LS, and S1S in PCNF0 respectively. If any of these fields are configured to be less than 8 bits long, the least significant bits of the fields, as seen from the RAM representation, are used.

If S0, LENGTH, or S1 are specified with zero length, their fields will be omitted in memory. Otherwise, each field is represented as a separate byte, regardless of the number of bits in their on-air counterpart.

#### 16.1.3 Maximum packet length

Independent of the configuration of MAXLEN, the combined length of S0, LENGTH, S1, and PAYLOAD cannot exceed 255 bytes.

#### 16.1.4 Address configuration

The on-air radio ADDRESS field is composed of two parts: the base address field and the address prefix field, see *Figure 19*. The size of the base address field is configurable through BALEN in PCNF1. The base address is truncated from LSByte if the BALEN is less than 4.



The on-air addresses are defined in the BASEn and PREFIXn registers. For other radio address registers such as the TXADDRESS, RXADDRESSES and RXMATCH registers, logical radio addresses ranging from 0 to 7 are being used. The relationship between the on-air radio addresses and the logical addresses are described in *Table 17*.

Logical address	Base address	Prefix byte
0	BASE0	PREFIX0.AP0
1	BASE1	PREFIX0.AP1
2	BASE1	PREFIX0.AP2
3	BASE1	PREFIX0.AP3
4	BASE1	PREFIX1.AP4
5	BASE1	PREFIX1.AP5
6	BASE1	PREFIX1.AP6
7	BASE1	PREFIX1.AP7

Table 17 Definition of logical addresses

#### 16.1.5 Received Signal Strength Indicator (RSSI)

The radio implements a mechanism for measuring the power in the received radio signal. This feature is called Received Signal Strength Indicator (RSSI).

Sampling of the received signal strength is started by using the RSSISTART task. The sample can be read from the RSSISAMPLE register.

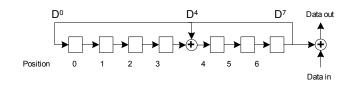
The sample period of the RSSI is defined by RSSI<sub>RESOLUTION</sub>, see the device specific product specification for details. The RSSI sample will hold the average received signal strength during this sample period.

For the RSSI sample to be valid the radio has to be enabled in receive mode (RXEN task) and the reception has to be started (READY event followed by START task).

#### 16.1.6 Data whitening

The RADIO is able to do packet whitening and de-whitening, see WHITEEN in PCNF1 register for how to enable whitening. When enabled, whitening and de-whitening are handled by the RADIO automatically as packets are sent and received, that is, radio packets located in RAM will not be whitened.

The whitening word is generated using polynomial  $g(D) = D^7 + D^4 + 1$ , which then is XORed with the data packet that is to be whitened, or de-whitened, see *Figure 21*.



*Figure 21* Data whitening and de-whitening.



Whitening and de-whitening will be performed over the whole packet, except for the preamble and the address field.

The linear feedback shift register, illustrated in *Figure 21* on page 75 can be initialized through the DATAWHITEIV register.

#### 16.1.7 CRC

The CRC generator in the RADIO calculates the CRC over the whole packet excluding the preamble. If desirable the address field can be excluded from the CRC calculation as well, see CRCCNF register for more information.

The CRC polynomial is configurable as illustrated in *Figure 22* where bit 0 in the CRCPOLY register corresponds to  $X^0$  and bit 1 corresponds to  $X^1$  etc. See CRCPOLY for more information.

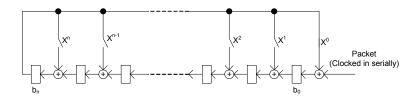


Figure 22 CRC generator for an n bit CRC

As illustrated in *Figure 17* on page 65, the CRC is calculated by feeding the packet serially through the CRC generator. Before the packet is clocked through the CRC generator, the CRC generator's latches  $b_0$  through  $b_n$  will be initialized with a predefined value specified in the CRCINIT register. When the whole packet is clocked through the CRC generator, latches  $b_0$  through  $b_n$  will hold the resulting CRC. This value will be used by the RADIO during both transmission and reception but it is not available to be read by the CPU at any time. A received CRC can however be read by the CPU through the RXCRC register independent of whether or not it has passed the CRC check.

The length (*n*) of the CRC is configurable, see CRCCNF for more information.

#### 16.1.8 Radio states

State Description DISABLED No operations are going on inside the radio and the power consumption is at a minimum. RXRU The radio is ramping up and preparing for reception. RXIDLE The radio is ready for reception to start. RX Reception has been started and the addresses enabled in the RXADDRESSES register are being monitored. TXRU The radio is ramping up and preparing for transmission. TXIDLE The radio is ready for transmission to start. ΤХ The radio is transmitting a packet.

The radio can enter the states described in *Table 18*.

Table 18 Radio states



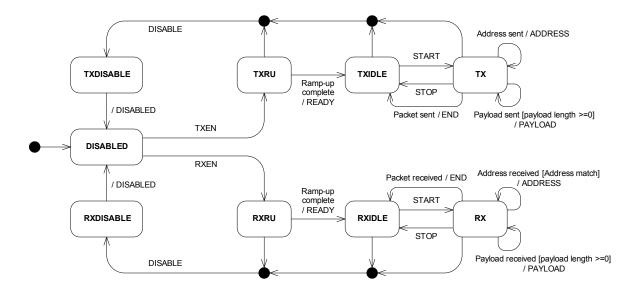


Figure 23 Radio state diagram

#### 16.1.9 Maximum consecutive transmission time

Maximum consecutive transmission time is defined as the longest time the RADIO can be active transmitting before it has to be disabled, that is, the longest possible time between READY event and DISABLE task.

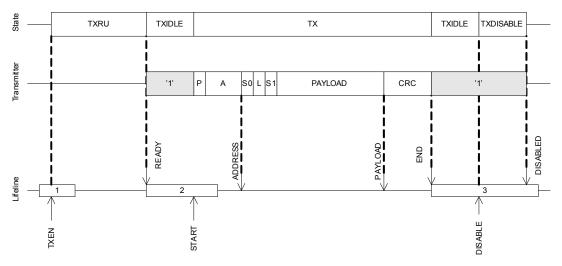
Maximum consecutive transmission time for the RADIO is 4 ms running off a 60 ppm crystal and 16 ms running off a 30 ppm crystal.



#### 16.1.10 Transmit sequence

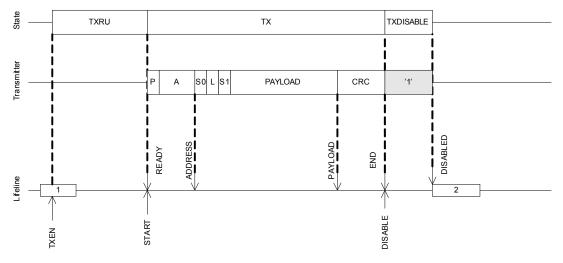
Before the RADIO is able to transmit a packet, it must first ramp-up in TX mode, see TXRU in *Figure 23* on page 77 and *Figure 24*. A TXRU ramp-up sequence is initiated when the TXEN task is triggered. After the radio has successfully ramped up it will generate the READY event indicating that a packet transmission can be initiated. A packet transmission is initiated by triggering the START task. As illustrated in *Figure 23* on page 77 the START task can first be triggered after the RADIO has entered into the TXIDLE state.

*Figure 24* illustrates a single packet transmission where the CPU manually triggers the different tasks needed to control the flow of the RADIO, that is, no shortcuts are used. If shortcuts are not used, a certain amount of delay caused by CPU execution is expected between READY and START, and between END and DISABLE. As illustrated in *Figure 23* on page 77 the RADIO will by default transmit '1's between READY and START and between READY and DISABLED.



*Figure 24 Transmit sequence* 

A slightly modified version of the transmit sequence from *Figure 24* is illustrated in *Figure 25* where the RADIO is configured to use shortcuts between READY and START, and between END and DISABLE, and therefore no delay is introduced.



*Figure 25 Transmit sequence using shortcuts to avoid delays.* 



The RADIO is able to send multiple packets one after the other without having to disable and re-enable the RADIO between packets, this is illustrated in *Figure 26*.

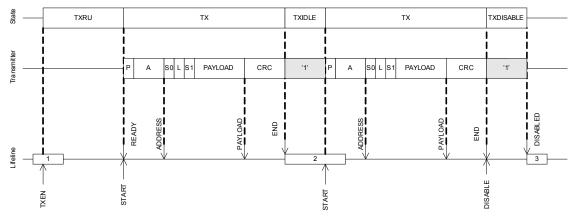


Figure 26 Transmission of multiple packets

#### 16.1.11 Receive sequence

Before the RADIO is able to receive a packet, it must first ramp-up in RX mode. An RXRU ramp-up sequence is initiated when the RXEN task is triggered. After the radio has successfully ramped up it will generate the READY event indicating that a packet reception can be initiated. A packet reception is initiated by triggering the START task. As illustrated in *Figure 27* the START task can first be triggered after the RADIO has entered into the RXIDLE state.

*Figure 28* on page 80 illustrates a single packet reception where the CPU manually triggers the different tasks needed to control the flow of the RADIO, that is, no shortcuts are used. If shortcuts are not used, a certain amount of delay, caused by CPU execution, is expected between READY and START, and between END and DISABLE. As illustrated in *Figure 28* on page 80 the RADIO will be listening and possibly receiving undefined data, illustrated with an 'X', from START and until a packet with valid preamble (P) is received.

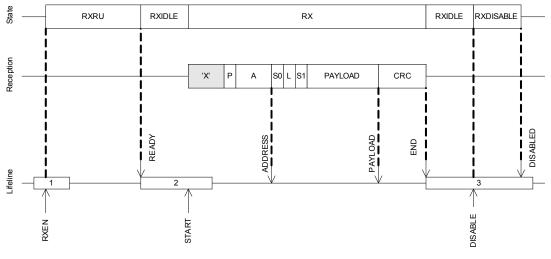


Figure 27 Receive sequence

A slightly modified version of the receive sequence from *Figure 27* is illustrated in *Figure 28* on page 80 where the RADIO is configured to use shortcuts between READY and START, and between END and DISABLE, and therefore no delay is introduced.



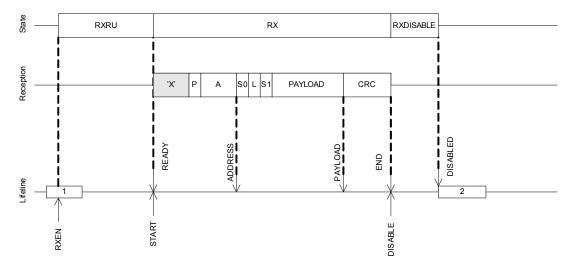


Figure 28 Receive sequence using shortcuts to avoid delays

The RADIO is able to receive multiple packets one after the other without having to disable and re-enable the RADIO between packets, this is illustrated *Figure 29*.

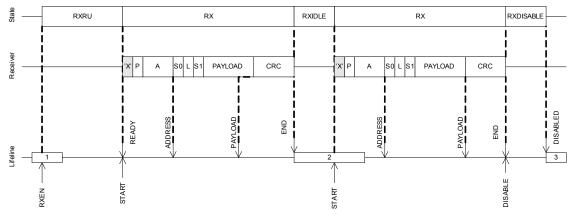


Figure 29 Reception of multiple packets

#### 16.1.12 Interframe spacing

Interframe spacing is the time interval between two consecutive packets. It is defined as the time, in micro seconds, from the end of the last bit of the previous packet received and to the start of the first bit of the subsequent packet that is transmitted. The RADIO is able to enforce this interval as specified in the TIFS register as long as TIFS is not specified to be shorter than the RADIO's startup time (see the radio timing parameters t<sub>TXEN</sub> and t<sub>RXEN</sub> in the product specification for details), that is, the time needed to switch off the receiver, and switch back on the transmitter. TIFS is only enforced if END\_DISABLE and DISABLED\_TXEN shortcuts are enabled. TIFS is only qualified for use in BLE\_1MBIT mode,

#### 16.1.13 Device address match

The device address match feature is tailored for address white listing in a *Bluetooth* low energy and similar implementations. This feature enables on-the-fly device address matching while receiving a packet on air. This feature only works in receive mode and as long as RADIO is configured for little endian.



The Device Address match unit assumes that the 48 first bits of the payload is the device address and that bit number 6 in S0 is the TxAdd bit. See the *Bluetooth* specification 4.0 for more information about device addresses, TxAdd, and whitelisting.

The RADIO is able to listen for 8 different device addresses at the same time. These addresses are specified in a DAB/DAP register pair, one pair per address, in addition to a TxAdd bit configured in the DACNF register. The DAB register specifies the 32 least significant bits of the device address, while the DAP register specifies the 16 most significant bits of the device address.

Each of the device addresses can be individually included or excluded from the matching mechanism. This is configured in the DACNF register.

#### 16.1.14 Bit counter

The RADIO implements a simple counter that can be configured to generate an event after a specific number of bits have been transmitted or received. By using shortcuts, this counter can be started from different events generated by the RADIO and then count relative to these events.

The bit counter can only be started after the RADIO has received the ADDRESS event.

The bit counter starts counting after it is started with the BCSTART task, and stopped by triggering the BCSTOP task. A BCMATCH event will be generated when the bit counter has counted the number of bits specified in the BCC register. The bit counter will continue to count bits until the DISABLED event is generated or until the BCSTOP task is triggered. The CPU can therefore, after a BCMATCH event, reconfigure the BCC value for new BCMATCH events within the same packet.

The bit counter will stop and reset on BCSTOP, STOP, and DISABLE tasks. The bit counter is also stopped and reset on END event unless the END\_START shortcut is enabled.

*Figure 30* illustrates how the bit counter can be used to generate a BCMATCH event in the beginning of the packet payload (BCC=12 bits), and generate a second BCMATCH event after sending 2 bytes (BCC=12 + 16 bits) of the payload.

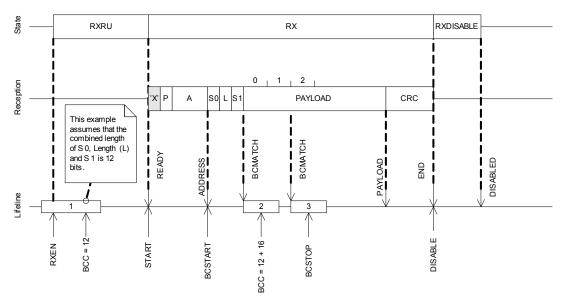


Figure 30 Bit counter example.





#### 16.1.15 Override registers

The OVERRIDE[n] (n=0...4) registers must be configured for a specific radio mode if override parameters have been written to the FICR during production calibration. It is not possible to partially override these parameters, meaning all five registers must either be used or left untouched.

#### 16.1.15.1 Using the Override registers

Before enabling the radio, check if the FICR contains override parameters for the selected radio mode by reading the FICR.OVERRIDDEN register. If the FICR contains override parameters for this mode, they have to be copied into the RADIO.OVERRIDE[n] (n=0...4) registers before the radio is enabled. See *Chapter 6 "Factory Information Configuration Registers (FICR)"* on page 18 for further details.

Override can be disabled by setting the value of OVERRIDE[4].OREN to 0. Override can be re-enabled by setting the value of OVERRIDE[4].OREN to 1.

Configuration of the OVERRIDE registers, including OVERRIDE[4].OREN, must occur before the radio is enabled using the RXEN or TXEN task. The registers must not be overwritten while the radio is enabled, that is, from RADIO.RXEN or RADIO.TXEN until the RADIO.DISABLED event.



# 16.2 Register

Register	Offset	Description
TASKS		
TXEN	0x000	Enable radio in TX mode.
RXEN	0x004	Enable radio in RX mode.
START	0x008	Start radio.
STOP	0x00C	Stop radio.
DISABLE	0x010	Disable radio.
RSSISTART	0x014	Task for starting the RSSI and take one single sample of the receive signal strength.
RSSISTOP	0x018	Task for stopping the RSSI measurement.
BCSTART	0x01C	Start bit counter.
BCSTOP	0x020	Stop bit counter.
EVENTS		
READY	0x100	Ready event.
ADDRESS	0x104	Address event.
PAYLOAD	0x108	Payload event.
END	0x10C	End event.
DISABLED	0x110	Disabled event.
DEVMATCH	0x114	A device address match occurred on the last received packet.
DEVMISS	0x118	No device address match occurred on the last received packet.
RSSIEND	0x11C	Sampling of receive signal strength complete. A new RSSI sample is ready for readout from the RSSISAMPLE register.
BCMATCH	0x128	Bit counter reached bit count value specified in BCC.
REGISTERS		
SHORTS	0x200	Shortcuts for the radio.
INTENSET	0x304	Interrupt enable set register.
INTENCLR	0x308	Interrupt enable clear register.
CRCSTATUS	0x400	CRC status.
RXMATCH	0x408	Received address.
RXCRC	0x40C	Received CRC.
DAI	0x410	Device address match index.
PACKETPTR	0x504	Packet pointer.
FREQUENCY	0x508	Frequency.
TXPOWER	0x50C	Output power.
MODE	0x510	Data rate and modulation.
PCNF0	0x514	Packet configuration 0.
PCNF1	0x518	Packet configuration 1.
BASE0	0x51C	Base address 0.
BASE1	0x520	Base address 1.
PREFIXO	0x524	Prefixes bytes for logical addresses 0-3.



Register	Offset	Description
PREFIX1	0x528	Prefixes bytes for logical addresses 4-7.
TXADDRESS	0x52C	Transmit address select.
RXADDRESSES	0x530	Receive address select.
CRCCNF	0x534	CRC configuration.
CRCPOLY	0x538	CRC polynomial.
CRCINIT	0x53C	CRC initial value.
TEST	0x540	Test features enable register.
TIFS	0x544	Interframe Spacing in µs.
RSSISAMPLE	0x548	RSSI sample.
STATE	0x550	Current radio state.
DATAWHITEIV	0x554	Data whitening initial value.
BCC	0x560	Bit counter compare.
DAB[0]	0x600	Device address 0 base segment.
DAB[1]	0x604	Device address 1 base segment.
DAB[7]	0x61C	Device address 7 base segment.
DAP[0]	0x620	Device address 0 prefix.
DAP[1]	0x624	Device address 1 prefix.
DAP[7]	0x63C	Device address 7 prefix.
DACNF	0x640	Device address match configuration.
	••	
OVERRIDE[0]	0x724	Override0 - Radio configuration parameters.
OVERRIDE[1]	0x728	Override1 - Radio configuration parameters.
OVERRIDE[2]	0x72C	Override2 - Radio configuration parameters.
OVERRIDE[3]	0x730	Override3 - Radio configuration parameters.
OVERRIDE[4]	0x734	Override4 - Radio configuration parameters.
POWER	0xFFC	Peripheral power control.

Table 19 Register overview



#### 16.2.1 SHORTS

Bit number	31 30 29 28	27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		I - GFE DCBA
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field Val	<sup>ie</sup> Value	Description
A RW READY_START		Enable or disable shortcut between READY event and START task.
B RW END_DISABLE		Enable or disable shortcut between END event and DISABLE task.
C RW DISABLED_TXEN		Enable or disable shortcut between DISABLED event and TXEN task.
D RW DISABLED_RXEN		Enable or disable shortcut between DISABLED event and RXEN task.
E RW ADDRESS_RSSISTART		Enable or disable shortcut between ADDRESS event and RSSISTART task.
F RW END_START		Enable or disable shortcut between END event and START task.
G RW ADDRESS_BCSTART		Enable or disable shortcut between ADDRESS event and BCSTART task.
I RW DISABLED_RSSISTOP		Enable or disable shortcut between DISABLED event and RSSISTOP task.

## 16.2.2 CRCSTATUS

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6	5	43	2	1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		• •			-	-	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0	0 0	) 0	0	0
ID RW Field Value ID	Va	lue			De	scri	ipti	on																							
A R					CI	RC s	tatu	is o	f pa	cke	t red	ceive	ed																		
	0				Pa	acke	et re	ceiv	/ed	witl	n CF	RC e	ror																		
	1				Pa	acke	et re	ceiv	/ed	witl	n CF	RC O	К																		

## 16.2.3 RXMATCH

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (	,
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A /	
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	,
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
A R						Lc	gic	al a	ddre	ess (	on v	vhic	h p	revi	ous	pac	ket	wa	s ree	eiv	ed.												

## 16.2.4 RXCRC

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (	
ID (Field ID)		-	-	-	-	-	-	-	-	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	,
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
A R						CF	RC f	ield	of p	orev	iou	sly r	ece	ivec	l pa	cke	t.																



## 16.2.5 PACKETPTR

Bit n	umb	er		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Fi	ield I	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A
Rese	t val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID RI	W	Field	Value ID	Va	lue			De	scri	ptio	on																							
A R	W							tra	ansr	nitt	ing	, the	e pa	cket	t po	inte	ed to	o by	/ thi	s ac	ldre	ss v	vill I	ecej be ti nis a	rans	smi	tte		nd					
								Th	nis a	ddr	ess	is a	byt	e ali	igne	ed R	AM	ade	dres	is.														
								De	ecisi	ion	poi	nt: S	TAF	RT ta	ask.																			

## 16.2.6 TXPOWER

Bit number		31 30 29 3	8 27 2	6 25	24 2	3 22	2 21	20	19	18	17	16	15	14	13	12	11	10	9 8	37	6	5	4	3	21	10
ID (Field ID)				-		-	-	-	-	-	-	-	-	-	-	-	-	-		A	A	A	A	A	A A	A A
Reset value		0 0 0	0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	) (	0	0	0	0	0 0	0 0
ID RW Field	Value ID	Value	Desc	riptio	on																					
A RW			Rad	io out	put p	owe	r. De	ecisi	on j	poir	it: T	XEN	l tas	sk.												
	POS_4_DBM	0x04	+ 4	dBm																						
	0_DBM	0	0 dE	ßm																						
	NEG_4_DBM	0xFC	-4 d	Bm																						
	NEG_8_DBM	0xF8	-8 d	Bm																						
	NEG_12_DBM	0xF4	-12	dBm																						
	NEG_16_DBM	0xF0	-16	dBm																						
	NEG_20_DBM	0xEC	-20	dBm																						
	NEG_30_DBM	0xD8	-30	dBm																						

### 16.2.7 MODE

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8 7	76	5	4	3	2	1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-	-	-	-	A	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	0	0	0	0	0	0
ID RW Field Value ID	Va	lue			De	scri	pti	on																							
A RW					Tł or	ne ra n set	dic	o suj g ar	ate a ppo e co <i>etoc</i>	rts F mp	- req atib	uen ole w	cy- /ith	shift eith	: Ke	ying									•		ng				
NRF_1MBIT	0				1	Mbi	t/s l	Nor	dic p	orop	oriet	tary	rad	io n	nod	e															
NRF_2MBIT	1				2	Mbi	t/s l	Nor	dic p	orop	oriet	tary	rad	io n	nod	le															
NRF_250_KBIT	2				25	50 kl	oit/	s No	ordio	: pr	opri	etar	y ra	ndio	mc	ode															
BLE_1MBIT	3				1	Mbi	t/s l	Blue	too	th L	ow l	Enei	rgy.																		



## 16.2.8 PCNF0

Bit ı	numb	per		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (I	ield	ID)		-	-	-	-	-	-	-	-	-	-	-	-	c	c	c	c	-	-	-	-	-	-	-	B	-	-	-	-	A	A	A A
Res	et va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID F	w	Field	Value ID	Va	lue			De	scri	ipti	on																							
А	RW	LFLEN		[0	8]			Le	engt	th o	f ler	ngtł	n fie	ld ir	n nu	mb	er o	fbi	ts. D	Decis	sion	ро	int:	STA	RT 1	tasł	٢.							
В	RW	SOLEN		[0	1]			Le	engt	th o	f SO	fiel	d in	nu	mbe	er of	f byt	tes.	Dec	isio	n po	oint	: ST	ART	tas	k.								
С	RW	S1LEN		[0	8]			Le	engt	th o	f S1	fie	ld in	nu	mbe	er of	bit	s. D	ecis	ion	poiı	nt: S	STA	RT t	ask.									



## 16.2.9 PCNF1

Bit nu	mber			31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5 4	43	3 2	1	0
ID (Fie	ld ID)			-	-	-	-	-	-	E	D	-	-	-	-	-	С	с	с	В	В	В	В	В	В	В	В	A	A	A	A A	A A	A	A
Reset	value			0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	0	0
ID RW	/ Fie	ld	Value ID	Val	lue			De	scri	ptio	on																							
A RV	V M	AXLEN		[0.	25	5]								•	ket chorly	• •				•	·					is la	rg	er tl	nar	n				
B RV	V ST	ATLEN		[0.	25	5]		th se pa of	an v nd/i cke N b	vha rece t. If yte	t is eive you s to	defi the uwa the	ineo e nu ant f e pa	d in mb		Len f by	gth tes	fie def	ld of inec	f the d by	e pa the	icke e Lei	et. U ngt	sua h fie	lly t eld i	he n tł	RA ne l	DIC RAE	) w DIO	ill )				
C RV	V BA	LEN		[2.	4]			ba to	se a tal a	add add	ress ress	an of i	d th 3 by													•								
D RV	V EN	IDIAN												f pa RT ta	cke <sup>.</sup> ask.	t ler	ngth	n fie	ld.															
			LITTLE	0				Le	ast s	sigr	nific	ant	bit	ona	air fi	rst																		
			BIG	1				M	ost s	sigr	hifica	ant	bit	on a	air fi	rst																		
E RV	v w	HITEEN						Ра	cket	t wl	nite	nin	g er	nabl	ed																			
				0				Di	sabl	led																								
				1				En	able	ed																								

#### 16.2.10 BASE0

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	5 5	5 4	3	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	4 <i>4</i>	A A	A	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	<b>b</b> (	) (	) ()	0	0	0
ID RW Field Va	alue ID	Val	ue			De	scri	ptie	on																							
A RW						Ra	ndio	bas	se a	ddre	ess	0. D	ecis	ion	poiı	nt: S	TAF	RT ta	isk.													

## 16.2.11 BASE1

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	10
ID (Field ID)	А	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field	Value ID Va	alue			De	scri	ptio	on																							
A RW					Ra	ndio	bas	se a	ddre	ess í	1. D	ecis	ion	poi	nt: S	STAF	RT ta	ask.													



## 16.2.12 PREFIX0

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	43	2	1	0
ID (Field	ID)		D	D	D	D	D	D	D	D	c	c	c	c	c	c	c	c	В	в	В	В	В	В	В	В	A	A	A	AA	A 4	A	A
Reset va	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0
ID RW	Field	Value ID	Va	lue			De	scri	ptic	on																							
A RW	AP0						Ac	ddre	ess p	orefi	x 0.	Dec	isio	n po	oint	: ST/	٩RT	tasł	κ.														
B RW	AP1																																
C RW	AP2																																
D RW	AP3																																

### 16.2.13 PREFIX1

Bit number	31 30	29 28	27 2	6 25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	5 !	5	43	2	1	0
ID (Field ID)	D D	D D	DC	D	D	с	c	с	c	с	c	c	С	В	в	в	В	В	в	B	B	A	A	A	A A	A	A	A
Reset value	0 0	0 0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	0	0 0	0	0	0
ID RW Field	Value ID Value		Desc	ripti	on																							
A RW AP4			Add	lress	prefi	x 4.	Dec	isio	n p	oint	: ST/	ART	tas	k.														
B RW AP5																												
C RW AP6																												
D RW AP7																												

### 16.2.14 TXADDRESS

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A RW							ogic sk.	al a	ddre	ess t	o be	eus	ed w	/her	n tra	insm	nitti	ng a	a pa	cke	t. D	ecis	ion	poi	nt:	ST	AR	Т				



## 16.2.15 RXADDRESSES

Bit nur	nber			31	30	) 29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	2 1	0
ID (Fiel	ld ID)			-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	н	G	F	E	D	: в	A
Reset v	/alue			0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) 0	0
ID RW	Fie	ld	Value ID	Va	lue	•		De	scri	ptio	on																							
A RW	/ A[	DR0						Er	abl	e re	cep	tion	on	log	ical	adc	lress	s 0.	Deci	isio	n po	oint	STA	RT t	ask									
				0				Di	sab	le																								
				1				Er	abl	e																								
B RW	/ A[	DR1																																
C RW	/ A[	DR2																																
D RW	/ A[	DR3																																
E RW	/ A[	DR4																																
F RW	/ A[	DR5																																
G RW	/ A[	DR6																																
H RW	/ A[	DR7																																

### 16.2.16 CRCCNF

Bit	numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8 3	76	5 5	4	3	2	1	0
ID	(Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- 1	B-		-	-	-	-	A	A
Re	set va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	0	0	0	0	0	0
ID	RW	Field	Value ID	Va	lue			De	scri	ptio	on																							
А	RW	LEN		[1	3]			CF	RC le	engt	th ir	n nu	mb	er o	fbyt	tes.	Dec	isio	n po	oint:	STA	RT	task											
				0				CF	RC le	engt	th is	zer	o, a	nd C	RC	calc	ulat	ion	is d	isab	led.													
В	RW	SKIP_ADR						Le	ave	pad	cket	ado	dres	s fie	ld o	ut o	f CF	RC c	alcu	latio	on. [	Deci	sior	n po	int:	STA	RT	tas	k.					
				0				CF	RC ca	alcu	ılati	on i	nclu	Ides	ado	dres	s fie	ld.																
				1								on c afte					e ad	ldre	ss fi	eld.	The	CR	C ca	lcul	atio	n w	rill s	star	t af					

### 16.2.17 CRCPOLY

Bit nun	nber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	B 7	6	5	4	3	2	1	0
ID (Fiel	d ID)		-	-	-	-	-	-	-	-	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	۱ A	A	A	A	A	A	-
Reset v	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	) 0	0	0	0	0	0	0
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																							
A RW							Ea co 1. Th X <sup>8</sup>	orres	tern spor	n in nds wing x <sup>3</sup> +	the to t g ex + x <sup>2</sup>	CR( he t am + 1	erm ple i = 1	ís fo 100	(por r an	nent 8-b	t. Tł	ippe ne le RC p	ast	sign	nifica	anti											



## 16.2.18 CRCINIT

Bit number	31 30 29 28	3 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1	0
ID (Field ID)		A A A A A A A A A A A A A A A A	A
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	0
ID RW Field	Value ID Value	Description	
A RW		Initial value for CRC calculation. Decision point: START task.	

## 16.2.19 FREQUENCY

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 :	3 3	21	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A A	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	) ()
ID RW Field	Value ID	Va	lue			Des	scri	ptic	n																							
A RW											•		v off: RXEI		n M	Hz:	RF f	requ	uen	cy =	240	)0 +	A (	MH	łz)							

## 16.2.20 TEST

Bit	numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 3	21	0
ID (	Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		В	A
Res	et va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	0
ID	RW	Field	Value ID	Va	lue			De	scri	ptio	on																							
А	RW	CONST_CARRIER						Co	nst	ant	car	rier.	Dec	isic	on po	oint	: TX	EN 1	tasł	κ.														
				0				Di	sab	le																								
				1				En	abl	e																								
В	RW	PLL_LOCK						PL	L lo	ock.	Dec	isio	n po	oint	: TXI	EN d	or R	XEN	tas	ik.														
				0				Di	sab	le																								
				1				En	abl	e																								

### 16.2.21 RSSISAMPLE

Bit numb	er		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 :	32	2 1	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A A	A	A
Reset val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0
ID RW	Field	Value ID	Va	lue			De	escri	ipti	on																							
A RW			[0	12	7]		ac st	tua reng	l reo gth	ceiv is as	ed s s fol	sign Iow	al st s:		gth	is a	neg							tive ceiv					th	e			



## 16.2.22 STATE

Bit nı	umber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8 7	76	5	4 3	32	1	0
ID (Fi	eld ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-	- 1	A A	A	A
Reset	value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0 0	0	0 (	0 0	0	0
ID	RW Field	Value ID	Valu	ue			De	scri	ptie	on																						
А	R						Cι	urrei	nt ra	adic	o sta	te.																				
		DISABLED	0				Ra	ndio	is i	n th	e D	ISAB	LEC	) sta	te.																	
		RXRU	1				Ra	ndio	is iı	n th	e R)	(RU	stat	te.																		
		RXIDLE	2				Ra	ndio	is iı	n th	e R)	KIDL	E st	ate.																		
		RX	3				Ra	ndio	is iı	n th	e R)	< sta	te.																			
		RXDISABLE	4				Ra	ndio	is iı	n th	e R)	(DIS	ABL	E st	ate	•																
		TXRU	9				Ra	ndio	is i	n th	e T)	(RU	stat	e.																		
		TXIDLE	10				Ra	adio	is iı	n th	e T)	(IDL	E st	ate.																		
		ТХ	11				Ra	ndio	is i	n th	e T)	< sta	te.																			
		TXDISABLE	12				Ra	adio	is iı	n th	e TX	(DIS	ABL	E st	ate																	

#### 16.2.23 DATAWHITEIV

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 ·	10	
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	1	A	A	A	A	A A	
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0 0	0 0	
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
A RW										-				. Bit nt: T			•			Posi	tior	n 6 c	of th	ne L	_SF	R, E	Bit	1 to	C				

## 16.2.24 DAI

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 '	1 (	
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A /	A /	L.
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0	0 0	0 (	
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
AR						In	dex	(n)	of d	evic	e ad	ddre	ess, s	see	DAE	8[n]	anc	d DA	P[n	], th	at g	got a	an a	dd	res	s m	ate	ch.					

#### 16.2.25 TIFS

Bit number		31	30	29	28	27	26	25	24	23	22	2 21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6 !	<b>5</b> (	4	3	2 1	1 0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A i	A	A	A /	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1 (	0 (	0 (	0 (	0 0	0 0
ID RW Field	Value ID	Va	lue			De	escr	ipti	on																							
A RW						ln de pa	iterf efin acke	fram ed a et to	as th o the	pace le ti e sta	e is t me, art c	the in n of th	nicro	e int o se st b	con	ds, f	fron	n the	e en	d of	the	e las	•					iou	IS			



## 16.2.26 DAB[n] (n=0..7)

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 '	1 (	)
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α.	A	A	A	A /	A /	
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0	0 (	0 0	0 (	)
ID RW Field	Value ID	Va	lue			De	scri	pti	on																								
A RW						D	evic	e ac	ldre	ess k	base	seg	gme	nt.																			

## 16.2.27 DAP[n] (n=0..7)

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6	5 4	4 3	3 2	2 1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1 (	0 (	0 (	0 0	) 0	0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A RW						D	evic	e ad	ddre	ss p	orefi	x.																				



#### 16.2.28 DACNF

Bit	numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	б	5 4	43	2	1	0
ID (	Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	Ρ	0	Ν	м	L	К	J	I	H	G	FI	E C	) C	В	A
Res	et va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	1	0 (	0	0	0	0
ID	RW	Field	Value ID	Va	lue			De	scri	pti	on																							
A	RW	ENA0		0 1				D	nabl isab nabl	le	r dis	able	e de	vice	e ad	dres	s m	atc	hing	g us	ing	dev	ice	add	lress	. 0.								
В	RW	ENA1						••																										
С	RW	ENA2																																
D	RW	ENA3						••																										
Е	RW	ENA4						••																										
F	RW	ENA5						••																										
G	RW	ENA6						••																										
н	RW	ENA7						••																										
L	RW	TXADD0						Ъ	Add	d fo	r de	vice	ad	dres	s 0.																			
J	RW	TXADD1						••																										
К	RW	TXADD2						••																										
L	RW	TXADD3						••																										
М	RW	TXADD4						••																										
Ν	RW	TXADD5																																
0	RW	TXADD6																																
Ρ	RW	TXADD7																																

#### 16.2.29 BCC

Bit number	31 30	29 28	27	26 2	5 24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)	A A	A A	A	A A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α.	A	A
Reset value	0 0	0 0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field Value I	Value		Des	scrip	tion																								
A RW			Bit	cour	nter c	omp	oare	•																					

## 16.2.30 OVERRIDE[n] (n=0..3)

Bit number	3	81 3	30 2	29 :	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	21	0
ID (Field ID)	A	•	A /	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α.	A	A	A A	A
Reset value	٥	) (	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0
ID RW Field Va	alue ID V	/alı	Je			De	scri	ptic	on																							
A R						Ra	ndio	ove	erric	le																						



## 16.2.31 OVERRIDE[4]

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field	ID)		В	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID RW	Field	Value ID	Va	lue			De	escri	ipti	on																							
A RW							Ra	adio	ovo	erric	de																						
B RW	OREN		0							con			RID	F[n]	(n=	:04	1) re	egist	ers.														

### 16.2.32 POWER

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	01
ID RW Field Value I	) Va	lue			De	escri	ipti	on																							
A RW	0 1				st Pe	ate erip	by s hera	swite al is	chir pov	ng ti vere	ontro he p ed o ed o	oerip off	•	•									et	to i	ts i	nit	ial				



# 17 Timer/counter (TIMER)

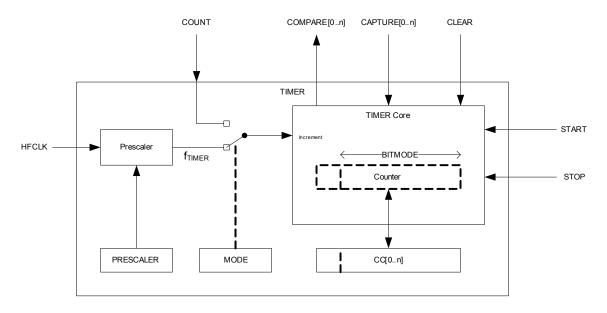


Figure 31 Block schematic for Time/counter

## 17.1 Functional description

The TIMER can operate in two modes, Timer mode and Counter mode. In both modes the TIMER is started by triggering the START task, and stopped by triggering the STOP task. The TIMER is a count-up timer.

In Timer mode, the TIMER's internal Counter register is incremented by one for every tick of the timer frequency f<sub>TIMER</sub> as illustrated in *Figure 31*. The timer frequency is derived from HFCLK as described below using the values specified in the PRESCALER register:

$$f_{TIMER} = \frac{HFCLK}{2^{PRESCALER}}$$

In counter mode, the TIMER's internal Counter register is incremented by one each time the COUNT task is triggered, that is, the timer frequency and the prescaler are not utilized in counter mode. Similarly, the COUNT task has no effect in Timer mode.

The TIMER's maximum value is configured by changing the bit-width of the timer in the BITMODE register. For details on which bit-widths that are supported on the different timers see the device product specification.

The PRESCALER register and the BITMODE register must only be updated when the timer is stopped. If these registers are updated while the TIMER is started then this may result in unpredictable behavior.

When the timer is incremented beyond its maximum value the Counter register will overflow and the TIMER will automatically start over from zero.

The Counter register can be cleared, that is, its internal value set to zero explicitly, by triggering the CLEAR task.



The TIMER implements multiple capture/compare registers, see the product specification for more information on how many capture/compare registers that are supported in the chip.

#### 17.1.1 Compare

The TIMER implements one COMPARE event for every available capture/compare register. A COMPARE event is generated when the Counter is incremented and then becomes equal to the value specified in one of the capture compare registers. When the Counter value becomes equal to the value specified in a capture compare register CC[n], the corresponding compare event COMPARE[n] is generated.

BITMODE specifies how many bits of the Counter register and the capture/compare register that are used when the comparison is performed. Other bits will be ignored.

#### 17.1.2 Capture

The TIMER implements one capture task for every available capture/compare register. Every time the CAPTURE[n] task is triggered the Counter value is copied to the CC[n] register.

#### 17.1.3 Task priority

If the START task and the STOP task are triggered at the same time, that is, within the same period of HFCLK, the STOP task will be prioritized.

#### 17.1.4 Task delays

The CLEAR task, COUNT task and the STOP task will guarantee to take effect within one clock cycle of the HFCLK. Depending on sub-power mode, the START task may require longer time to take effect, see product specification for more information. See POWER chapter, *Chapter 11 "Power management (POWER)"* on page 40, for more information about sub-power modes.



# 17.2 Registers

Register	Offset	Description
TASKS		
START	0x000	Start Timer
STOP	0x004	Stop Timer
COUNT	0x008	Increment Timer (Counter mode only)
CLEAR	0x00C	Clear timer
CAPTURE[0]	0x040	Capture Timer value to CC0 register
CAPTURE[1]	0x044	Capture Timer value to CC1 register
CAPTURE[2]	0x048	Capture Timer value to CC2 register
CAPTURE[3]	0x04C	Capture Timer value to CC3 register
EVENTS		
COMPARE[0]	0x140	Compare event on CC[0] match
COMPARE[1]	0x144	Compare event on CC[1] match
COMPARE[2]	0x148	Compare event on CC[2] match
COMPARE[3]	0x14C	Compare event on CC[3] match
REGISTERS		
SHORTS	0x200	Shortcuts
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
MODE	0x504	Timer mode selection
BITMODE	0x508	Configure the number of bits used by the TIMER
PRESCALER	0x510	Timer prescaler register
CC[0]	0x540	Capture/Compare register 0
CC[1]	0x544	Capture/Compare register 1
CC[2]	0x548	Capture/Compare register 2
CC[3]	0x54C	Capture/Compare register 3

Table 20 Register overview



## 17.2.1 SHORTS

Bit	num	ber	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (
ID	(Field	ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	н	G	F	E	-	-	-	-	D	c	B /
Re	set va	lue	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID		Field	Va	lue			De	scri	ptic	on																							
Α	RW	COMPARE0_CLEAR					En	nabl	e or	disa	able	e sho	orto	ut k	oetw	/eer	n CC	DMP	ARE	0 ar	nd C	CLEA	۹R ta	ask.									
В	RW	COMPARE1_CLEAR					En	nabl	e or	disa	able	e sho	orto	ut k	oetw	/eei	n CC	DMP	ARE	1 ar	nd (	CLEA	AR ta	ask.	•								
С	RW	COMPARE2_CLEAR					Er	nabl	e or	disa	able	e sho	orto	ut k	oetw	/eei	۱CC	DMP	ARE	2 ar	nd (	CLEA	AR ta	ask.									
D	RW	COMPARE3_CLEAR					En	nabl	e or	disa	able	e sho	orto	ut k	oetw	/eei	n CC	DMP	ARE	3 ar	nd (	CLEA	۹R ta	ask.	•								
Е	RW	COMPARE0_STOP					En	nabl	e or	disa	able	e sho	orto	ut k	oetw	/eer	n CC	OMP	ARE	0 ar	nd S	TO	<sup>o</sup> ta	sk.									
F	RW	COMPARE 1_STOP					En	nabl	e or	disa	able	e sho	orto	ut k	oetw	/eer	n CC	OMP	ARE	1 ar	nd S	TO	<sup>o</sup> ta	sk.									
G	RW	COMPARE2_STOP					Er	nabl	e or	disa	able	e sho	orto	ut k	oetw	/eer	n CC	OMP	ARE	2 ar	nd S	TO	P ta	sk.									
Н	RW	COMPARE3_STOP					En	nabl	e or	disa	able	e sho	orto	ut b	oetw	/eer	n CC	OMP	ARE	3 ar	nd S	TO	<sup>o</sup> ta	sk.									

#### 17.2.2 MODE

Bit numl	ber	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A
Reset va	lue	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW	Field	Val	ue			De	scri	ptic	on																								
A RW						Tii	mer	mo	de																								
	TIMER	0				Se	lect	tim	ier r	nod	e																						
	COUNTER	1				Se	lect	cou	unte	r m	ode																						

## 17.2.3 PRESCALER

Bit number	31 30 29 28	27 26 25 24 23 22	22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)			A A A A
Reset value	0 0 0 0	0 0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field	Value	Description	
A RW	[09]	Prescaler value	

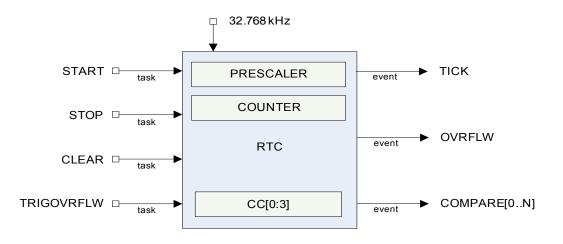


## 17.2.4 BITMODE

Bit number	31	30	29	9 2	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)	-	-	-	-		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A
Reset value	0	0	0	C	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field	Va	lue				De	scri	pti	on																								
A RW						Tir	ner	bit	wic	lth																							
16BIT	0					16	bit	tim	ner b	oit v	vidt	h																					
08BIT	1					8 k	oit t	ime	er bi	t wi	dth																						
24BIT	2					24	bit	tim	ner l	oit v	vidt	h																					
32BIT	3					32	bit	tim	ner k	oit v	vidt	h																					



## 18 Real Time Counter (RTC)



*Figure 32 RTC block schematic* 

## 18.1 Functional description

The RTC is a 24 bit low-frequency clock with frequency prescaling and tick, compare, and overflow events.

#### 18.1.1 Clock Source

The RTC will run off a low-frequency clock (LFCLK) running at 32.768 kHz. This clock may be either a RC oscillator or a crystal oscillator. The COUNTER resolution will therefore be 30.517 µs. The RTC must be able to run while the 16 MHz system clock (SysClk) source is OFF.

#### 18.1.2 Resolution versus overflow and the PRESCALER

COUNTER increment frequency:

 $f_{RTC} = \frac{32,768kHz}{PRESCALER + 1}$ 

The PRESCALER register is read/write when the RTC is stopped. The PRESCALER register is read-only once the RTC is STARTed. Writing to the PRESCALER register when the RTC is started has no effect.

The PRESCALER is restarted on START, CLEAR and TRIGOVRFLW, that is, the prescaler value is latched to an internal register (<<PRESC>>) on these tasks.



#### Examples

- Desired COUNTER frequency 100 Hz (10 ms per counter period) PRESCALER = round(32.768 kHz / 100 Hz) - 1 = 327 f<sub>RTC</sub> = 99.9 Hz 10009.576 μs counter period
   Desired COUNTER frequency 8 Hz (125 millisecond counter period)
- Desired COUNTER frequency 8 Hz (125 millisecond counter period) PRESCALER = round (32.768 kHz / 8 Hz) – 1 = 4095 f<sub>RTC</sub> = 8 Hz 125 ms counter period

Prescaler	<b>Counter resolution</b>	Overflow
0	30.517 μs	512 seconds
2 <sup>8</sup> - 1	7812.5 µs	131072 seconds
2 <sup>12</sup> - 1	125 ms	582.542 hours

Table 21 RTC Resolution versus Overflow

#### 18.1.3 The COUNTER register

The COUNTER increments on LFCLK when the internal PRESCALER register (<<PRESC>>) is 0x00. The internal <<PRESC>> register is reloaded from the PRESCALER register. If enabled, the TICK event occurs on each increment of the COUNTER. The TICK event can be disabled.

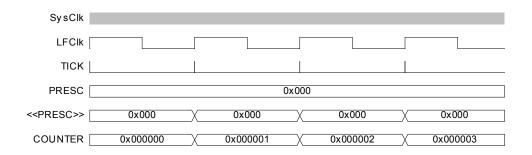
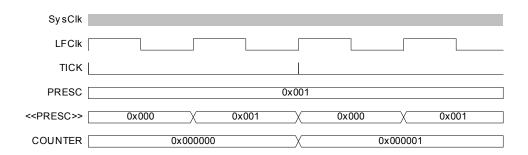


Figure 33 Timing diagram - COUNTER\_PRESCALER\_0



*Figure 34 Timing diagram - COUNTER\_PRESCALER\_1* 

#### 18.1.4 Overflow features

The TRIGOVRFLW task sets the COUNTER value to 0xFFFFF0 to allow SW test of the overflow condition. OVRFLW occurs when COUNTER overflows from 0xFFFFFF to 0x000000.

**Note:** The OVRFLW event is disabled by default.

#### 18.1.5 The TICK event

The TICK event enables low power "tick-less" RTOS implementation as it optionally provides a regular interrupt source for a RTOS without the need to use the ARM<sup>®</sup> SysTick feature. Using the RTC TICK event rather than the SysTick allows the CPU to be powered down while still keeping RTOS scheduling active.

**Note:** The TICK event is disabled by default.

#### 18.1.6 Event Control feature

To optimize RTC power consumption, events in the RTC can be prevented from being routed to the PPI by configuring the EVTEN register. When routing an event to the PPI is disabled the HFCLK will no longer be required and power consumption will be reduced.

For example, if the TICK event is not required for an application, routing of this event to the PPI should be disabled as it is frequently occurring and may raise power consumption if HFCLK can otherwise be powered down for long durations.

#### 18.1.7 Compare feature

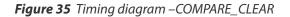
There are three supported compare registers and up to one optional. See product specification for details on available compare registers.

When setting a compare register, the following behavior of the RTC compare event should be noted:

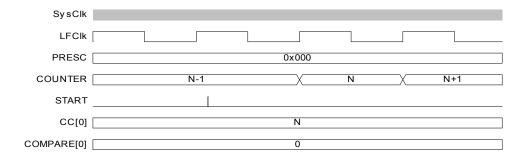
• If a CC register value is 0 when a CLEAR task is set, this will not trigger a COMPARE event.

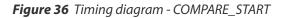


SysClk	
LFClk	
PRESC	0x000
COUNTER	X X 0x000000
CLEAR	
CC[0]	0x000000
COMPARE[0]	0

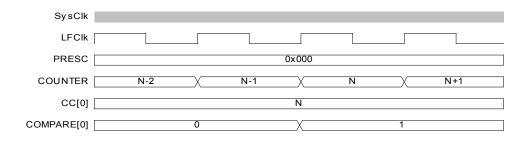


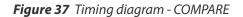
• If a CC register is N and the COUNTER value is N when the START task is set, this will **not** trigger a COMPARE event.





• COMPARE occurs when a CC register is N and the COUNTER value transitions from N-1 to N.





• If the COUNTER is N, writing N+2 to a CC register is guaranteed to trigger a COMPARE event at N+2.



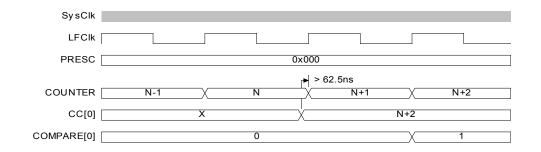
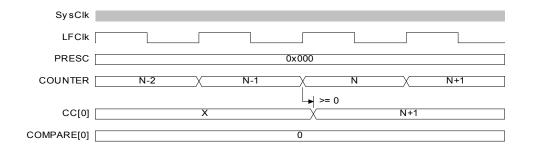
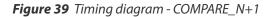


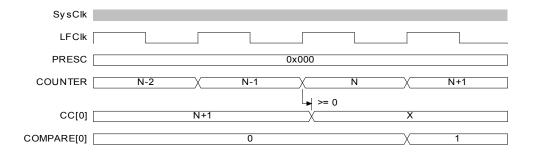
Figure 38 Timing diagram - COMPARE\_N+2

• If the COUNTER is N, writing N or N+1 to a CC register may not trigger a COMPARE event.





 If the COUNTER is N and the current CC register value is N+1 or N+2 when a new CC value is written, a match may trigger on the previous CC value before the new value takes effect. If the current CC value greater than N+2 when the new value is written, there will be no event due to the old value.



*Figure 40 Timing diagram - COMPARE\_N-1* 



## 18.1.8 TASK and EVENT jitter/delay

The source of jitter or delay in the RTC is due to the peripheral clock being a low frequency clock (LFCLK) which is not synchronous to the faster system clock (SysClk). Registers in the peripheral interface, part of the SysClk domain, have a set of mirrored registers in the LFCLK domain. For example, the COUNTER value accessible from the CPU is in the SysClk domain and is latched on read from an internal register called COUNTER in the LFCLK domain. COUNTER is the register which is actually modified each time the RTC ticks. These registers must be synchronised between clock domains (SysClk and LFCLK).

The following is a summary of the jitter introduced on tasks and events. Figures illustrating jitter follow.

Task	Delay
CLEAR, STOP, START, TRIGOVRFLW	+15 to 46 μs

 Table 22
 RTC jitter magnitudes on tasks

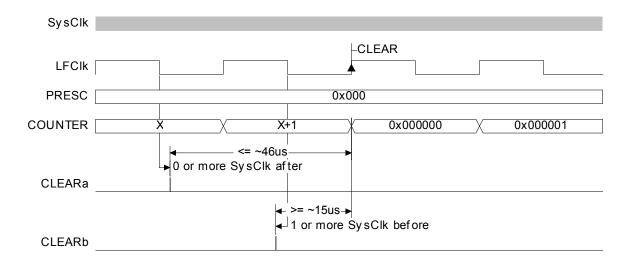
Operation/Function	Jitter
START to COUNTER increment	+/- 15 μs
COMPARE to COMPARE <sup>1</sup>	+/- 62.5 ns

1.Assumes RTC runs continuously between these events.

Note: 32.768 kHz clock jitter is additional to the above provided numbers

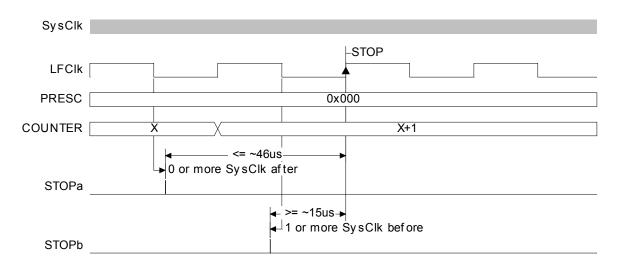
Table 23 RTC jitter magnitudes on events

1. CLEAR and STOP (and TRIGOVRFLW; not shown) will be delayed as long as it takes for the peripheral to clock a falling edge and rising of the LFCLK. This is between 15.2585  $\mu$ s and 45.7755  $\mu$ s – rounded to 15  $\mu$ s and 46  $\mu$ s for the remainder of the section.



*Figure 41 Timing diagram - DELAY\_CLEAR* 





*Figure 42 Timing diagram - DELAY\_STOP* 

2. The START task will start the RTC. The first increment of COUNTER (and instance of TICK event) will be after 30. 5 μs +/-15 μs, again because at least 1 falling edge must occur after the START TASK before the rising edge causes events and COUNTER increment. The figures show the smallest and largest delays to on the START task which appears as a +/-15 μs jitter on the first COUNTER increment.

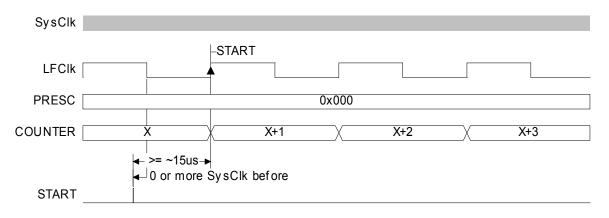
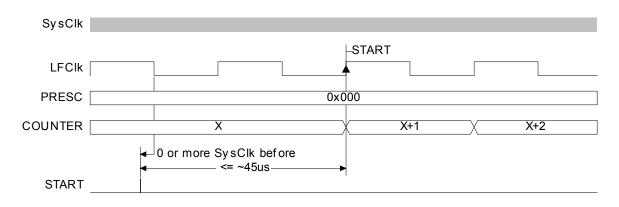


Figure 43 Timing diagram - JITTER\_START-







#### 18.1.9 Reading the COUNTER register

To read the COUNTER register, the internal <<COUNTER>> value is sampled. To ensure <<COUNTER>> is safely sampled (considering a LFCLK transition may occur during a read), the CPU and core memory bus are halted for 3 cycles by lowering the core PREADY signal. The Read takes the CPU 2 cycles in addition resulting in the COUNTER register read taking a fixed five SysClk clock cycles.

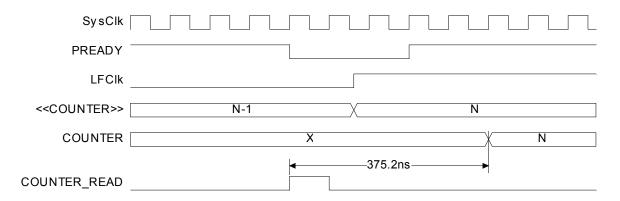
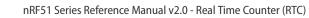


Figure 45 Timing diagram - COUNTER\_READ





# 18.2 Registers

Register	Offset	Description
TASKS		
START	0x000	Start RTC COUNTER.
STOP	0x004	Stop RTC COUNTER.
CLEAR	0x008	Clear RTC COUNTER.
TRIGOVRFLW	0x00C	Set COUNTER to 0xFFFFF0.
EVENTS		
TICK	0x100	Event on COUNTER increment.
OVRFLW	0x104	Event on COUNTER overflow.
COMPARE[0]	0x140	Compare event on CC[0] match.
COMPARE[1]	0x144	Compare event on CC[1] match.
COMPARE[2]	0x148	Compare event on CC[2] match.
COMPARE[3]	0x14C	Compare event on CC[3] match.
REGISTERS		
INTENSET	0x304	Configures which events shall generate a RTC interrupt.
INTENCLR	0x308	Configures which events shall not generate a RTC interrupt.
EVTEN	0x340	Enable or disable event routing to PPI.
EVTENSET	0x344	Enable event routing to PPI.
EVTENCLR	0x348	Disable event routing to PPI.
COUNTER	0x504	Current COUNTER value.
PRESCALER	0x508	12 bit prescaler for COUNTER frequency (32768/(PRESCALER+1)). Must be written when RTC is stopped.
CC[0]	0x540	Compare register.
CC[1]	0x544	Compare register.
CC[2]	0x548	Compare register.
CC[3]	0x54C	Compare register.



#### 18.2.1 EVTEN

Bit	numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5 4	3	2	1	0
ID	(Field	ID)		-		-	-	-	-	-	-	-	-	-	-	F	E	D	С	-	-	-	-	-	-	-	-	-	-		-	-	B	A
Re	set va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) 0	0	0	0
ID	RW	Field V	Value ID	Val	ue			De	scri	iptio	on																							
А	RW	TICK						En	abl	e or	<sup>,</sup> dis	abl	e ro	utin	g of	TIC	Ke	vent	t to	PPI.														
				1				En	abl	e																								
				0				Di	sab	le																								
В	RW	OVERFLW						En	abl	e or	' dis	abl	e ro	utin	g of	OV	RFL	We	ven	t to	PPI	•												
				1				En	abl	e																								
				0				Di	sab	le																								
C	RW	COMPARE0						En	abl	e or	' dis	abl	e ro	utin	g of	CO	MP	ARE	[0]e	ven	t to	PPI.	•											
				1				En	abl	e																								
				0				Di	sab	le																								
D	RW	COMPARE1						En	abl	e or	<sup>r</sup> dis	abl	e ro	utin	g of	CO	MP	ARE	[1]e	ven	t to	PPI.	•											
				1				En	abl	e																								
				0				Di	sab	le																								
Е	RW	COMPARE2						En	abl	e or	dis	abl	e ro	utin	g of	CO	MP	٩RE	[2]e	ven	t to	PPI.												
				1				En	abl	e																								
				0				Di	sab	le																								
F	RW	COMPARE3						En	abl	e or	<sup>,</sup> dis	abl	e ro	utin	g of	CO	MP	ARE	[3]e	ven	t to	PPI.												
				1				En	abl	e																								
				0				Di	sab	le																								



## 18.2.2 EVTENSET

Bit number	31 30 29 28	27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		F E D C B A
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field Value ID	Value	Description
A TICK		Enable routing of TICK event to PPI.
W	1	Enable
R	EVTEN.0	Readback bit 0 of EVTEN
B OVRFLW		Enable routing of OVRFLW event to PPI.
W	1	Enable
R	EVTEN.1	Readback bit 1 of EVTEN
C COMPARE0		Enable routing of COMPARE[0] event to PPI.
W	1	Enable
R	EVTEN.16	Readback bit 16 of EVTEN
D COMPARE1		Enable routing of COMPARE[1] event to PPI.
W	1	Enable
R	EVTEN.17	Readback bit 17 of EVTEN
E COMPARE2		Enable routing of COMPARE[2] event to PPI.
W	1	Enable
R	EVTEN.18	Readback bit 18 of EVTEN
F COMPARE3		Enable routing of COMPARE[3] event to PPI.
W	1	Enable
R	EVTEN.19	Readback bit 19 of EVTEN



## 18.2.3 EVTENCLR

| nun   | nber  |  | 31  | 30  | 29   | 28   
   
   
  | 27   
   
   | 26   | 25   | 24         | 23   | 22   | 21   | 20  
  | 19   | 18   | 17   | 16   | 15   | 14   | 13   | 12  
  | 11  | 10  | 9   | 8  | 7 (  | 5 5  
  | 5 4   | 3   | 2   | 1 (   | D  |
|-------|---|--|---|---|--
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--|--|--|------------|--|--|--
--|--|--|--|--|--|--|--
--	---	---	---	--	--
(Fiel	d ID)		-	-	-
   
   
  | -  
   
   | -  | -  | -          | -  | -  | -  | -   
  | F  | Е  | D  | С  | -  | -  | -  | -   
  | -   | -   | -   | -  |  | -  
  | -   | -   | -   | В   | A  |
| set v | alue  |  | 0   | 0   | 0  | 0  
   
   
  | 0  
   
   | 0  | 0  | 0          | 0  | 0  | 0  | 0   
  | 0  | 0  | 0  | 0  | 0  | 0  | 0  | 0   
  | 0   | 0   | 0   | 0  | 0 (  | ) (  
  | 0   | 0   | 0   | 0   | D  |
| RW    | Field   | Value ID   | Va  | lue   |  |  
   
   
  | De   
   
   | scri   | pti  | on         |  |  |  | |
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
|       | TICK  |  |   |   |  |  
   
   
  | Di   
   
   | sab  | le re  | outi       | ng   | of T   | ICK  | eve   
  | nt t   | o Pf   | PI.  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| W     |   |  | 1   |   |  |  
   
   
  | Di   
   
   | sab  | le   |            |  |  |  | |
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| R     |   |  | E١  | VTEN  | 1.0  |  
   
   
  | Re   
   
   | adł  | bacl   | < bit      | t 0 c  | of E\  | /TEI   | N   
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
|       | OVRFLW  |  |   |   |  |  
   
   
  | Di   
   
   | sab  | le re  | outi       | ng   | of C   | VRF  | LW  
  | eve  | ent 1  | to P   | PI.  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| W     |   |  | 1   |   |  |  
   
   
  | Di   
   
   | sab  | le   |            |  |  |  | |
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| R     |   |  | E١  | VTEN  | <b>I</b> .1  |  
   
   
  | Re   
   
   | adł  | bacl   | < bit      | t 1 c  | of E\  | /TEI   | N   
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
|       | COMPARE   | 0  |   |   |  |  
   
   
  | Di   
   
   | sab  | le re  | outi       | ng   | of C   | ОМ   | PAR   
  | RE[0]  | ev   | ent  | to F   | PI.  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| W     |   |  | 1   |   |  |  
   
   
  | Di   
   
   | sab  | le   |            |  |  |  | |
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| R     |   |  | E\  | VTEN  | I.16   |  
   
   
  | Re   
   
   | adł  | bacl   | < bit      | t 16   | of E   | EVT  | ΞN  
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
|       | COMPARE   | 1  |   |   |  |  
   
   
  | Di   
   
   | sab  | le re  | outi       | ng   | of C   | OM   | PAR   
  | RE[1]  | ev   | ent  | to F   | PI.  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| W     |   |  | 1   |   |  |  
   
   
  | Di   
   
   | sab  | le   |            |  |  |  | |
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| R     |   |  | E١  | VTEN  | l.17   |  
   
   
  | Re   
   
   | adł  | bacl   | < bit      | t 17   | of   | EVT  | ΞN  
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
|       | COMPARE   | 2  |   |   |  |  
   
   
  | Di   
   
   | sab  | le re  | outi       | ng   | of C   | ОМ   | PAR   
  | RE[2]  | ev   | ent  | to F   | PI.  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| W     |   |  | 1   |   |  |  
   
   
  | Di   
   
   | sab  | le   |            |  |  |  | |
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| R     |   |  | E١  | VTEN  | I.18   |  
   
   
  | Re   
   
   | adł  | bacl   | < bit      | t 18   | of E   | VT   | ΞN  
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
|       | COMPARE   | 3  |   |   |  |  
   
   
  | Di   
   
   | sab  | le re  | outi       | ng   | of C   | ОМ   | PAR   
  | RE[3]  | ev   | ent  | to F   | PPI.   |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| W     |   |  | 1   |   |  |  
   
   
  | Di   
   
   | sab  | le   |            |  |  |  | |
  |  |  |  |  |  |  |  |   
  |   |   |   |  |  |  
  |   |   |   |   |  |
| R     |   |  | E١  | VTEN  | l.19   |  
   
   
  | Re   
   
   | adł  | bacl   | < bit      | t 19   | of E   | EVT  | EN  
  |  |  |  |  |  |  |  | | | | | | | | | | | | | | | | | | | | | | | | | | | |
  |   |   |   |  |  |  
  |   |   |   |   |  |
|       | (Fiel<br>set v<br>RW<br>W<br>R<br>W<br>R<br>W<br>R<br>R<br>W<br>R<br>W<br>R | <ul> <li>W</li> <li>R</li> <li>OVRFLW</li> <li>W</li> <li>COMPARE</li> <li>M</li> <li>COMPARE</li> <li>M</li> <li>COMPARE</li> <li>M</li> <li>COMPARE</li> <li>M</li> <li>COMPARE</li> <li>M</li> <li>COMPARE</li> </ul> | (Field ID)         relation         TICK         TICK         TICK         ID         ID | (Field ID)       -         set value       0         RW       Field       Value ID       Value ID         RW       CONFLW       11         R       COMPAREO       11         R       COMPAREO | (Field ID)       - | (Field ID)       - <td< th=""><th>(Field ID)       -       <td< th=""><th>(Field ID)       <math>  -</math></th><th>(Field ID)<math>I = I = I = I = I = I = I = I = I = I =</math></th><th>(Field ID)</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       -</th><th>relation of the interview of the interview</th><th>General and a construction of the second of the second</th><th>General and a construction of the second of the second</th><th>Grand and and any operation of the second of the second operation o</th><th>Grand and and any operation of the second of the second operation o</th><td>General Control       Price of the control Control       Price of the contro       &lt;</td><td>Yeak       1</td><th>General Problem Prob</th><th>Gried ID<th>(Field ID) 0&lt;</th><th>(Field ID) 0&lt;</th><th>Criped ID 0<!--</th--><th>Control Control Control</th><td>(Field Value 10 Val</td></th></th></td<></th></td<> | (Field ID)       - <td< th=""><th>(Field ID)       <math>  -</math></th><th>(Field ID)<math>I = I = I = I = I = I = I = I = I = I =</math></th><th>(Field ID)</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       0</th><th>(Field ID)       -</th><th>relation of the interview of the interview</th><th>General and a construction of the second of the second</th><th>General and a construction of the second of the second</th><th>Grand and and any operation of the second of the second operation o</th><th>Grand and and any operation of the second of the second operation o</th><td>General Control       Price of the control Control       Price of the contro       &lt;</td><td>Yeak       1</td><th>General Problem Prob</th><th>Gried ID<th>(Field ID) 0&lt;</th><th>(Field ID) 0&lt;</th><th>Criped ID 0<!--</th--><th>Control Control Control</th><td>(Field Value 10 Val</td></th></th></td<> | (Field ID) $  -$ | (Field ID) $I = I = I = I = I = I = I = I = I = I =$ | (Field ID) | (Field ID)       0 | (Field ID)       0 | (Field ID)       0 | (Field ID)       0 | (Field ID)       0 | (Field ID)       0 | (Field ID)       0 | (Field ID)       0 | (Field ID)       - | relation of the interview | General and a construction of the second | General and a construction of the second | Grand and and any operation of the second of the second operation o | Grand and and any operation of the second of the second operation o | General Control       Price of the control Control       Price of the contro       < | Yeak       1 | General Problem Prob | Gried ID <th>(Field ID) 0&lt;</th> <th>(Field ID) 0&lt;</th> <th>Criped ID 0<!--</th--><th>Control Control Control</th><td>(Field Value 10 Val</td></th> | (Field ID) 0< | (Field ID) 0< | Criped ID 0 </th <th>Control Control Control</th> <td>(Field Value 10 Val</td> | Control | (Field Value 10 Val |

#### 18.2.4 COUNTER

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)		-	-	-	-	-	-	-	-	A	A	A	A	A	A	A	A	A	Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset va	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW	Field \	Value ID	Va	lue			De	scri	ptio	on																								
А							Co	oun	terv	/alu	e																							

## 18.2.5 PRESCALER

Bit numb	er	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	A	A	A	A	A	A	A
Reset val	lue	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW	Field Value ID	Va	lue			De	escri	ptio	on																								
А						PI	RESC	CALI	ER v	alue	e																						



# 18.2.6 CC[N] (n=0..3)

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)	-	-	-	-	-	-	-	-	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field Value ID	) Va	lue			De	scri	ptic	on																								
A					C	omp	are	valu	Je																							



# 19 Watchdog timer (WDT)

# 19.1 Functional description

The watchdog is implemented as a down-counter that generates a TIMEOUT event when it wraps over after counting down to 0. The watchdog timer is started by triggering the START task, whereupon the watchdog counter is loaded with the value specified in the CRV register. This counter is also reloaded with the value specified in the CRV register when a reload request is granted.

The counter reload value is specified in the CRV register, and the timer is started using the START task.

The watchdog's timeout period is given by:

$$\frac{CRV+1}{32768}[s]$$

When started, the watchdog will automatically force the 32.768 kHz RC oscillator on.

#### 19.1.1 Reload criteria

The watchdog has 8 separate reload request registers which shall be used to request the watchdog to reload its counter with the value specified in the CRV register. To reload the watchdog counter, the special value 0x6E524635 needs to be written to all enabled reload registers. One or more RR registers can be individually enabled via the RREN register.

#### 19.1.2 Temporarily pausing the watchdog

By default the watchdog will be active counting down the down-counter while the CPU is sleeping and when it is halted by the debugger. It is however possible to configure the watchdog to automatically pause while the CPU is sleeping as well as when it is halted by the debugger.

#### 19.1.3 Watchdog reset

A TIMEOUT event will automatically lead to a watchdog reset equivalent to a system reset, see *Chapter 11* "*Power management (POWER)*" on page 40. If the watchdog is configured to generate an interrupt on the TIMEOUT event, the watchdog reset will be postponed with two 32.768 kHz clock cycles after the TIMEOUT event has been generated. Once the TIMEOUT event has been generated, the impending watchdog reset will always be effectuated.

The watchdog must be configured before it is started. After it is started, the watchdog's configuration registers, which comprises registers CRV, RREN, and CONFIG, will be blocked for further configuration.

The watchdog is reset when the device is put into System OFF mode. The watchdog is also reset when the whole system is reset, except for when the system is reset through a soft reset, see *Chapter 11 "Power management (POWER)"* on page 40 for more information about reset types.

When the device starts running again, after a reset, or waking up from OFF mode, the watchdog configuration registers will be available for configuration again.



# 19.2 Registers

Register	Offset	Description	
TASKS			
START	0x000	Start the watchdog	
EVENTS			
TIMEOUT	0x100	Watchdog timeout	
REGISTERS			
INTENSET	0x304	Interrupt enable set register	
INTENCLR	0x308	Interrupt enable clear register	
RUNSTATUS	0x400	Run status	
REQSTATUS	0x404	Request status	
CRV	0x504	Counter reload value	
RREN	0x508	Reload request enable	
CONFIG	0x50C	Configuration register	
RR[0]	0x600	Reload request 0	
RR[1]	0x604	Reload request 1	
••	••		
RR[7]	0x61C	Reload request 7	

#### Table 24 Register overview

## 19.2.1 RUNSTATUS

Bit number	:	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A
Reset value	(	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field	Value ID	Val	ue			De	scri	ptic	on																								
A R		0				W	dica 'atch 'atch	ndog	g is i	not	runı	ning	g is r	unn	ing.																		



## 19.2.2 REQSTATUS

Bit num	nber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	55	5 4	3	2	10
ID (Field	d ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	H (	G F	E	D	c	ΒA
Reset v	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	0	0	0	01
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																						
A R	RRO						Re	que	est s	stat	us fo	or R	R[0]	reg	iste	r																
			0				RF	R[0]	reg	iste	r is ı	not	ena	blec	l, or	are	alr	ead	y re	que	stin	g re	eloa	d								
			1				RF	R[0]	reg	iste	r is e	enal	olec	l, an	d a	re n	ot y	et r	equ	esti	ng r	elo	ad									
BR	RR1						••																									
C R	RR2																															
DR	RR3						••																									
RR	RR4						••																									
FR	RR5						••																									
G R	RR6																															
ΗR	RR7																															

#### 19.2.3 CRV

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 ·	10
ID (Field ID)	А	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	0 0
ID RW Field	Value ID Va	lue			De	scri	ptic	on																							
A RW					Co	ount	er r	eloa	id va	alue	in n	um	ber	of cy	ycle	s of	the	32.7	768	kHz	clo	ck									

#### 19.2.4 RREN

Bit	numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	B 7	6	5	4	3	2	1 0
ID (	Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- +	I G	F	E	D	c	B A
Res	et va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	0	0	0	0	01
ID	RW	Field	Value ID	Va	lue			De	scri	pti	on																						
А	RW	RR0						Er	abl	e oi	r dis	able	e RF	R[0]	regi	ster																	
				0				Di	sab	le R	R[0]	] reg	giste	er																			
				1				Er	abl	e Rl	R[0]	reg	iste	r																			
В	RW	RR1																															
C	RW	RR2																															
D	RW	RR3																															
R	RW	RR4																															
F	RW	RR5																															
G	RW	RR6																															
н	RW	RR7																															



# 19.2.5 CONFIG

Bit numb	er		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 E	5 5	4	3	21	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-	B		A
Reset val	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) ()	0	0	0 0	) 1
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																						
A RW	SLEEP							onfi eepi			e w	atcl	hdo	g to	eitl	ner	be p	baus	ed,	or l	ept	run	nin	g, w	/hil	e tł	ne C	PU	is			
		PAUSE	0				Pa	ause	e wa	tch	dog	g wł	nile t	the	CPL	l is s	slee	ping	J.													
		RUN	1				Ke	eep	the	wat	tch	dog	run	nin	g w	hile	the	CPI	J is	slee	pin	g.										
B RW	HALT								•				hdo ggei	•	eitl	ner	be p	baus	ed,	or l	ept	run	nin	g, w	/hil	e tł	ne C	PU	is			
		PAUSE	0				Pa	ause	e wa	tch	dog	g wł	nile t	the	CPL	l is l	halt	ed b	y th	e d	ebu	gge	er.									
		RUN	1				Ke	eep	the	wat	tch	dog	run	nin	g w	hile	the	CPI	J is	halt	ed l	oy tl	he d	lebu	ugg	ger.						

# 19.2.6 RR[n] (n=0..7)

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (	D
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	4
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	D
ID RW Field	Value ID	Va	lue			De	scri	ptie	on																								
A W						Re	eloa	d re	que	st re	egis	ter																					



# 20 Random Number Generator (RNG)

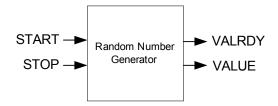


Figure 46 Random Number Generator

# 20.1 Functional description

The Random Number Generator (RNG) generates true non-deterministic random numbers based on internal thermal noise.

The RNG is started by triggering the START task. When started, new random numbers are generated continuously and written to the VALUE register when ready. A VALRDY event is generated for every new random number that is written to the VALUE register.

#### 20.1.1 Digital error correction

A digital corrector algorithm is employed on the internal bit stream to remove any bias toward '1' or '0'. The bits are then queued into an 8 bit register for parallel readout from the VALUE register.

It is possible to disable the bias in the CONFIG register. This offers a substantial speed advantage, but may result in a statistical distribution that is not perfectly uniform.

#### 20.1.2 Speed

The time needed to generate one random byte of data is unpredictable, and may vary from one byte to the next. This is especially true when digital error correction is enabled.



# 20.2 Registers

Register	Offset	Description
TASKS		
START	0x000	Task starting the random number generator.
STOP	0x004	Task stopping the random number generator.
EVENTS		
VALRDY	0x100	Event being generated for every new random number written to the VALUE register.
REGISTERS		
SHORTS	0x200	Shortcut register.
INTENSET	0x304	Interrupt enable set register.
INTENCLR	0x308	Interrupt enable clear register.
CONFIG	0x504	Configuration register.
VALUE	0x508	Output random number.

Table 25 Register overview

#### 20.2.1 SHORTS

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-	-	-	-	- A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID RW Field Value ID	Va	lue			De	scri	pti	on																							
A RW VALRDY_STOP					Sh	ort	cut	bet	wee	en V	ALR	DY	eve	nt a	nd	sto	P ta	sk.													
	0				Di	sab	le																								
	1				En	nabl	e																								

#### 20.2.2 VALUE

Bit number	31 30 29 2	8 27 26 25 24 23 22 21 20 19	18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)			A A A A A A A A
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field	Value ID Value	Description	
A R	[0255]	Generated random number.	



# 20.2.3 CONFIG

Bit	num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (	Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- A
Res	et va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID	RW	Field	Value ID	Va	lue			De	scri	pti	on																							
А	RW	DERCEN						D	igita	al er	ror	corr	ecti	ion																				
				0				D	isab	led																								
				1				Er	nabl	ed																								



# 21 Temperature sensor (TEMP)

## 21.1 Functional description

The temperature sensor measures the silicon die temperature.

The TEMP is started by triggering the START task. When the temperature measurement is completed, a DATARDY event will be generated and the result of the measurement can be read from the TEMP register.

When the temperature measurement is completed, the TEMP analog electronics power down to save power.

The TEMP only supports one-shot operation, meaning that every TEMP measurement has to be explicitly started using the START task.

## 21.2 Registers

Register	Offset	Description
TASKS		
START	0x000	Start temperature measurement.
STOP	0x004	Stop temperature measurement.
EVENTS		
DATARDY	0x100	Temperature measurement complete, data ready.
REGISTERS		
INTENSET	0x304	Interrupt enable set register.
INTENCLR	0x308	Interrupt enable clear register.
TEMP	0x508	Temperature.

Table 26 Register overview

#### 21.2.1 TEMP

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
R	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
ID Field	Va	lue			De	scri	ptio	on																									
A					pr	recis	t of t ion. ion p		•				uren	nen	t. Di	e te	mp	erat	ure	in °(	C, 2'	's co	mp	lem	nent	t fo	rma	nt, 0	.25	°C			



# 22 AES Electronic Codebook Mode Encryption (ECB)

# 22.1 Functional description

AES ECB is a single AES block encrypt hardware module.

AES ECB features:

- 128 bit AES encryption
- Supports standard AES ECB block encryption
- Memory pointer support
- DMA data transfer

AES ECB performs a 128 bit AES block encrypt. At the STARTECB task, data and key is loaded into the algorithm by EasyDMA. When output data has been written back to memory, the ENDECB event is triggered.

AES ECB can be stopped by triggering the STOPECB task.

#### 22.1.1 EasyDMA

The ECB implements an EasyDMA mechanism for reading and writing to and from the RAM. When the ENDECB or ERRORECB is generated EasyDMA stops accessing the RAM.

#### 22.1.2 ECB Data Structure

Input to the block encrypt and output from the block encrypt are stored in the same data structure. ECBDATAPTR should point to this data structure before STARTECB is initiated.

Property	Address offset	Description
KEY	0	16 byte AES key
CLEARTEXT	16	16 byte AES cleartext input block
CIPHERTEXT	32	16 byte AES ciphertext output block

#### Table 27 ECB data structure overview

#### 22.1.3 Shared resources

The ECB, CCM, and AAR share the same AES module. The ECB will always have lowest priority and if there is a sharing conflict during encryption, the ECB operation will be aborted and an ERRORECB event will be generated.



# 22.2 Registers

Register	Offset	Description
TASKS		
STARTECB	0x000	Start ECB block encrypt. If a crypto operation is already running in the AES core, the STARTECB task will not start a new encryption and an ERRORECB event will be triggered.
STOPECB	0x004	Abort a possible executing ECB operation. If a running ECB operation is aborted by STOPECB, the ERRORECB event is triggered.
EVENTS		
ENDECB	0x100	ECB block encrypt complete.
ERRORECB	0x104	ECB block encrypt aborted because of a STOPECB task or due to an error.
REGISTERS		
INTENSET	0x304	Interrupt enable set register.
INTENCLR	0x308	Interrupt enable clear register.
ECBDATAPTR	0x504	ECB block encrypt memory pointers.

Table 28 Register overview

## 22.2.1 ECBDATAPTR

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α.	A	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field	Value ID	Va	lue			De	scri	ptie	on																							
A						Pc	ointe	er to	o the	e EC	Bd	ata	stru	ctu	re (s	see	Tab	le 2	7 on	pa	ge 1	122)										



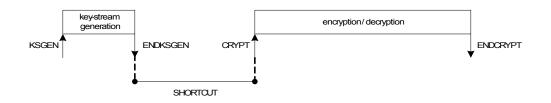
# 23 AES CCM Mode Encryption (CCM)

# 23.1 Functional description

The AES CCM supports three operations: key-stream generation, packet encryption, and packet decryption. All these operations are done in compliance with the *Bluetooth* specification<sup>1</sup>. A new key-stream must be generated before a new packet encryption or packet decryption operation can be started.

A key-stream is generated by triggering the KSGEN task. An ENDKSGEN event will be generated when the new key-stream has been generated. The key-stream will be stored in the AES CCM's temporary memory area, specified by the SCRATCHPTR, where it will be used in subsequent encryption and decryption operations.

Encryption is started by triggering the CRYPT task with the MODE register set to ENCRYPTION. Similarly, decryption is started by triggering the same task with MODE set to DECRYPTION. An ENDCRYPT event will be generated when packet encryption is completed as well as when packet decryption is completed, see *Figure 47*.



*Figure 47 Key-stream generation followed by encryption or decryption. The shortcut is optional.* 

Key-stream generation, packet encryption, and packet decryption operations utilize the configuration specified in the data structure pointed to by the CNFPTR pointer. It is necessary to configure this pointer and its underlying data structure, and the MODE register before the KSGEN task is triggered. It is also necessary to configure the INPTR pointer and the OUTPTR pointer before the CRYPT task is triggered.

If a shortcut is used between ENDKSGEN event and CRYPT task, the INPTR pointer and the OUTPTR pointer must be configured before the KSGEN task is triggered.

#### 23.1.1 Encryption

During packet encryption the AES CCM will read the unencrypted packet located in RAM at the address specified in the INPTR pointer, encrypt the packet and append a four byte long Message Integrity Check (MIC) field to the packet. The AES CCM will also modify the length field of the packet to adjust for the appended MIC field, that is, add four bytes to the length, and store the resulting packet back into RAM at the address specified in the OUTPTR pointer, see *Figure 48* on page 125.

Empty packets (length field is set to 0) will not be encrypted but instead moved unmodified through the AES CCM.

The AES CCM is limited to read maximum 27 bytes of the unencrypted payload (PL) regardless of what is specified in the length field of the unencrypted packet.

<sup>1.</sup> Bluetooth AES CCM 128 bit block encryption, see Bluetooth specification Version 4.0.



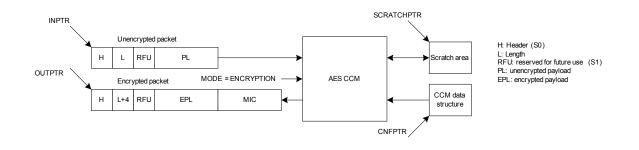


Figure 48 Encryption

#### 23.1.2 Decryption

During packet decryption the AES CCM will read the encrypted packet located in RAM at the address specified in the INPTR pointer, decrypt the packet, authenticate the packet's MIC field and generate the appropriate MIC status. The AES CCM will also modify the length field of the packet to adjust for the MIC field, that is, subtract four bytes from the length, and then store the decrypted packet back into RAM at the address pointed to by the OUTPTR pointer, see *Figure 49*.

The CCM is only able to decrypt packets that are at least 5 bytes long, that is, 1 byte encrypted payload (EPL) and 4 bytes of MIC. The CCM will therefore generate a MIC error for packets where the length field is set to 1, 2, 3 or 4.

Empty packets (length field is set to 0) will not be decrypted but instead moved unmodified through the AES CCM, these packets will always pass the MIC check.

The AES CCM is limited to read maximum 27 bytes of the encrypted payload and four bytes of the MIC regardless of what is specified in the length field of the encrypted packet.

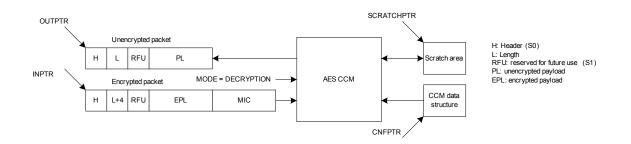


Figure 49 Decryption



#### 23.1.3 AES CCM and RADIO concurrent operation

The AES CCM is designed to run in parallel with the RADIO to enable on-the-fly encryption and decryption of RADIO packets without CPU involvement. To facilitate this, the RADIO has to be configured with the following settings:

Radio parameter	Value	Description
PCNF0.S0LEN	1	S0 field = 1 byte. Must match with the HEADER property in the data structure associated with INPTR and OUTPTR, see <i>Table 31</i> on page 129 and <i>Table 32</i> on page 129
PCNF0.LFLEN	5	Length field = 5 bit. Must match with the LENGTH property in the data structure associated with INPTR and OUTPTR, see <i>Table 31</i> on page 129 and <i>Table 32</i> on page 129
PCNF0.S1LEN	3	S1 field = 3 bit. Must match with the RFU property in the data structure associated with INPTR and OUTPTR, see <i>Table 31</i> on page 129 and <i>Table 32</i> on page 129
MODE	1_MBIT_QPSK	1 Mbps data rate
PCNF1.BALEN	3	Length of address (32 bit)
CRCCNF.LEN	3	Length of CRC (24 bit)

Table 29 Radio configuration settings

#### 23.1.4 Encrypting packets on-the-fly in radio transmit mode

When the AES CCM is encrypting a packet on-the-fly at the same time as the RADIO is transmitting it, the RADIO must read the encrypted packet from the same memory location as the AES CCM is writing to. The OUTPTR pointer in the AES CCM must therefore point to the same memory location as the PACKETPTR pointer in the RADIO, see *Figure 50*.

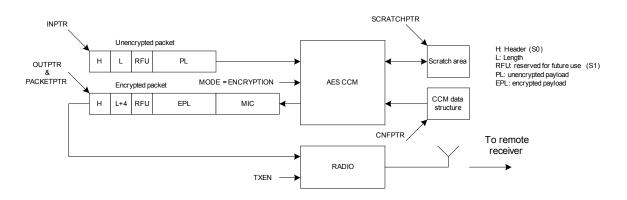


Figure 50 Configuration of on-the-fly encryption

In order to match the RADIO's timing, the KSGEN task must be triggered no later than when the START task in the RADIO is triggered, in addition the shortcut between the ENDKSGEN event and the CRYPT task must be enabled. This use-case is illustrated in *Figure 51* on page 127 using a PPI connection between the READY event in the RADIO and the KSGEN task in the AES CCM.



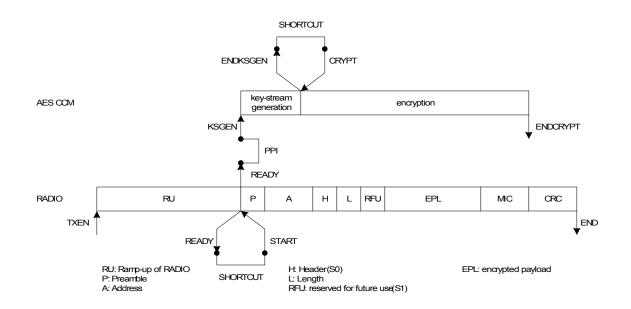


Figure 51 On-the-fly encryption using a PPI connection

#### 23.1.5 Decrypting packets on-the-fly in radio receive mode

When the AES CCM is decrypting a packet on-the-fly at the same time as the RADIO is receiving it, the AES CCM must read the encrypted packet from the same memory location as the RADIO is writing to. The INPTR pointer in the AES CCM must therefore point to the same memory location as the PACKETPTR pointer in the RADIO, see *Figure 52*.

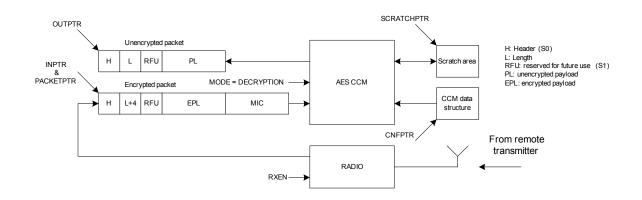


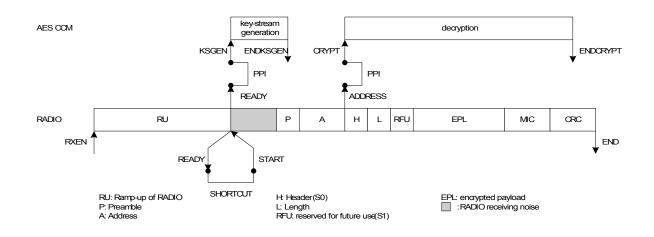
Figure 52 Configuration of on-the-fly decryption

In order to match the RADIO's timing, the KSGEN task must be triggered no later than when the START task in the RADIO is triggered. In addition, the CRYPT task must be triggered no earlier than when the ADDRESS event is generated by the RADIO.

If the CRYPT task is triggered exactly at the same time as the ADDRESS event is generated by the RADIO, the AES CCM will guarantee that the decryption is completed no later than when the END event in the RADIO is generated.



This use-case is illustrated in *Figure 53* using a PPI connection between the ADDRESS event in the RADIO and the CRYPT task in the AES CCM. The KSGEN task is trigger from the READY event in the RADIO through a PPI connection.



*Figure 53* On-the-fly decryption using a PPI connection between the READY event in the RADIO and the KSGEN task in the AES CCM

#### 23.1.6 CCM data structure

The CCM data structure specified in *Table 30* is located in RAM at the memory location specified by the CNFPTR pointer register.

Property	Address offset	Description
KEY	0	16 byte AES key
PKTCTR	16	Octet0 (LSO) of packet counter
	17	Octet1 of packet counter
	18	Octet2 of packet counter
	19	Octet3 of packet counter
	20	Bit 6 – Bit 0: Octet4 (7 most significant bits of packet counter, with Bit 6 being the most significant bit) Bit7: Ignored
	21	Ignored
	22	Ignored
	23	Ignored
DIRECTION	24	Bit 0: Direction bit Bit 7 – Bit 1: Zero padded
IV	25	8 byte initialization vector (IV) Octet0 (LSO) of IV, Octet1 of IV, , Octet7 (MSO) of IV

Table 30 CCM data structure overview



The NONCE vector (as specified by the *Bluetooth* Core Specification) will be generated by hardware based on the information specified in the CNFPTR data structure.

Property	Address offset	Description	
HEADER	0	Packet Header	
LENGTH	1	Number of bytes in unencrypted payload	
RFU	2	Reserved Future Use	
PAYLOAD	3	0 to 27 bytes unencrypted payload	

 Table 31 Data structure for unencrypted packet

Property	Address offset	Description	
HEADER	0	Packet Header	
LENGTH	1	Number of bytes in encrypted payload including length of MIC Note: LENGTH will be 0 for empty packets since the MIC is not added to empty packets	
RFU	2	Reserved Future Use	
PAYLOAD	3	0 to 27 bytes encrypted payload	
MIC	3 + payload length	ENCRYPT: 4 bytes encrypted MIC Note: MIC is not added to empty packets	

 Table 32 Data structure for encrypted packet

#### 23.1.7 EasyDMA and ERROR event

The CCM implements an EasyDMA mechanism for reading and writing to the RAM.

In some scenarios where the CPU and other DMA enabled peripherals are accessing the RAM at the same time, the CCM DMA could experience some bus conflicts which may also result in an error during encryption. If this happens, the ERROR event will be generated.

EasyDMA finishes accessing the RAM when the ENDKSGEN and ENDCRYPT events are generated.

#### 23.1.8 Shared resources

The CCM shares registers and other resources with other peripherals that have the same ID as the CCM. The user must therefore disable all peripherals that have the same ID as the CCM before the CCM can be configured and used. Disabling a peripheral that have the same ID as the CCM will not reset any of the registers that are shared with the CCM. It is therefore important to configure all relevant CCM registers explicitly to secure that it operates correctly.

The Instantiation table, *Table 3* on page 13, shows which peripherals have the same ID as the CCM.



# 23.2 Registers

Register	Offset	Description
TASKS		
KSGEN	0x000	Start generation of key-stream. This operation will stop by itself when completed.
CRYPT	0x004	Start encryption/decryption. This operation will stop by itself when completed.
STOP	0x008	Stop encryption/decryption
EVENTS		
ENDKSGEN	0x100	Key-stream generation complete
ENDCRYPT	0x104	Encrypt/decrypt complete
ERROR	0x108	CCM error event
REGISTERS		
SHORTS	0x200	Shortcut register
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
MICSTATUS	0x400	MIC check result
ENABLE	0x500	Enable
MODE	0x504	Operation mode
CNFPTR	0x508	Pointer to data structure holding AES key and NONCE vector
INPTR	0x50C	Input pointer
OUTPTR	0x510	Output pointer
SCRATCHPTR	0x521	Pointer to data area used for temporary storage

#### 23.2.1 SHORTS

Bit ı	numb	er		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	21	0
ID (I	Field I	D)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		A
Res	et val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0
ID F	RW	Field	Value ID	Va	lue			De	scri	ipti	on																							
			10							-																								

## 23.2.2 MICSTATUS

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field II	/alue D	Val	ue			De	scri	ipti	on																								
A R						op	oera	tio	า.	the iled		C ch	eck	perf	forn	ned	duı	ring	the	pre	evio	us d	lecr	ypt	ion	I							



#### 23.2.3 ENABLE

Bit number	3	13	0 29	28	3 27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9 8	8 7	76	5 5	5 4	43	32	1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-				-				A	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	0 0	) (	) (	0 0	) (	0	0
ID RW Field ID	alue )	alu	9		De	escri	ipti	on																							

#### 23.2.4 MODE

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	21	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) 1
ID RW	Field	Value ID	Va	lue			De	scri	iptio	on																							

#### 23.2.5 CNFPTR

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW	Field	Value ID	Va	lue			De	escr	ipti	on																								
A RW								oint ee 1						ctur												NCI	Ev	ect	or					

#### 23.2.6 INPTR

Bit numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (	)
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	١
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	)
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																								
A RW							In	put	poi	nte	r																							

#### 23.2.7 OUTPTR

Bit numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	1 0	
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A	
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																								
A RW							0	utpi	ut p	oint	er																							



# 23.2.8 SCRATCHPTR

Bit numb	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 :	3 2	2 1	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	۱ A	A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	) 0	0
ID RW	Field	Value ID	Va	lue			De	scri	ipti	on																							
A													dat nera						•				e dı	urin	g k	ey-	stre	ear	n				
													ised rypt			•			-					-					I.				



# 24 Accelerated Address Resolver (AAR)

# 24.1 Functional description

#### 24.1.1 Resolving a resolvable address

A private resolvable address shall be composed of 6 bytes as illustrated in *Figure 54*.

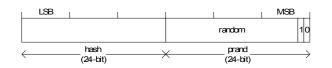


Figure 54 Resolvable address

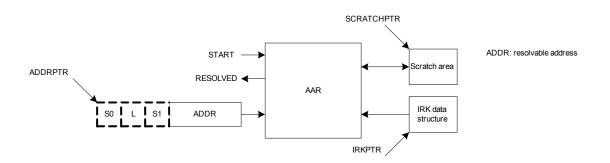
To resolve an address the ADDRPTR pointer must point to the least significant byte (LSB) of the resolvable address offset by 3 bytes to accommodate the packet header. The resolver is started by triggering the START task. A RESOLVED event is generated when and if the AAR manages to resolve the address using one of the Identity Resolving Keys (IRK) found in the IRK data structure. The AAR will use the IRK specified in the register IRK0 to IRK15 starting from IRK0. How many to be used is specified by the NIRK register. The AAR module will generate a NOTRESOLVED event if it is not able to resolve the address using the specified list of IRKs.

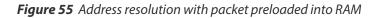
The AAR will go through the list of available IRKs in the IRK data structure and for each IRK try to resolve the address according to the Resolvable Private Address Resolution Procedure described in the Bluetooth

Specification<sup>2</sup>. The time it takes to resolve an address may vary depending on where in the list the resolvable address is located. The resolution time will also be affected by RAM accesses performed by other peripherals and the CPU. See product specification for more information about resolution time.

The AAR will not distinguish between public and random addresses. The AAR will also not distinguish between static and private addresses, or between private resolvable and private non-resolvable addresses.

The AAR will stop as soon as it has managed to resolve the address, or after trying to resolve the address using NIRK number of IRKs from the IRK data structure. The AAR will generate an END event after it has stopped.





<sup>2.</sup> Bluetooth Specification Version 4.0 [Vol 3] chapter 10.8.2.3.



# 24.1.2 Use case example for chaining RADIO packet reception with resolving addresses with the AAR

The AAR may be started as soon as the 6 bytes required by the AAR has been received by the RADIO and stored in RAM. The ADDRPTR pointer must point to the least significant byte of the resolvable address within the received packet.

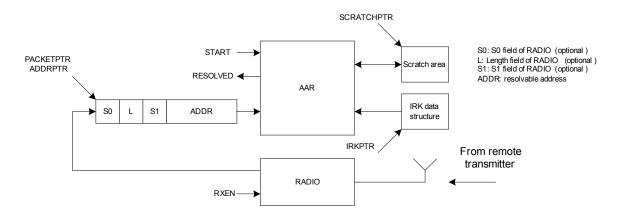


Figure 56 Address resolution with packet loaded into RAM by the RADIO

#### 24.1.3 IRK data structure

The IRK data structure specified in *Table 33* is located in RAM at the memory location specified by the CNFPTR pointer register.

Property	Address offset	Description	
IRK0	0	IRK number 0 (16 – byte)	
IRK1	16	IRK number 1 (16 – byte)	
IRK15	240	IRK number 15 (16 – byte)	

 Table 33 IRK data structure overview

#### 24.1.4 EasyDMA

The AAR implements EasyDMA for reading and writing to the RAM. EasyDMA finishes accessing the RAM when the END, RESOLVED, and NOTRESOLVED events are generated.

#### 24.1.5 Shared resources

The AAR shares registers and other resources with other peripherals that have the same ID as the AAR. The user must therefore disable all peripherals that have the same ID as the AAR before the AAR can be configured and used. Disabling a peripheral that have the same ID as the AAR will not reset any of the registers that are shared with the AAR. It is therefore important to configure all relevant AAR registers explicitly to secure that it operates correctly. The Instantiation table, *Table 3* on page 13, shows which peripherals have the same ID as the AAR.



# 24.2 Registers

Register	Offset	Description
TASKS		
START	0x000	Start resolving addresses based on IRKs specified in the IRK data structure
STOP	0x008	Stop resolving addresses
EVENTS		
END	0x100	Address resolution procedure complete
RESOLVED	0x104	Address resolved
NOTRESOLVED	0x108	Address not resolved
REGISTERS		
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
STATUS	0x400	Resolution status
ENABLE	0x500	Enable AAR
NIRK	0x504	Number of IRKs
IRKPTR	0x508	Pointer to IRK data structure
ADDPTR	0x510	Pointer to the resolvable address
SCRATCHPTR	0x514	Pointer to data area used for temporary storage

## 24.2.1 STATUS

Bit numb	er		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	54	3	2	1	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-			A	A	A	A
Reset val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) ()	0	0	0
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																							
A R			0.	.15			Th	ne IF	RK tł	nat	was	use	d la	st ti	me	an a	ddi	ess	was	res	olv	ed											

#### 24.2.2 ENABLE

Bit numb	er		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- /	A A
Reset val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW	Field	Value	Va	lue			De	scri	inti	on																							
		ID							Pu																								



#### 24.2.3 NIRK

Bit numb	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 :	3 2	2 1	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	A	A /	A A	A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	) ()	1
ID RW	Field	Value ID	Va	lue			De	scri	ptie	on																							
A RW			1.	.16			Νι	umb	oer o	of id	lent	ity r	oot	key	's av	aila	ble	in tl	he ll	RK d	lata	stru	uctu	ire									

## 24.2.4 IRKPTR

Bit numb	per		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	10
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A
Reset val	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																							
A RW							Po	ointe	er to	o th	e IRI	K da	ita s	truc	ture	e																	

#### 24.2.5 ADDRPTR

Bit numbe	er	3	81	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 3	3	21	0
ID (Field II	)	ł	١	A	A	A	A	Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α .	A A	AA
Reset valu	e	C	)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0	0 0	0
ID RW	Field Val	ue ۱	/al	ue			De	scri	ptio	on																							
A RW							Рс	ointe	er to	o res	solv	able	ad	dres	s (6	byt	es)																

#### 24.2.6 SCRATCHPTR

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6	5 4	4 :	3 2	2 1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0	0	0 0	) ()	0
ID RW Field	Value ID	Va	lue			De	scri	ptio	on																							
A RW						Po	ointe	er to	o a 's	scra	tch'	dat	a ar	e us	ed f	for t	tem	pora	ary s	stora	age	dur	ing	res	olu	itio	n					
						A	spa	ce o	of m	inin	num	13 k	yte	s mi	ust l	oe r	esei	rvec														



# 25 Serial Peripheral Interface (SPI) Master

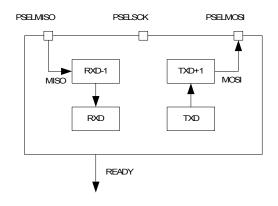


Figure 57 SPI master

**Note:** RXD-1 and TXD+1 illustrate the double buffered version of RXD and TXD respectively.

# 25.1 SPI master - functional description

The SPI master provides a simple CPU interface which includes a TXD register for sending data and an RXD register for receiving data. These registers are double buffered to enable some degree of uninterrupted data flow in and out of the SPI master. The SPI master does not implement support for chip select directly. Therefore, the CPU must use available GPIOs to select the correct slave and control this independently of the SPI master. The SPI master supports SPI modes 0 through 3.

#### 25.1.1 SPI master mode pin configuration

The different signals SCK, MOSI, and MISO associated with the SPI master are mapped to physical pins according to the configuration specified in the PSELSCK, PSELMOSI, and PSELMISO registers respectively. If a value of 0xFFFFFFFF is specified in any of these registers, the associated SPI master signal is not connected to any physical pin. The PSELSCK, PSELMOSI, and PSELMISO registers and their configurations are only used as long as the SPI master is enabled, and retained only as long as the device is in ON mode,

To secure correct behavior in the SPI, the pins used by the SPI must be configured in the GPIO peripheral as described in *Table 34* prior to enabling the SPI. The SCK must always be connected to a pin, and that pin's input buffer must always be connected for the SPI to work. This configuration must be retained in the GPIO for the selected IOs as long as the SPI is enabled.

Only one peripheral can be assigned to drive a particular GPIO pin at a time, failing to do so may result in unpredictable behavior.

SPI master signal	SPI master pin	Direction	Output value
SCK	As specified in PSELSCK	Output	Same as CONFIG.CPOL
MOSI	As specified in PSELMOSI	Output	0
MISO	As specified in PSELMISO	Input	Not applicable

Table 34 GPIO configuration



#### 25.1.2 Shared resources

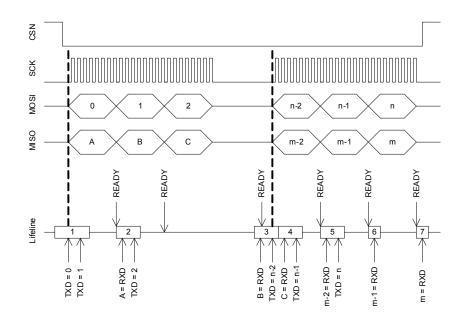
The SPI shares registers and other resources with other peripherals that have the same ID as the SPI. The user must therefore disable all peripherals that have the same ID as the SPI before the SPI can be configured and used. Disabling a peripheral that have the same ID as the SPI will not reset any of the registers that are shared with the SPI. It is therefore important to configure all relevant SPI registers explicitly to secure that it operates correctly.

The Instantiation table in *Section 4.2 "Instantiation"* on page 13 shows which peripherals have the same ID as the SPI.

#### 25.1.3 SPI master transaction sequence

An SPI master transaction is started by writing the first byte, which is to be transmitted by the SPI master, to the TXD register. Since the transmitter is double buffered, the second byte can be written to the TXD register immediately after the first one. The SPI master will then send these bytes in the order they are written to the TXD register.

The SPI master is a synchronous interface, and for every byte that is sent, a different byte will be received at the same time; this is illustrated in *Figure 58*. Bytes that are received will be moved to the RXD register where the CPU can extract them by reading the register. The RXD register is double buffered in the same way as the TXD register, and a second byte can therefore be received at the same time as the first byte is being extracted from RXD by the CPU. The SPI master will generate a READY event every time a new byte is moved to the RXD register. The double buffered byte will be moved from RXD-1 to RXD as soon as the first byte is extracted from RXD. The SPI master will stop when there are no more bytes to send in TXD and TXD+1.

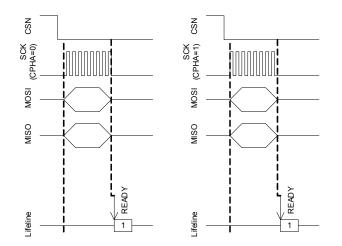


#### Figure 58 SPI master transaction.

The READY event of the third byte transaction is delayed until B is extracted from RXD in occurrence number 3 on the horizontal lifeline. The reason for this is that the third event is generated first when C is moved from RXD-1 to RXD after B is read.



The SPI master will move the incoming byte to the RXD register after a short delay following the SCK clock period of the last bit in the byte. This also means that the READY event will be delayed accordingly, see *Figure 59*. Therefore, it is important that you always clear the READY event, even if the RXD register and the data that is being received is not used.



*Figure 59* READY event timing details for SPI master mode for CPHA = 0 and CPHA = 1.

# 25.2 Registers

Register	Offset	Description
EVENTS		
READY	0x108	TXD byte sent and RXD byte received
REGISTERS		
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
ENABLE	0x500	Enable SPI
PSELSCK	0x508	Pin select for SCK
PSELMOSI	0x50C	Pin select for MOSI
PSELMISO	0x510	Pin select for MISO
RXD	0x518	RXD register
TXD	0x51C	TXD register
FREQUENCY	0x524	SPI frequency
CONFIG	0x554	Configuration register

Table 35Register overview



#### 25.2.1 RXD

Bit numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (	,
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	<b>A</b> 1	A A	•
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0 0	,
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																								
A RW							R۶	۲ da	ta r	ecei	ved	l. Do	bubl	e bı	uffe	red.																		

#### 25.2.2 TXD

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 <sup>.</sup>	1 0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	Α.	A	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field	Value ID	Va	lue	•		De	scri	ipti	on																							
A RW						T)	۲ da	ta t	o se	end.	Do	uble	e bu	ffer	ed																	

## 25.2.3 ENABLE

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6	5	4 :	32	1	0
ID (Field ID)	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 A	A	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	0	0
ID RW Field Value ID	Val	lue			De	scri	pti	on																							
A RW					Er	nabl	e or	<sup>r</sup> dis	able	e SP	'I																				
DISABLE	0				Di	sab	le S	ΡI																							
ENABLE	1				Er	nabl	e SF	יו																							

## 25.2.4 PSELSCK

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6 !	5	4 :	3 2	2 1	0
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	۹ A	A
Reset value	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	1	1
ID RW Field Value ID	Va	lue			De	scri	ptio	on																							
А					Pi	n nu	ımb	oer o	onf	igu	rati	on fe	or S	PI S	CK :	sign	al														
RW	[0.	31]			Pi	n nu	ımb	oer t	o rc	oute	the	e SPI	SC	K sig	gna	l to															
W	0x	FFF	FFFF	F	Di	scor	nne	ct																							



## 25.2.5 PSELMOSI

Bit numbe	r		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 :	3 2	2 1	0
ID (Field ID	<b>)</b> )		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α.	Α	A	A /	A A	A
Reset valu	e		1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	1	1
ID RW Fi	ield	Value ID	Va	lue			De	scri	ptic	on																							
А							Pir	n nu	mb	er c	onfi	gur	atio	n fo	r SP	IMC	DSI s	ign	al														
RW			[0.	31]			Pir	n nu	mb	er to	o ro	ute	the	SPI	MO	SI sig	gna	l to															
W			0x	FFF	FFF	F	Di	scor	nne	ct																							

#### 25.2.6 PSELMISO

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8 7	7 E	5 5	4	3	2	10
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A	۱, ł	A	A	A	A A
Reset value	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 '	11	1	1	1	1	11
ID RW Field Value ID	Val	lue			De	scri	ptio	on																						
A					Pi	n nu	ımb	oer c	onf	igur	ratic	on fo	or SF	'l m	aste	er M	ISO	sigr	al											
RW	[0.	31]			Pi	n nı	ımb	oer t	o ro	ute	the	SPI	mas	ster	MIS	50 s	igna	l to												
W	0x	FFF	FFFF	F	Di	sco	nne	ct																						

#### 25.2.7 CONFIG

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	54	3	2	10
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		• •	-	С	ΒA
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0 0
ID RW Field	Value ID	Va	lue			De	scri	ptic	on																						
A RW ORD	ER					Bit	or	der																							
	MSBFIRST	0				Mo	ost	sign	ifica	ant l	oit s	hift	ed o	ut f	irst																
	LSBFIRST	1				Le	ast	sign	ifica	ant	bit s	hift	ed c	out f	first																
B RW CPO	-					Se	rial	cloc	:k (S	SCK)	pol	arit	у																		
	ACTIVEHIGH	0				Ac	tive	e hig	h																						
	ACTIVELOW	1				Ac	tive	lov	v																						
C RW CPH	Ą					Se	rial	cloc	:k (S	SCK)	pha	ase																			
	LEADING	0				Sa	mp	le oi	n le	adir	ng e	dge	of	loc	k, sł	nift s	seria	al da	ta c	on tr	ailin	g e	dge								
	TRAILING	1				Sa	mp	le o	n tra	ailin	g eo	dge	of c	lock	k, sh	ift s	eria	l da	ta o	n le	adin	g e	dge								



# 25.2.8 FREQUENCY

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field Value ID	Va	lue			De	scri	pti	on																							
A RW					SF	۷ m	aste	er d	ata i	rate																					
K125	0>	x020	000	000	12	25 k	bps																								
K250	0>	x040	000	000	25	50 k	bps																								
K500	0>	x080	000	000	50	00 k	bps																								
M1	0>	x100	000	000	1	Mbp	os																								
M2	0>	x200	000	000	2	Mbp	os																								
M4	0>	x400	000	000	4	Mbp	os																								
M8	0>	x800	000	000	8	Mbp	os																								



# 26 SPI Slave (SPIS)

SPIS is a SPI slave with EasyDMA support for ultra low power serial communication from an external SPI master. EasyDMA in conjunction with hardware based semaphore mechanisms removes all real-time requirements associated with controlling the SPI slave from a low priority CPU execution context.

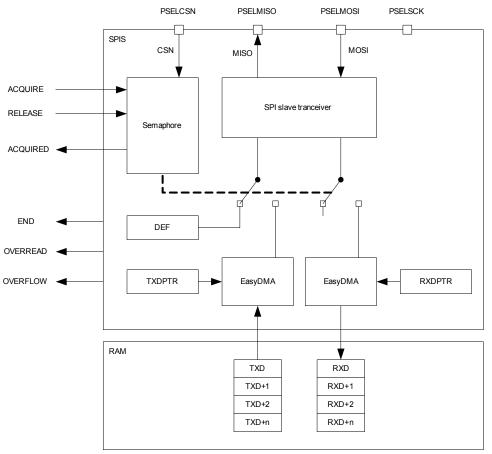


Figure 60 SPI slave

# 26.1 Pin configuration

The different signals CSN, SCK, MOSI, and MISO associated with the SPI slave are mapped to physical pins according to the configuration specified in the PSELCSN, PSELSCK, PSELMOSI, and PSELMISO registers respectively. If a value of 0xFFFFFFFF is specified in any of these registers, the associated SPI slave signal will not be connected to any physical pins.

The PSELCSN, PSELSCK, PSELMOSI, and PSELMISO registers and their configurations are only used as long as the SPI slave is enabled, and retained only as long as the device is in System ON mode, see *Chapter 11 "Power management (POWER)"* on page 40 for more information about power modes.

To secure correct behavior in the SPI slave, the pins used by the SPI slave must be configured in the GPIO peripheral as described in *Table 36* on page 144 prior to enabling the SPI slave. This is to secure that the pins used by the SPI slave are driven correctly if the SPI slave itself is temporarily disabled or the device temporarily enters System OFF. This configuration must be retained in the GPIO for the selected I/Os as long as the SPI slave is to be recognized by an external SPI master.

The MISO line is set in high impedance as long as the SPI slave is not selected with CSN.



Only one peripheral can be assigned to drive a particular GPIO pin at a time, failing to do so may result in unpredictable behavior.

SPI signal	SPI pin	Direction	Output value	Comment
CSN	As specified in PSELCSN	Input	Not applicable	
SCK	As specified in PSELSCK	Input	Not applicable	
MOSI	As specified in PSELMOSI	Input	Not applicable	
MISO	As specified in PSELMISO	Input	Not applicable	Emulates that the SPI slave is not selected.

Table 36 Pin configuration

# 26.2 Shared resources

The SPI slave shares registers and other resources with other peripherals that have the same ID as the SPI slave. You must, therefore, disable all peripherals that have the same ID as the SPI slave before the SPI slave can be configured and used. Disabling a peripheral that has the same ID as the SPI slave will not reset any of the registers that are shared with the SPI slave. It is important to configure all relevant SPI slave registers explicitly to secure that it operates correctly.

The Instantiation table in *Section 4.2 "Instantiation"* on page 13 shows which peripherals have the same ID as the SPI slave.

# 26.3 EasyDMA

The SPI Slave implements EasyDMA for reading and writing to and from the RAM. The EasyDMA will have finished accessing the RAM when the END event is generated.

# 26.4 SPI slave operation

SPI slave uses two memory pointers, RXDPTR and TXDPTR, that point to the RXD buffer (receive buffer) and TXD buffer (transmit buffer) respectively, see *Figure 60* on page 143. Since these buffers are located in RAM, which can be accessed by both the SPI slave and the CPU, a hardware based semaphore mechanism is implemented to enable safe sharing.

Before the CPU can safely update the RXDPTR and TXDPTR pointers it must first acquire the SPI semaphore. The CPU can acquire the semaphore by triggering the ACQUIRE task and then receiving the ACQUIRED event. When the CPU has updated the RXDPTR and TXDPTR pointers the CPU must release the semaphore before the SPI slave will be able to acquire it. The CPU releases the semaphore by triggering the RELEASE task. This is illustrated in *Figure 61* on page 146. Triggering the RELEASE task when the semaphore is not granted to the CPU will have no effect.

The semaphore mechanism does not, at any time, prevent the CPU from performing read or write access to the RXDPTR register, the TXDPTR registers, or the RAM that these pointers are pointing to. The semaphore is only telling when these can be updated by the CPU so that safe sharing is achieved.

The semaphore is by default assigned to the CPU after the SPI slave is enabled. No ACQUIRED event will be generated for this initial semaphore handover. An ACQUIRED event will be generated immediately if the ACQUIRE task is triggered while the semaphore is assigned to the CPU.

The SPI slave will try to acquire the semaphore when CSN goes low. If the SPI slave does not manage to acquire the semaphore at this point, the transaction will be ignored. This means that all incoming data on



MOSI will be discarded, and the DEF (default) character will be clocked out on the MISO line throughout the whole transaction. This will also be the case even if the semaphore is released by the CPU during the transaction. In case of a race condition where the CPU and the SPI slave try to acquire the semaphore at the same time, as illustrated in lifeline item 2 in *Figure 61* on page 146, the semaphore will be granted to the CPU.

If the SPI slave acquires the semaphore, the transaction will be granted. The incoming data on MOSI will be stored in the RXD buffer and the data in the TXD buffer will be clocked out on MISO.

When a granted transaction is completed and CSN goes high, the SPI slave will automatically release the semaphore and generate the END event.

As long as the semaphore is available the SPI slave can be granted multiple transactions one after the other. If the CPU is not able to reconfigure the TXDPTR and RXDPTR between granted transactions, the same TX data will be clocked out and the RX buffers will be overwritten. To prevent this from happening, the END\_ACQUIRE shortcut can be used. With this shortcut enabled the semaphore will be handed over to the CPU automatically after the granted transaction has completed, giving the CPU the ability to update the TXDPTR and RXDPTR between every granted transaction.

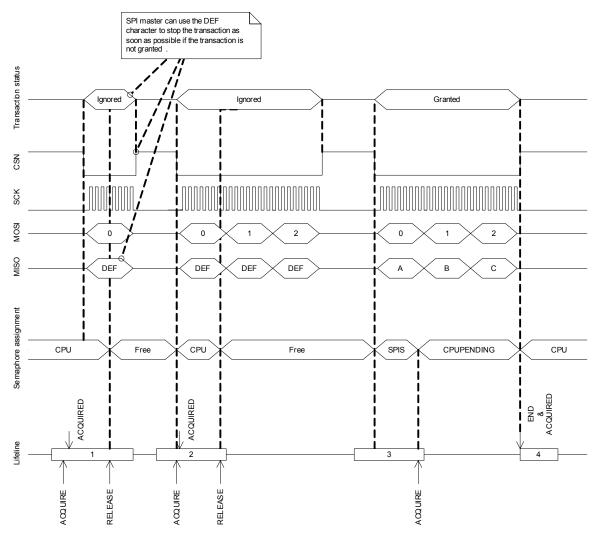
If the CPU tries to acquire the semaphore while it is assigned to the SPI slave, an immediate handover will not be granted. However, the semaphore will be handed over to the CPU as soon as the SPI slave has released the semaphore after the granted transaction is completed. If the END\_ACQUIRE shortcut is enabled and the CPU has triggered the ACQUIRE task during a granted transaction, only one ACQUIRE request will be served following the END event.

The MAXRX register specifies the maximum number of bytes the SPI slave can receive in one granted transaction. If the SPI slave receives more than MAXRX number of bytes, an OVERFLOW will be indicated in the STATUS register and the incoming bytes will be discarded.

The MAXTX parameter specifies the maximum number of bytes the SPI slave can transmit in one granted transaction. If the SPI slave is forced to transmit more than MAXTX number of bytes, an OVERREAD will be indicated in the STATUS register and the ORC character will be clocked out.

The AMOUNTRX and AMOUNTTX registers are updated when a granted transaction is completed. The AMOUNTTX register indicates how many bytes were read from the TX buffer in the last transaction, that is, ORC (over-read) characters are not included in this number. Similarly, the AMOUNTRX register indicates how many bytes were written into the RX buffer in the last transaction.





*Figure 61* SPI transaction when shortcut between END and ACQUIRE is enabled.



# 26.5 Registers

Register	Offset	Description	Support
TASKS			
ACQUIRE	0x024	Acquire SPI semaphore.	S
RELEASE	0x028	Release SPI semaphore, enabling the SPI slave to acquire it.	S
EVENTS			
END	0x104	Granted transaction completed.	S
ACQUIRED	0x128	Semaphore acquired.	S
REGISTERS			
SHORTS	0x200	Shortcut for the SPI slave.	S
INTENSET	0x304	Interrupt enable set register.	S
INTENCLR	0x308	Interrupt enable clear register.	S
SEMSTAT	0x400	Semaphore status register.	S
STATUS	0x440	Status from last transaction.	S
ENABLE	0x500	Enable SPI slave.	S
PSELSCK	0x508	Pin select for SCK.	S
PSELMISO	0x50C	Pin select for MISO.	S
PSELMOSI	0x510	Pin select for MOSI.	S
PSELCSN	0x514	Pin select for CSN.	S
RXDPTR	0x534	RXD data pointer.	S
MAXRX	0x538	Maximum number of bytes in receive buffer.	S
AMOUNTRX	0x53C	Number of bytes received in last granted transaction.	S
TXDPTR	0x544	TXD data pointer.	S
MAXTX	0x548	Maximum number of bytes in transmit buffer.	S
AMOUNTTX	0x54C	Number of bytes transmitted in last granted transaction.	S
CONFIG	0x554	Configuration register.	S
DEF	0x55C	Default character. Character clocked out in case of an ignored transaction.	S
ORC	0x5C0	Over-read character. Character clocked out after an over-read of the transmit buffer.	S

Table 37 Register overview

## 26.5.1 SHORTS

Bit number	31 30 29 28	27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		A
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field Value ID	Value	Description
A RW END_ACQUIRE		Enable or disable shortcut between END event and ACQUIRE task.



### 26.5.2 ENABLE

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (Field ID)		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	A A	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field	Value ID		Va	lue		De	scri	pti	on																							
A RW						Er	nabl	e or	<sup>,</sup> dis	abl	e SF	9 sla	ave																			
	DISABLE			0		Di	sab	le S	PI s	lave	2																					
	ENABLE			2		Er	nabl	e SF	PI sl	ave																						

#### 26.5.3 SEMSTAT

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	2 1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		• -	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	<b>D</b> (	) ()	1
ID RW Field	Value ID	Va	lue			De	scr	ipti	on																							
A R						Se	ema	pho	ore	stat	us.																					
	Free	0				Se	ema	pho	ore i	is fre	ee.																					
	CPU	1				Se	ema	pho	ore i	is as	sigr	ned	to C	PU.																		
	SPIS	2				Se	ema	pho	ore i	is as	sigr	ned	to S	PI s	lave	<b>.</b>																
	CPUENDING	3				Se	ema	pho	ore i	is as	sigr	ned	to S	PI b	ut a	a ha	ndo	over	to t	he (	CPU	l is p	bend	din	g.							

## 26.5.4 STATUS

Individual status bits are cleared by writing a '1' to the bits that shall be cleared.

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	В	A
Reset va	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW	Field	Value ID	Va	lue			De	scr	ipti	on																								
A RW	OVERREAD						T	(hi	iffoi		ar_ra	hee	dot	ecte	nd ai	nd r	nev	/ent	-ed															

### 26.5.5 RXDPTR

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field Value ID	Va	lue			De	scri	pti	on																							
A RW					R)	KD c	lata	poi	inte	r																					



## 26.5.6 TXDPTR

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6	5	4 :	3 2	2 1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α.	A	A	A /	A A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0	0	0 (	0 0	) (	0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A RW						T)	KD o	lata	poi	nte	r																					

## 26.5.7 PSELCSN

Bit number	31	<b>30</b> (	29 2	28 2	7 26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	2 1	0
ID (Field ID)	Α	A	A A	A A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A	AA
Reset value	1	1	11	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 '	1	1
ID RW Field Value I	Val	ue		D	escr	ipti	on																							
A RW W		.31] FFFF	FFF	F	Pin n Pin n Disco	umb	oer t																							

## 26.5.8 PSELSCK

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A
Reset value		1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11
ID RW Field	Value ID	Va	lue			De	scri	ptio	on																							
A RW W		-	31 (FFF	-	FF	Pi		ımb	er t		-		on fe SP				-															

## 26.5.9 PSELMOSI

Bit number	31 30 29 28	8 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)	<b>A A A A</b>	A A A A A A A A A A A A A A A A A A A
Reset value	1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
ID RW Field Value	ID Value	Description
A RW W	[031] 0xFFFFFFFF	Pin number configuration for SPI MOSI signal. Pin number to route the SPI MOSI signal to. Disconnect.



## 26.5.10 PSELMISO

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /
Reset value	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1
ID RW Field Value I	D Va	lue			De	scri	ipti	on																							
A RW	ſ	)31	1									on f e SP																			

#### 26.5.11 CONFIG

Bit	num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	98	; 7	6	5	4	3	2	10
ID	(Field	IID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-	-	-	-	c	ΒA
Re	set va	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	0 0
ID	RW	Field	Value ID	Va	lue			De	scri	pti	on																						
A	RW	ORDER						Bi	t or	der.																							
			MSBFIRST	0				Μ	ost	sigr	nific	ant	bit	shif	ted	out	firs	t.															
			LSBFIRST	1				Le	ast	sigi	nific	cant	bit	shif	ted	out	firs	t.															
В	RW	CPHA						Se	erial	clo	ck (	SCK	) pł	nase																			
			LEADING	0				Sa	mp	le o	n le	eadi	ng e	edg	e of	clo	ck, s	shift	t ser	ial c	data	a on	trai	iling	edg	e.							
			TRAILING	1				Sa	mp	le o	on ti	railir	ng e	dge	of	cloc	:k, s	hift	ser	ial d	lata	on	lead	ding	edg	e.							
С	RW	CPOL						Se	erial	clo	ck (	SCK	) pc	olari	ty.									-	-								
			ACTIVEHIGH	0				Ac	tive	e hig	gh.																						
			ACTIVELOW	1				Ac	tive	e lov	w.																						

### 26.5.12 DEF

Bit number	31 30 29 2	8 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		AAAAAAAA
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field Value	ID Value	Description
A RW		Default character. Clocked out on MISO during an ignored transaction.

# 26.5.13 ORC

Bit number	31 30 29 28	27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		AAAAAAAA
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field Valu	ue ID Value	Description
A RW		Over-read character. Character clocked out after an over-read of the transmit buffer.



#### 26.5.14 MAXRX

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 3	32	1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	A	A	A	A	A A	A	Α
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	) (	0	0
ID RW Field	Value ID	Va	lue			De	scri	ipti	on																							
A RW						М	axir	nur	ո ու	ımb	er c	of by	ytes	in r	ece	ive	buf	fer.														

# 26.5.15 MAXTX

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	A	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A RW						М	axir	nun	n nı	ımb	er c	of by	ytes	in t	ran	smi	t bu	ffer														

## 26.5.16 AMOUNTRX

Bit number	31 30 29 28	27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)		A A A A
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field Va	lue ID Value	Description
A RW		Number of bytes received in the last granted transaction.

## 26.5.17 AMOUNTTX

Bit number	31 30 29	28 27 26 25 24	23 22 21 2	0 19 18 17 16 1	5 14 13 12 11 10 9 8 7	7 6 5 4 3 2 1 0
ID (Field ID)					4	<b>A A A A A A A</b>
Reset value	0 0 0	0 0 0 0 0	0 0 0 0	0 0 0 0 0	0 0 0 0 0 0 0 0	0 0 0 0 0 0 0
ID RW Field Valu	e ID Value	Description				
A RW		Number of b	ytes transmit	ted in the last grar	nted transaction.	



# 27 I<sup>2</sup>C compatible Two Wire Interface (TWI)

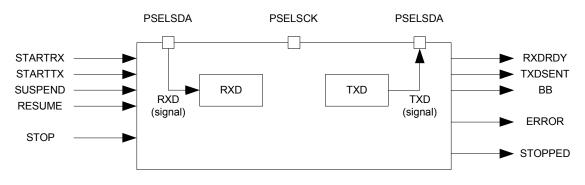
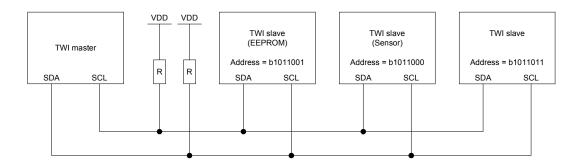


Figure 62 TWI master's main features

# 27.1 Functional description

The TWI master is compatible with I<sup>2</sup>C operating at 100 kHz and 400 kHz. This TWI master is not compatible with CBUS. As illustrated in *Figure 62*, the TWI transmitter and receiver are single buffered.

A TWI setup comprising one master and three slaves is illustrated in *Figure 63*. This TWI master is only able to operate as the only master on the TWI bus.





This TWI master supports clock stretching performed by the slaves. The TWI master is started by triggering the STARTTX or STARTRX tasks, and stopped by triggering the STOP task.

If a NACK is clocked in from the slave, the TWI master will generate an ERROR event.

# 27.2 Master mode pin configuration

The different signals SCL and SDA associated with the TWI master are mapped to physical pins according to the configuration specified in the PSELSCL and PSELSDA registers respectively. If a value of 0xFFFFFFF is specified in any of these registers, the associated TWI master signal is not connected to any physical pin. The PSELSCL and PSELSDA registers and their configurations are only used as long as the TWI master is enabled, and retained only as long as the device is in ON mode. To secure correct signal levels on the pins used by the TWI master when the system is in OFF mode, and when the TWI master is disabled, these pins must be configured in the GPIO peripheral as described in *Table 38* on page 153.



Only one peripheral can be assigned to drive a particular GPIO pin at a time, failing to do so may result in unpredictable behavior.

TWI master signal	TWI master pin	Direction	Drive strength	Output value
SCL	As specified in PSELSCL	Input	S0D1	Not applicable
SDA	As specified in PSELSDA	Input	S0D1	Not applicable

Table 38 GPIO configuration

# 27.3 Shared resources

The TWI shares registers and other resources with other peripherals that have the same ID as the TWI. Therefore, you must disable all peripherals that have the same ID as the TWI before the TWI can be configured and used. Disabling a peripheral that has the same ID as the TWI will not reset any of the registers that are shared with the TWI. It is therefore important to configure all relevant TWI registers explicitly to secure that it operates correctly.

The Instantiation table in *Section 4.2 "Instantiation"* on page 13 shows which peripherals have the same ID as the TWI.

# 27.4 Master write sequence

A TWI master write sequence is started by triggering the STARTTX task. After the STARTTX task has been triggered, the TWI master will generate a start condition on the TWI bus, followed by clocking out the address and the READ/WRITE bit set to 0 (WRITE=0, READ=1). The address must match the address of the slave device that the master wants to write to. The READ/WRITE bit is followed by an ACK/NACK bit (ACK=0 or NACK=1) generated by the slave.

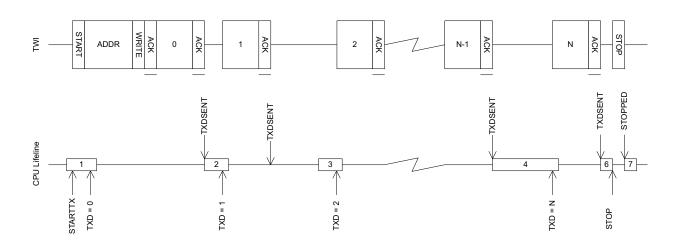
After receiving the ACK/NACK bit, the TWI master will clock out the data bytes that are written to the TXD register. Each byte clocked out from the master will be followed by an ACK/NACK bit clocked in from the slave. A TXDSENT event will be generated each time the TWI master has clocked out a TXD byte, and the associated ACK/NACK bit has been clocked in from the slave.

The TWI master transmitter is single buffered, and a second byte can only be written to the TXD register after the previous byte has been clocked out and the ACK/NACK bit clocked in, that is, after the TXDSENT event has been generated.

If the CPU is prevented from writing to TXD when the TWI master is ready to clock out a byte, the TWI master will stretch the clock until the CPU has written a byte to the TXD register.

A typical TWI master write sequence is illustrated in *Figure 64* on page 154. Occurrence 3 in *Figure 64* on page 154 illustrates delayed processing of the TXDSENT event associated with TXD byte 1. In this scenario the TWI master will stretch the clock to prevent writing erroneous data to the slave.





*Figure 64* The TWI master writing data to a slave.

The TWI master write sequence is stopped when the STOP task is triggered. When the STOP task is triggered, the TWI master will generate a stop condition on the TWI bus.

# 27.5 Master read sequence

A TWI master read sequence is started by triggering the STARTRX task. After the STARTRX task has been triggered the TWI master will generate a start condition on the TWI bus, followed by clocking out the address and the READ/WRITE bit set to 1 (WRITE = 0, READ = 1). The address must match the address of the slave device that the master wants to read from. The READ/WRITE bit is followed by an ACK/NACK bit (ACK=0 or NACK = 1) generated by the slave.

After having sent the ACK/NACK bit the TWI slave will send data to the master using the clock generated by the master.

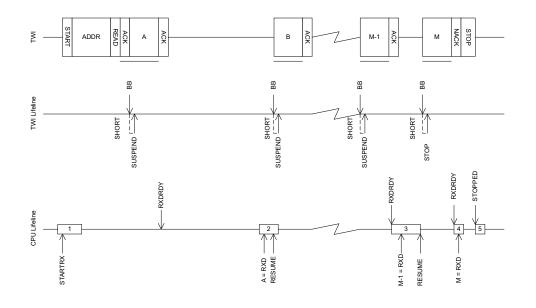
The TWI master will generate a RXDRDY event every time a new byte is received in the RXD register.

After receiving a byte, the TWI master will delay sending the ACK/NACK bit by stretching the clock until the CPU has extracted the received byte, that is, by reading the RXD register.

The TWI master read sequence is stopped by triggering the STOP task. This task must be triggered before the last byte is extracted from RXD to ensure that the TWI master sends a NACK back to the slave before generating the stop condition.

A typical TWI master read sequence is illustrated in *Figure 65* on page 155. Occurrence 3 in *Figure 65* on page 155 illustrates delayed processing of the RXDRDY event associated with RXD byte B. In this scenario the TWI master will stretch the clock to prevent the slave from overwriting the contents of the RXD register.



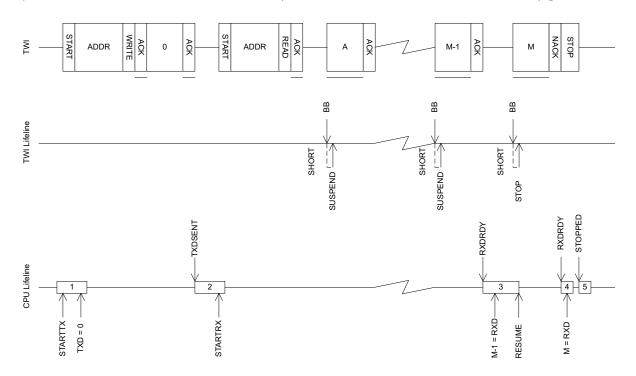


*Figure 65* The TWI master reading data from a slave.



# 27.6 Master repeated start sequence

*Figure 66* illustrates a typical repeated start sequence where the TWI master writes one byte to the slave followed by reading M bytes from the slave. Any combination and number of transmit and receive sequences can be combined in this fashion. Only one shortcut to STOP can be enabled at any given time.



**Figure 66** A repeated start sequence, where the TWI master writes one byte, followed by reading M bytes from the slave without performing a stop in-between

To generate a repeated start after a read sequence, a second start task must be triggered instead of the STOP task, that is, STARTRX or STARTTX. This start task must be triggered before the last byte is extracted from RXD to ensure that the TWI master sends a NACK back to the slave before generating the repeated start condition.



# 27.7 Registers

Register	Offset	Description
TASKS		
STARTRX	0x000	Start TWI receive sequence
STARTTX	0x008	Start TWI transmit sequence
STOP	0x014	Stop TWI transaction
SUSPEND	0x01C	Suspend TWI transaction
RESUME	0x020	Resume TWI transaction
EVENTS		
STOPPED	0x104	TWI stopped
RXDRDY	0x108	TWI RXD byte received
TXDSENT	0x11C	TWI TXD byte sent
ERROR	0x124	TWI error
BB	0x138	TWI byte boundary, generated before each byte that is sent or received
REGISTERS		
SHORTS	0x200	Shortcut register
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
ERRORSRC	0x4C4	TWI error source
ENABLE	0x500	Enable TWI master
PSELSCL	0x508	Pin select for SCL
PSELSDA	0x50C	Pin select for SDA
RXD	0x518	RXD register
TXD	0x51C	TXD register
FREQUENCY	0x524	TWI frequency
ADDRESS	0x588	Address used in the TWI transfer

Table 39 Register overview

# 27.7.1 SHORTS

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	87	6	5	4	3	2	1	D
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-	-	-	-	В	A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	D
ID RW	Field	Value ID	Va	lue			De	scri	pti	on																							
A RW	BB_SUSPEND		0 1				Di	nort sab nabl	le	: be	twe	en E	3B e	ven	t an	d Sl	JSP	ENC	) tas	sk													
B RW	BB_STOP		0 1				Di	nort sab nabl	le	be	twe	en E	3B e	ven	t an	d ST	ΓOP	tas	k														



#### 27.7.2 ERRORSRC

Bit numb	er		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-		•	-	-	-	-	-	-			-	-	-	-	-	CF	3 -
Reset val	ue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW	Field	Value ID	Va	lue			De	scri	ptic	on																							
							Description																										
B RW	ANACK		1			<b>Description</b> NACK received after sending the address (write '1' to clear)																											

### 27.7.3 ENABLE

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (Field	ID)		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	A /	A A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0 (
ID RW	Field	Value ID	Va	lue			De	scri	ipti	on																							
A RW	ENABLE						Er	nabl	e o	r dis	abl	e TV	٧I																				
		DISABLE	0				Di	isab	le T	WI																							
		ENABLE	5				Er	nabl	le T	WI																							

# 27.7.4 PSELSCL

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	32	: 1	0
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	Α	A	A	A	A	A	A	A	A	A	Α.	A	A	A A	A	A
Reset value	ue 11									1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11	1	1
ID RW Field Value ID	Val	ue			De	scri	pti	on																							
A					Pi	ո ու	umb	ber (	con	figu	rati	on f	or T	WI S	SCL	sig	nal														
RW	[0.	.31]			Pi	n nı	umb	bert	to ro	oute	e the	e TV	VI SO	CL si	igna	al to	)														
W	0x	FFFI	FFFF	F	Di	sco	nne	ct																							

# 27.7.5 PSELSDA

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	98	3 7	76	5	4	3	2	1	0
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A /	A A	A	A	A	A	A	A
Reset value	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	1 1	11	1	1	1	1	1	1
ID RW Field Value ID	Va	lue			De	scri	ptic	on																							
Α					Pi	n nu	ımb	er c	onf	igur	atio	n fo	or TV	VI SI	DAs	sign	al														
RW	[0	31]			Pi	n nı	ımb	er t	o ro	ute	the	τw	I SD	A sig	gna	l to															
W	0>	FFF	FFFF	F	Di	scoi	nne	ct																							



#### 27.7.6 RXD

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	б	5	4 3	32	2 1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	Α.	A /	A A	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	D	0	0 0	) (	0	0
ID RW Field	Value ID	Va	lue			De	scri	ipti	on																							
A RW						RX	K da	ta f	rom	last	t tra	nsfe	er																			

#### 27.7.7 TXD

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	1 (	)
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	A /	A /	
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	
ID RW Field	Value ID	Va	lue			De	scri	ipti	on																								
A RW						T)	( da	ta f	or n	ext	trar	nsfe	r																				

# 27.7.8 FREQUENCY

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 0
ID (Field ID)	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A
Reset value	0	0	0	0	0	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field Value ID	Va	lue			De	scri	ptio	on																							
A RW					T٧	VI m	nast	er c	lock	free	que	ncy																			
K100	0x	0198	8000	00	10	00 kl	ops																								
K250	0x	400	0000	)	25	0 kl	ops																								
K400	0x	0668	8000	00	40	00 kl	ops																								

#### 27.7.9 ADDRESS

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	A	A	A	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID RW Field	Value ID	Va	lue			De	scr	ipti	on																							
A RW						T١	NI a	ddr	ess																							



# 28 Universal Asynchronous Receiver/Transmitter (UART)

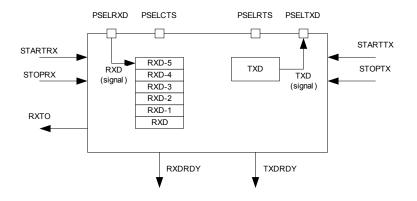


Figure 67 UART configuration

# 28.1 Functional description

The UART implements support for the following features:

- Full-duplex operation
- Automatic flow control
- Parity checking and generation for the 9<sup>th</sup> data bit

As illustrated in *Figure 67*, the UART uses the TXD and RXD registers directly to transmit and receive data. The UART uses one stop bit.

# 28.2 Pin configuration

The different signals RXD, CTS (Clear To Send, active low), RTS (Request To Send, active low), and TXD associated with the UART are mapped to physical pins according to the configuration specified in the PSELRXD, PSELCTS, PSELRTS, and PSELTXD registers respectively. If a value of 0xFFFFFFFF is specified in any of these registers, the associated UART signal will not be connected to any physical pin. The PSELRXD, PSELCTS, PSELRTS, and PSELTXD registers and their configurations are only used as long as the UART is enabled, and retained only for the duration the device is in ON mode. To secure correct signal levels on the pins by the UART when the system is in OFF mode, the pins must be configured in the GPIO peripheral as described in *Table 40*.

Only one peripheral can be assigned to drive a particular GPIO pin at a time. Failing to do so may result in unpredictable behavior.

UART pin	Direction	Output value
RXD	Input	Not applicable
CTS	Input	Not applicable
RTS	Output	1
TXD	Output	1

Table 40 GPIO configuration



The UART shares registers and other resources with other peripherals that have the same ID as the UART. Therefore, you must disable all peripherals that have the same ID as the UART before the UART can be configured and used. Disabling a peripheral that has the same ID as the UART will not reset any of the registers that are shared with the UART. It is therefore important to configure all relevant UART registers explicitly to ensure that it operates correctly.

See the Instantiation table in Section 4.2 "Instantiation" on page 13 for details on peripherals and their IDs.

# 28.4 Transmission

A UART transmission sequence is started by triggering the STARTTX task. Bytes are transmitted by writing to the TXD register. When a byte has been successfully transmitted the UART will generate a TXDRDY event after which a new byte can be written to the TXD register. A UART transmission sequence is stopped immediately by triggering the STOPTX task.

If flow control is enabled a transmission will be automatically suspended when CTS is deactivated and resumed when CTS is activated again, as illustrated in *Figure 68*. A byte that is in transmission when CTS is deactivated will be fully transmitted before the transmission is suspended.

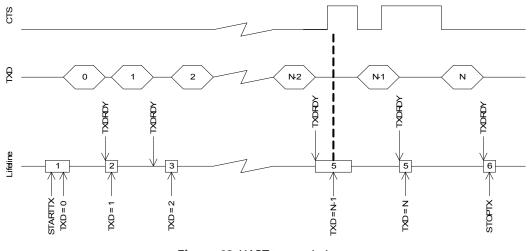


Figure 68 UART transmission

# 28.5 Reception

A UART reception sequence is started by triggering the STARTRX task. The UART receiver chain implements a FIFO capable of storing six incoming RXD bytes before data is overwritten. Bytes are extracted from this FIFO by reading the RXD register. When a byte is extracted from the FIFO a new byte pending in the FIFO will be moved to the RXD register. The UART will generate an RXDRDY event every time a new byte is moved to the RXD register.

When flow control is enabled, the UART will deactivate the RTS signal when there is only space for four more bytes in the receiver FIFO. The counterpart transmitter is therefore able to send up to four bytes after the RTS signal is deactivated before data is being overwritten. To prevent overwriting data in the FIFO, the counterpart UART transmitter must therefore make sure to stop transmitting data within four bytes after the RTS line is deactivated.



The RTS signal will first be activated again when the FIFO has been emptied, that is, when all bytes in the FIFO have been read by the CPU, see *Figure 69*.

The RTS signal will also be deactivated when the receiver is stopped through the STOPRX task as illustrated in *Figure 69*. The UART will be able to receive up to four bytes if they are sent immediately after the RTS signal has been deactivated. This is possible because the UART is, even after the STOPRX task is triggered, able to receive bytes for an extended period equal to the time it takes to send four bytes on the configured baud rate. The UART will generate a receiver timeout event (RXTO) when this period has elapsed.

To prevent loss of incoming data the RXD register must only be read once following every RXDRDY event.

If there is more than one byte available in the FIFO, RXDRDY event will be triggered immediately when the RXD register is read. To ensure that you have control of any new RXDRDY event arriving you must clear the EVENT before you read the RXD register.

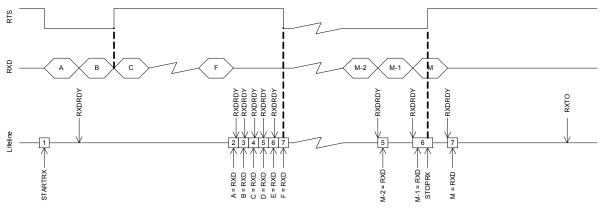


Figure 69 UART reception.

As indicated in occurrence 2 in *Figure 69*, the RXDRDY event associated with byte B is generated first after byte A has been extracted from RXD.

#### 28.6 Error conditions

An ERROR event, in the form of a framing error, will be generated if a valid stop bit is not detected in a frame. Another ERROR event, in the form of a break condition, will be generated if the RXD line is held active low for longer than the length of a data frame. Effectively, a framing error is always generated before a break condition occurs.

## 28.7 Using the UART without flow control

If flow control is not enabled the interface will behave as if the CTS and RTS lines are kept active all the time.

## 28.8 Parity configuration

When parity is enabled, the parity will be generated automatically from the even parity of TXD and RXD for transmission and reception respectively.



# 28.9 Registers

Register	Offset	Description
TASKS		
STARTRX	0x000	Start UART receiver
STOPRX	0x004	Stop UART receiver
STARTTX	0x008	Start UART transmitter
STOPTX	0x00C	Stop UART transmitter
EVENTS		
RXDRDY	0x108	Data received in RXD
TXDRDY	0x11C	Data sent from TXD
ERROR	0x124	Error detected
RXTO	0x144	Receiver timeout
REGISTERS		
INTENSET	0x304	Interrupt enable set register
INTENCLR	0x308	Interrupt enable clear register
ERRORSRC	0x480	Error source
ENABLE	0x500	Enable UART
PSELRTS	0x508	Pin select for RTS
PSELTXD	0x50C	Pin select for TXD
PSELCTS	0x510	Pin select for CTS
PSELRXD	0x514	Pin select for RXD
RXD	0x518	RXD register
TXD	0x51C	TXD register
BAUDRATE	0x524	Baud rate
CONFIG	0x56C	Configuration of parity and hardware flow control

Table 41 Register overview

## 28.9.1 RXD

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 3	2 1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	Α.	A	A	A /	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	<b>)</b> (	0 (
ID RW Field	Value ID	Va	lue			De	scri	ptic	on																							
A R						R۶	( da	ta re	ecei	ved	in p	orevi	ious	tra	nsfe	rs.																



#### 28.9.2 TXD

Bit number	31	30 2	29	28 27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	5	5	4	3	2	1 (	)
ID (Field ID)			•		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	A	Α.	A	A	A	A	A i	A
Reset value	0	0 0	D	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	D	0	0	0	0	0 0	)
ID RW Field	Value ID Valu	ue		De	scri	ptic	on																								
A W				T)	( da	ta to	o be	trai	nsfe	rrec	ł.																				

#### 28.9.3 ENABLE

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (Field ID)		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	A A	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) ()
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A RW						Er	nabl	e oi	<sup>r</sup> dis	able	e UA	ART																				
	DISABLE	0				Di	sab	le U	IAR	-																						
	ENABLE	4				Er	nabl	e U	ART																							

### 28.9.4 PSELRTS

Bit number	3	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9 8	B 7	76	5	4	3	2	1	0
ID (Field ID)	ŀ	۹.	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A	A	A	A	A	A	A
Reset value	1	l	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	11	1	1	1	1	1	1
ID RW Field	Value ID \	/alı	Je			De	scri	ptic	on																							
А						Piı	ի ու	ımb	er c	onf	iguı	atic	on fo	or U	ART	RTS	5 sig	nal														
RW		[0]	31]			Piı	n nu	ımb	er t	o ro	ute	the	UAI	RT R	TS s	sign	al to	C														
W		0xF	FFF	FFF	F	Di	scoi	nne	ct																							

#### 28.9.5 PSELTXD

Bit numl	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 1	0
ID (Field	ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	4 <i>4</i>	A
Reset va	lue		1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11	1
ID RW	Field	Value ID	Va	ue			De	scri	ptic	on																							
А							Pi	n nı	ımb	er c	onfi	igur	atio	n fo	r UA	RT	ТХС	) sig	nal														
RW					Pi	n nı	ımb	er t	o ro	ute	the	UAF	RT T	XD s	sign	al to	c																
w			0x	FFF	FFFF	F	Di	sco	nne	ct																							



## 28.9.6 PSELCTS

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 <sup>.</sup>	1 0
ID (Field ID)	· ·								A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A
Reset value	· · ·								1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 '	11
ID RW Field	Value ID	Va	lue			De	scri	ptio	on																							
А						Pir	n nu	ımb	oer c	onf	igur	atic	on fo	or U/	ART	СТ	5 sig	nal														
RW		[0	31]			Pir	n nu	ımb	oer t	o ro	ute	the	UA	RT C	TS s	sign	al to	C														
W		0x	FFF	FFFF	F	Di	sco	nne	ct																							

### 28.9.7 PSELRXD

Bit number		28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6 !	5	4	3 2	2 1	0				
ID (Field ID)			A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A A	A
Reset value	et value 1 1 1 1								1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1 1	1	1
ID RW Fie	d Va	alue ID		De	scri	ptic	n																										
А							Piı	n nu	ımb	er c	onfi	gur	atio	n fo	r UA	RT	RXD	sig	nal														
RW			[0.	31]			Piı	n nu	ımb	er to	o ro	ute	the	UAF	RT R	KD s	igna	al to	)														
W			0x	FFF	FFF	F	Di	scoi	nne	ct																							

### 28.9.8 ERRORSRC

Bit nur	nber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	76	5	4	3	2	10
ID (Fiel	d ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-				-	-	D	c	ΒA
Reset v	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	<b>D</b> (	0	0	0	0	0 0
ID RW	Field	Value ID	Va	lue			De	scri	ptic	n																						
A	OVERRUN									rror. is re		ved	whi	ile tl	he p	revi	ous	dat	a sti	ill lie	es in	RXI	D. (P	revio	ous	da	ta i	s los	st.)			
R			0				No	ot od	ccur	red																						
R			1				00	cur	red																							
W			1				Cle	ear																								
В	PARITY								erro racto		ith l	bad	par	ity i	s red	ceive	ed, i	f HV	V pa	arity	che	ck i	s en	able	d.							
С	FRAMING						A١	valio	d sto	rror op bi n rec	it is	not		ecte	ed o	n th	ie se	erial	dat	a in	outa	afte	r all	bits	in a	a cł	ara	icte	r			
D	BREAK		•				Th	e se	erial		a inp					-				-				rame y bit		'he	dat	a				



# 28.9.9 CONFIG

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	б	5	4	3 2	2 1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-		-	- E	3 B	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 1	0	0
ID RW Field Value ID	Val	lue			De	scri	ptie	on																							
A RW HWFC					Ha	ardv	vare	e flo	w co	ontr	ol																				
	0				Di	sab	led																								
	1				Er	abl	ed																								
B RW PARITY					Pa	rity																									
	0x	0			Ex	clud	de p	arit	y bit	t																					
	0x	7			In	cluc	le p	arity	/ bit																						

### 28.9.10 BAUDRATE

Bit number	31	30	29	28	27	26	25 2	4 2	3 22	2 21	20	19	18	17	16	15	14	13	12	11	10	9	8	76	5	4	3	2	1 (	)
ID (Field ID)	Α	A	A	A	A	A	A A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A	A	A	A	A	A /	١
Reset value	0	0	0	0	0	0	01	1	0	0	1	1	0	0	0	0	0	0	0	0	0	0	0 (	0	0	0	0	0	0 (	)
ID RW Field Value ID	Va	lue			De	scri	ption																							
A RW					Ва	ud-	rate																							
BAUD120	0 0:	x000	04F0	00	12	00 k	baud																							
BAUD240	0 0:	x000	09D0	00	24	00 I	baud																							
BAUD480	0 0	x00 <sup>-</sup>	13B0	00	48	00 ł	baud																							
BAUD960	0 0	x002	2750	00	96	00 ł	baud																							
BAUD144	00 00	x003	3B00	00	14	400	bauc	I																						
BAUD192	00 02	×004	1EA0	00	19	200	bauc	ł																						
BAUD288	00 02	x007	75F0	00	28	800	bauc	ł																						
BAUD384	00 02	x009	9D50	00	38	400	bauc	ł																						
BAUD576	00 0	x00I	EBFO	00	57	600	bauc	ł																						
BAUD768	00 02	x013	3A90	00	76	800	bauc	ł																						
BAUD115	200 0	x01l	D7E0	00	11	520	0 bau	ıd																						
BAUD230	400 02	x03/	AFB0	00	23	040	0 bau	ıd																						
BAUD250	000 0:	x04(	0000	00	25	000	0 bau	ıd																						
BAUD460	800 0:	x07	5F70	00	46	080	0 bau	ıd																						
BAUD921	600 02	X0E	BEDF	A4	92	160	0 bau	ıd																						
BAUD1M	0:	x100	0000	00	1 r	neg	abau	d																						



# 29 Quadrature Decoder (QDEC)

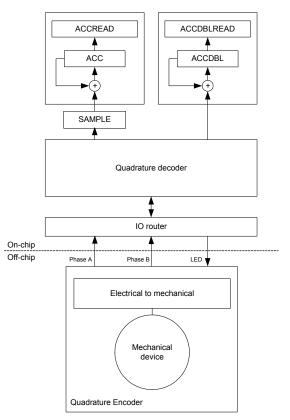


Figure 70 Quadrature decoder configuration

# 29.1 Functional description

The Quadrature Decoder (QDEC) can be used for decoding the output of an off-chip quadrature encoder. The QDEC provides the following:

- Decoding of digital waveform from off-chip quadrature encoder.
- Sample accumulation eliminating hard real-time requirements to be enforced on application.
- Optional input debounce filters.
- Optional LED output signal for optical encoders.

#### 29.1.1 Pin configuration

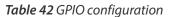
The different signals: Phase A, Phase B, and LED, are mapped to physical pins according to the configuration specified in the PSELA, PSELB, and PSELLED registers respectively. If a value of 0xFFFFFFFF is specified in any of these registers, the associated signal will not be connected to any physical pin. The PSELA, PSELB, and PSELLED registers and their configurations are only used as long as the QDEC is enabled, and retained only as long as the device is in ON mode.

To secure correct behavior in the QDEC, the pins used by the QDEC must be configured in the GPIO peripheral as described in *Table 42* on page 168 prior to enabling the QDEC. This configuration must be retained in the GPIO for the selected IOs as long as the QDEC is enabled.



Only one peripheral can be assigned to drive a particular GPIO pin at a time, failing to do so may result in unpredictable behavior.

QDEC signal	QDEC pin	Direction	Output value
Phase A	As specified in PSELA	Input	Not applicable
Phase B	As specified in PSELB	Input	Not applicable
LED	As specified in PSELLED	Input	Not applicable



#### 29.1.2 Sampling and decoding

The off-chip quadrature encoder is an incremental motion encoder outputting two waveforms; phase A and phase B. The two output waveforms are always 90 degrees out of phase, meaning that one always changes level before the other. The direction of movement is indicated by which of these two waveforms that changes level first. Invalid transitions may occur, that is when the two waveforms switch simultaneously. This may occur if the wheel rotates too fast relative to the sample rate set for the decoder.

The QDEC decodes the output from the off-chip encoder by sampling the QDEC phase input pins (A and B) at a fixed rate as specified in the SAMPLEPER register.

When started, the decoder continuously samples the two input waveforms and decodes these by comparing the current sample pair (n) with the previous sample pair (n-1). The decoding of the sample pairs is described in *Table 43*.

Previou (n - 1)	is sample pair	Curr pair		SAMPLE register	ACC operation	ACCDBL operation	Description
A	В	Α	В	register	operation	operation	
0	0	0	0	0	No change	No change	No movement
		0	1	1	Increment	No change	Movement in positive direction
		1	0	-1	Decrement	No change	Movement in negative direction
		1	1	2	No change	Increment	Error: Double transition
0	1	0	0	-1	Decrement	No change	Movement in negative direction
		0	1	0	No change	No change	No movement
		1	0	2	No change	Increment	Error: Double transition
		1	1	1	Increment	No change	Movement in positive direction
1	0	0	0	1	Increment	No change	Movement in positive direction
		0	1	2	No change	Increment	Error: Double transition
		1	0	0	No change	No change	No movement
		1	1	-1	Decrement	No change	Movement in negative direction
1	1	0	0	2	No change	Increment	Error: Double transition
		0	1	-1	Decrement	No change	Movement in negative direction
		1	0	1	Increment	No change	Movement in positive direction
		1	1	0	No change	No change	No movement

 Table 43
 Quadrature decoder input decoding



## 29.1.3 LED output

The LED output follows the sample period and the LED is switched on a given period prior to sampling and switched off immediately after the inputs are sampled. The period the LED is switched on prior to sampling is given in the LEDPRE register.

The LED output pin polarity is specified in the LEDPOL register.

For using off-chip mechanical encoders not requiring a LED, the LED output can be disabled by writing 0xFFFFFFF to the PSELLED register. In this case the QDEC will not acquire access to a LED output pin and the pin can be used for other purposes by the CPU.

#### 29.1.4 Debounce filters

Each of the two phase inputs have digital debounce filters. When enabled, the filter inputs are sampled at a fixed 1 MHz frequency during the entire sample period (which is specified in the SAMPLEPER register), and the filters require all of the samples within this sample period to equal before the input signal is accepted and transferred to the output of the filter.

As a result, only input signal with a steady state longer than twice the period specified in SAMPLEPER are guaranteed to pass through the filter, and any signal with a steady state shorter than SAMPLEPER will always be suppressed by the filter. (This is assumed that the frequency during the debounce period never exceeds 500 kHz (as required by the Nyquist theorem when using a 1 MHz sample frequency).

**Note:** The LED will always be ON when the debounce filters are enabled, as the inputs in this case will be sampled continuously.

#### 29.1.5 Accumulators

The quadrature decoder contains two the accumulator registers ACC and ACCDBL that accumulates valid motion sample values and the number of detected invalid samples (double transitions), respectively.

The ACC register will accumulate all valid values (1/-1) written to the SAMPLE register. This can be useful for preventing hard real-time requirements from being enforced on the application. When using the ACC register the application does not need to read every single sample from the SAMPLE register, but can instead fetch the ACC register whenever it fits the application. The ACC register will always hold the relative movement of the external mechanical device since the previous clearing of the ACC register. Sample values indicating a double transition (2) will not be accumulated in the ACC register.

An ACCOF event will be generated if the ACC receives a SAMPLE value that would cause the register to overflow or underflow. Any SAMPLE value that would cause the an ACC overflow or underflow will be discarded, but any samples not causing the ACC to overflow or underflow will still be accepted.

The accumulator ACCDBL accumulates the number of detected double transitions since the previous clearing of the ACCDBL register.

The ACC and ACCDBL registers can be cleared by the READCLRACC and subsequently read using the ACCREAD and ACCDBLREAD registers.



#### 29.1.6 Output/input pins

The QDEC uses a 3 pin interface to the off-chip quadrature encoder.

These pins will be acquired when the QDEC is enabled in the ENABLE register. The pins acquired by the QDEC cannot be written by the CPU, but they can still be read by the CPU.

The pin numbers to be used for the QDEC are selected using the PSELn registers.

# 29.2 Registers

Register	Offset	Description
TASKS		
START	0x000	Task starting the quadrature decoder. When started, the SAMPLE register will be continuously updated at the rate given in the SAMPLEPER register.
STOP	0x004	Task stopping the quadrature decoder.
READCLRACC	0x008	Task transferring the content of ACC to ACCREAD and the content of ACCDBL to ACCDBLREAD, and then clearing the ACC and ACCDBL registers. These read-and-clear operations will be done automatically.
EVENTS		
SAMPLERDY	0x100	Event being generated for every new sample value written to the SAMPLE register.
REPORTRDY	0X104	Event being generated when REPORTPER number of samples has been accumulated in the ACC register and the content of the ACC register does not equal 0. (Thus, this event is only generated if a motion is detected since the previous clearing of the ACC register).
ACCOF	0X108	ACC or ACCDBL register overflow.
REGISTERS		
SHORTS	0x200	Shortcut register.
INTENSET	0x304	Interrupt enable set register.
INTENCLR	0x308	Interrupt enable clear register.
ENABLE	0x500	ADC enable.
LEDPOL	0x504	ADC configuration.
SAMPLEPER	0x508	ADC conversion result.
SAMPLE	0X50C	Motion sample value.
REPORTPER	0X510	Number of samples to be taken before a REPORTRDY event is generated.
ACC	0X514	Register accumulating the valid transitions.
ACCREAD	0X518	Snapshot of the ACC register updated by the READCLRACC task.
PSELLED	0X51C	GPIO pin number to be used as LED output.
PSELA	0X520	GPIO pin number to be used as Phase A input.
PSELB	0X524	GPIO pin number to be used as Phase B input.
DBFEN	0X528	Enable input debounce filters.
LEDPRE	0X540	Time period the LED is switched ON prior to sampling.
ACCDBL	0X544	Register accumulating the number of detected double transitions.
ACCDBLREAD	0X548	Snapshot of the ACCDBL.



# 29.2.1 Shorts

Bit nun	nber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6	5	4	3 2	2 1	C
ID (Fiel	d ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-		• •	В	A
Reset v	alue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0 0	0
ID RW	Field	Value ID	Va	lue			De	scri	ipti	on																							
A RW	REPORTRDY_ READCLRACC	Disable Enable	0 1				Sł	ort	REF	POR	TRD	Y ev	/ent	t to	REA	DCI	RA	CC t	ask														
B RW	SAMPLERDY_ STOP	Disable Enable	0 1				Sł	ort	SAI	MPI	ERC	)Y e	vent	t to	STC	)P ta	isk.																

## 29.2.2 ENABLE

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6 !	5 4	4 3	3 2	2 1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-				• -	-	-	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	0 0	) (	0	) 0
ID RW Field Value ID	Val	ue			De	scri	pti	on																							
A RW VAL					W			•					der. ure									•									
DISABLE	0																														
ENABLE	1																														

# 29.2.3 LEDPOL

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 3	32	2 1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-			-		-	-	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0
ID RW Field Value ID	Va	lue			De	scri	pti	on																							
A RW					LE	Do	utp	ut p	ola	rity.																					
ACTIVELOW	0				LE	D a	ctiv	e or	n ou	tpu	t pir	n lov	Ν.																		
ACTIVEHIGH	1				LE	D a	ctiv	e or	n ou	tpu	t pir	n hig	gh.																		



# 29.2.4 SAMPLEPER

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4 3	3 2	1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		ŀ	A	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	0	0
ID RW Field Value ID	Va	lue			De	scri	ptic	on																							
A RW VAL					Sa	mp	le p	erio	d. T	he S	AM	PLE	regi	ster	wil	l be	upc	late	d fo	r ev	ery i	new	sa	mp	ole.						
128_US	0				12	.8 μ	5																								
256_US	1				25	6 μ	5																								
512_US	2				51	2 μ	5																								
1024_US	3				10	24	JS																								
2048_US	4				20	48	JS																								
4096_US	5				40	96 J	JS																								
8192_US	6				81	92 J	JS																								
16384_US	7				16	384	μs																								

### 29.2.5 SAMPLE

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	6 !	5 4	13	2	1	0
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A /	A /	A A		A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	) (	0	0	0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A R		(	12	)		La	ast n	noti	ons	sam	ple.																					
						m	otic	on.				•	eme dou					he s	ign	give	es th	ne d	irec	tio	n o	f th	e					



# 29.2.6 REPORTPER

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	87	76	5	4	3	2	1 (	)
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-	-	-	-	A	A	A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) 0	0	0	0	0	0	)
ID RW Field Value ID	Va	lue	!		De	scri	ptic	on																							
A RW VAL					th Th Rf W Rf Sf	pecif ne re PUS PUS PUS Pis t	EPO epor = SI e is th	RTR rt pe P * F ne re sam	EPY Erio RP Epo Ple	eve d in rt p	nt c µs i erio iod	an k s gi d in in [µ	· pe g ven [μs/ us/sa	ene as: /rep amp	oort	ed.	cifie	ed ir	n SA	.MP	LEP	ER		iter	be	fore	5				
10_SMPL	0				10	) sar	nple	es/r	еро	rt																					
40_SMPL	1				40	) sar	nple	es/r	epo	rt																					
80_SMPL	2				80	) sar	nple	es/r	еро	rt																					
120_SMPL	3				12	20 sa	amp	oles/	/rep	ort																					
160_SMPL	4				16	50 sa	amp	oles/	/rep	ort																					
200_SMPL	5				20	)0 sa	amp	les/	/rep	ort																					
240_SMPL	6				24	10 sa	amp	oles/	/rep	ort																					
280_SMPL	7				28	30 sa	amp	oles/	/rep	ort																					

### 29.2.7 ACC

Bit number	31 30 29 28	7 26 25 24 23 22 21 20 19 18 17 1	6 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)	A A A A	A A A A A A A A A A A A A A A A A A A	
Reset value	0 0 0 0	0 0 0 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
ID RW Field	Value ID Value	escription	
A R	(-10241023)	Register accumulating all valid samples (r RENC (in the SAMPLE register). Double tra accumulated in this register. The value is a 32 bit 2's complement value f a sample that would cause this register ample will be ignored and an overflow e The ACC register is cleared by triggering t	e. e. to overflow or underflow is received, the vent (ACCOF) will be generated.



# 29.2.8 ACCREAD

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1
ID (Field ID)		A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
AR		(-1	024	10	23)	Sr	naps	shot	t of t	the <i>i</i>	ACC	reg	iste	r.																		

#### 29.2.9 PSELLED

Bit number	31	1 30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	)
ID (Field ID)	Α	A	A	A	A	A	A	A	A	Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A i	A	١
Reset value	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	I
ID RW Field	Value ID Va	alue			De	scri	ptio	on																								
A RW	((	031)				PIO   sabl				r to ut.	be	use	d as	s LEI	Ο οι	ιtpι	ıt. V	/riti	ng t	the	valu	ıe 0	xFF	FF	FFI	FF۱	wil	I				

## 29.2.10 PSELA

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)		A	A	A	A	A	A	Α	Α	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A	A A
Reset value		1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	11
ID RW Field	Value ID	Va	lue			De	scri	ptic	on																							

# 29.2.11 PSELB

Bit number	31 30 29 28	3 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
ID (Field ID)	A A A A	<b>A A A A A A A A A A</b>
Reset value	1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
ID RW Field Value ID	Value	Description
A RW	(031)	GPIO pin number to be used as Phase B input. Writing the value 0xFFFFFFFF will disable this input.

## 29.2.12 LEDPRE

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	5 5	54	3	2	1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A	۹ <i>۱</i>	۹ <i>4</i>	A	A	A	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	) (	) (	0	0	0	0
ID RW Field Va	lue ID	Val	ue			De	scri	pti	on																							
A RW		(0.	.511	)		Pe	eriod	d in	µs t	he l	ED	is s	witc	hed	on	prio	or to	o sai	mpl	ing.												



#### 29.2.13 DBFEN

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9 8	B 7	7 6	5 5	5 4	13	2	1	0
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		• •	• •		-	-	-	-	A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 (	<b>b</b> (	) (	) (	) ()	0	0	0
ID RW Field V	/alue ID	Val	ue			De	scri	ipti	on																							
A RW						Er	nabl	e in	nput	del	oou	nce	filte	ers.																		
		0				Di	isab	le																								
		1				Er	nabl	e																								

### 29.2.14 ACCDBL

Bit number		31	30	29	28	27	26	25	5 24	23	22	2 21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 (	55	4	3	2	1	C
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-			-	A	A	F	\ A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	) (	0	0	0	0	) (
D RW Field	Value ID	Val	lue			De	escri	pti	on																							
A R		(0.	15)			(S W ill A aı	egist AMI /hen legal n ov re de	PLE th tra erf	is regansit	). gist ion eve afte	er h s wi nt ( er th	ill sto ACC	eacł op. :OF) naxi	ned will mur	its r be n or	nax ger r mi	imu nera nim	m v ted um	ʻalu if ai val	e th ny d ue v	e ac ouk vas	cum ole c	nula or ille	tior ega	n of	fdc	bub					

# 29.2.15 ACCDBLREAD

Bit number	3	1 30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 0	
ID (Field ID)	-	-	-	-	-	-	-		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	A /	A A	
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	
ID RW Field	Value ID V	alue			De	scri	ptio	n																								
A R	(	(015)	)		Sn	aps	hot	of th	ne A	CCI	OBL	regi	ister	r.																		
					Th	is fi	eld i	s up	date	ed ۱	whe	n th	e RE	EAD	CLR	ACC	2 tas	sk is	trig	gere	ed.											



# 30 Analog to Digital Converter (ADC)

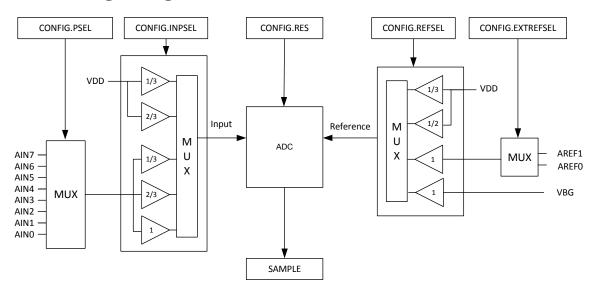


Figure 71 Analog to Digital converter

# 30.1 Functional description

#### 30.1.1 Configuration

All parameters such as input selection, reference selection, resolution, pre-scaling etc. are configured using the CONFIG register.

Note: It is not allowed to configure the ADC during an on-going ADC conversion (ADC busy).

#### 30.1.2 Usage

An ADC conversion is started by using the START task, either by writing the task register directly from the CPU or by triggering the task through the PPI.

During sampling the ADC will enter a busy state. The ADC busy/ready state can be monitored via the BUSY register.

When the ADC conversion is completed, an END event will be generated and the result of the conversion can be read from the RESULT register.

When the ADC conversion is completed, the ADC analog electronics power down to save power.

#### 30.1.3 One-shot / continuous operation

The ADC itself only supports one-shot operation, this means every single conversion has to be explicitly started using the START task.

However, continuous ADC operation can be achieved by continuously triggering the START task from, for example, a timer through the PPI.



#### 30.1.4 Pin configuration

The user can use the PSEL register to select one of the analog input pins, AIN0 through AIN7, as input for the ADC. See the device product specification for more information about which analog pins are available on a particular device. The selected analog pin will be acquired by the ADC when it is enabled through the ENABLE register, see *Chapter 13 "GPIO"* on page 58 for more information on how analog pins are selected.

#### 30.1.5 Shared resources

The ADC shares registers and other resources with peripherals that have the same ID as the ADC. The user must therefore disable all peripherals that have the same ID as the ADC before the ADC can be configured and used. The ADC is using the same analog pins as the LPCOMP. The LPCOMP must therefore be disabled before the ADC can be enabled. It is important to configure all relevant ADC registers explicitly to secure that it operates correctly.

The Instantiation table in *Section 4.2 "Instantiation"* on page 13 shows which peripherals have the same ID as the ADC.

Register	Offset	Description
TASKS		
START	0x000	Start a new ADC conversion.
STOP	0x004	Stop ADC.
EVENTS		
END	0x100	An ADC conversion is completed.
REGISTERS		
INTENSET	0x304	Interrupt enable set register.
INTENCLR	0x308	Interrupt enable clear register.
BUSY	0x400	ADC busy (conversion in progress).
ENABLE	0x500	Enable ADC. When enabled, the ADC will acquire access to the analog input pins specified in the CONFIG register.
CONFIG	0x504	ADC configuration.
RESULT	0x508	Result of the previous ADC conversion.

# 30.2 Registers

 Table 44
 Register overview

#### 30.2.1 BUSY

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5 4	4 3	32	2 1	0
ID (Field	ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		-		-	A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0 (	0 0	0	0
ID RW	Field	Value ID	Va	lue			De	scri	ptic	on																							
A R																																	
			0				A	DC is	s rea	ady.	No	ong	oing	g coi	nve	rsio	n.																
			1				A	DC is	s bu	sy. (	Conv	/ersi	ion i	in pı	rogr	ess.																	



## 30.2.2 ENABLE

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3 2	21	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-		• 4	A A
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	) (	) 0
ID RW Field Value ID	Va	lue			De	escr	ipti	on																							
A RW VAL																															
	0				А	DC	disa	bleo	d.																						
					۸		~ <b>~</b> ~	blec	J																						



# 30.2.3 CONFIG

Bit n	umbe	r		31 30	29	28	27	26 2	25 2	42	23 22	21	20	19	18	17	16	15	14	13	12	11	10	9	B 7	6	54	3 2	10
ID (F	ield IC	<b>D</b> )			-	-	-		-	-	-	-	-	-	-	Е	Е	D	D	D	D	D	D	DI	D -	с	СВ	BE	3 A A
Rese	t valu	e		0 0	0	0	0	0 0	) ()	0	) ()	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0	0	0 0	0 0	000
ID	RW	Field	Value ID	Value			De	scrip	tion																				
А	RW	RES					A	DC res	solut	ion																			
			8_BIT	0			8	bit																					
			9_BIT	1			9	bit																					
			10_BIT	2			10	) bit																					
В	RW	INPSEL					A	DC in	out s	eleo	ction.																		
			AIN_NO_ PS	0			Ar	nalog	inpu	ıt pi	in spe	cifie	d by	CC	NFI	G.PS	SEL	with	no	pre	scal	ing.							
			AIN_ 2_3_PS	1			Ar	nalog	inpu	ıt pi	in spe	cifie	d by	CC	NFI	G. P	SEL	witł	ר 2/:	3 pr	esca	aling	g.						
			AIN_ 1_3_PS	2			Ar	nalog	inpu	ıt pi	in spe	cifie	d by	CC	NFI	G. P	SEL	witł	n 1/:	3 pr	esca	aling	g.						
			VDD 2_3_PS	5			V	DD wi	th 2/	′3 p	rescal	ing.																	
			VDD 1_3_PS	6			V	DD wi	th 1/	′3 p	rescal	ing.																	
С	RW	REFSEL					A	DC ref	feren	ice :	select	ion																	
			VBG	0			Us	se inte	ernal	1.2	2 V bai	nd g	ap re	efer	enc	e.													
			EXT	1			Us	se ext	erna	l ref	ferenc	e sp	ecifi	ed	by C	ON	CFIC	G EX	TRE	FSE	L.								
			VDD_1_2 _PS	2				se VD 6V).	D wi	th 1	l/2 pre	escal	ing.	(Or	nly a	ppli	cab	le w	hen	ND	D is	in t	he	rang	je 1	.7V	-		
			VDD_1_3 _PS	3				se VD 6V).	D wi	th 1	/3 pre	escal	ing.	(Or	nly a	ppli	cab	le w	hen	ı VD	D is	in t	he	rang	je 2	2.5V	_		
D	RW	PSEL					Se	elect p	oin to	b be	e used	as A	NDC i	inp	ut p	in.													
			DISABLE	0			Ar	nalog	inpu	ıt pi	ins dis	able	ed.																
			AIN0	1			Us	se AlN	10 as	ana	alog ir	nput	•																
			AIN1	2			Us	se AlN	l1 as	ana	alog ir	nput																	
			AIN2	4			Us	se AIN	l2 as	ana	alog ir	nput	•																
			AIN3	8			Us	se AlN	I3 as	ana	alog ir	nput	•																
			AIN4	16			Us	se AlN	l4 as	ana	alog ir	nput	•																
			AIN5	32			Us	se AlN	l5 as	ana	alog ir	nput	•																
			AIN6	64			Us	se AlN	l6 as	ana	alog ir	nput	•																
			AIN7	128			Us	se AlN	l7 as	ana	alog ir	nput																	
Е	RW	EXTREFSEL					Ex	terna	l refe	erer	nce pi	n sel	ectio	on.															
			NONE	0			Ar	nalog	refe	ren	ce inp	uts o	disak	olec	ł.														
			AREF0	1			Us	se ARI	EF0 a	is ai	nalog	refe	renc	e.															
			AREF1	2			Us	se ARI	EF1 a	is ai	nalog	refe	renc	e.															



#### 30.2.4 RESULT

Bit num	ber		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7 6	6 !	5 4	4 3	3 2	2 1	0
ID (Field	ID (Field ID)					-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A	Α.	A /	A /	A /	A /	۹ <i>۱</i>	A A	A A
Reset va	lue		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	0 0	0 0	) (	) (	) 0
ID RW	Field	Value ID	Va	lue			De	scri	ptic	on																							
A R			A[ th sig	DC c e sa gnal	onv mpl , the	the p versi led a e res ne va	on. <sup>.</sup> anal ult v	The og i /alu	resi inpu e wi	ult v it sig Il be	alue gnal e set	e is r is e to t	relat qua the i	ive l to max	to th or g imu	ne s rea m (l	elec ter t limit	ted han ted l	AD( the cy th	C ref AD ne se	fere C re elec	ence efer cteo	e in reno d Al	ipu ce DC	it. II : bit	f							



# 31 Low Power Comparator (LPCOMP)

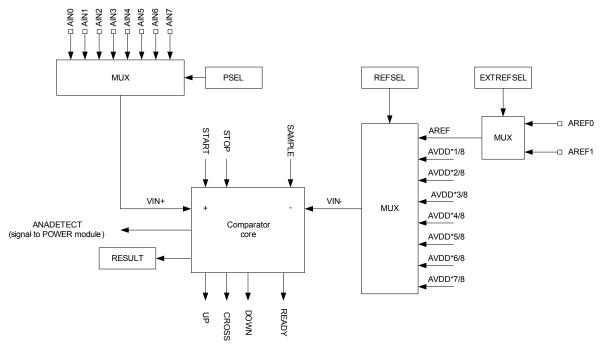


Figure 72 Low power comparator

# 31.1 Functional description

The low power comparator (LPCOMP) compares an input voltage (VIN+), which comes from an analog input pin selected through the PSEL register against a reference voltage (VIN-) selected through the REFSEL and EXTREFSEL registers.

The PSEL, REFSEL, and EXTREFSEL registers must be configured before the LPCOMP is enabled through the ENABLE register.

The LPCOMP is started by triggering the START task. After a start-up time of t<sub>LPCOMPSTARTUP</sub> the LPCOMP will generate a READY event to indicate that the comparator is ready to use and the output of the LPCOMP is correct. After the READY event is generated, the LPCOMP will generate events every time VIN+ crosses VIN-, that is, every time VIN+ rises above VIN- (upward crossing) and every time VIN+ falls below VIN- (downward crossing), see *Figure 72*. The LPCOMP is stopped by triggering the STOP task.

An upward crossing will generate an UP event and a downward crossing will generate a DOWN event. The CROSS event will be generated every time there is a crossing, independent of direction.

LPCOMP will be operational in both System ON and System OFF mode when it is enabled through the ENABLE register, see *Chapter 11 "Power management (POWER)"* on page 40 for more information about power modes. All LPCOMP registers including the ENABLE register are classified as retained registers when the LPCOMP is enabled. However, when the device wakes up from System OFF, all LPCOMP registers will be reset.

The LPCOMP can wake up the system from System OFF by asserting the ANADETECT signal. The ANADETECT signal can be derived from any of the event sources that generate the UP, DOWN, and CROSS events. In case of wakeup from System OFF, no events will be generated, only the ANADETECT signal. See the ANADETECT register *Section 31.4.7 "ANADETECT"* on page 185 for more information on how to



configure the ANADETECT signal.

The immediate value of the LPCOMP can be sampled to the RESULT register by triggering the SAMPLE task.

See the *Section 11.2.1 "RESETREAS"* on page 50 on the RESETREAS register in the POWER module for more information on how to detect a wakeup from LPCOMP.

## 31.2 Pin configuration

You can use the PSEL register to select one of the analog input pins, AIN0 through AIN7, as analog input pin for the LPCOMP, see *Figure 72* on page 181. Similarly, you can use the EXTREFSEL register to select one of the analog reference input pins, AREF0 and AREF1, as input for AREF in case AREF is selected in REFSEL. The selected analog pins will be acquired by the LPCOMP when it is enabled through the ENABLE register. See the product specification for more information about which analog pins are available on a particular device.

## 31.3 Shared resources

The LPCOMP shares registers and other resources with peripherals that have the same ID as the LPCOMP. You must disable all peripherals that have the same ID as the LPCOMP before the LPCOMP can be configured and used. Disabling a peripheral that has the same ID as the LPCOMP will not reset any of the registers that are shared with the LPCOMP. Therefore, it is important to configure all relevant LPCOMP registers explicitly to secure that it operates correctly.

See the Instantiation table in Section 4.2 "Instantiation" on page 13 for details on peripherals and their IDs.

**Note:** The LPCOMP is using the same analog pins as the ADC. The ADC must be disabled before the LPCOMP can be enabled.



# 31.4 Registers

Register	Offset	Description
TASKS		
START	0x000	Start comparator.
STOP	0x004	Stop comparator.
SAMPLE	0x008	Sample comparator value.
EVENTS		
READY	0x100	LPCOMP is ready and output is valid.
DOWN	0x104	Downward crossing.
UP	0x108	Upward crossing.
CROSS	0x10C	Downward or upward crossing.
REGISTERS		
SHORTS	0x200	Shortcuts for LPCOMP.
INTENSET	0x304	Interrupt enable set register.
INTENCLR	0x308	Interrupt enable clear register.
RESULT	0x400	Compare result.
ENABLE	0x500	Enable LPCOMP.
PSEL	0x504	Input pin select.
REFSEL	0x508	Reference select.
EXTREFSEL	0x50C	External reference select.
ANADETECT	0x520	Analog detect configuration.

Table 45 Register overview

## 31.4.1 SHORTS

Bit number	31	30 2	29 2	28 2	7 26	25	24	23 2	22 2	1 2	20 19	18	17	16	15	14	13	12	11	10	9	87	6	5	4	3 2	2 '	10
ID (Field ID)	-		-	-	-	-	-	-				-	-	-	-	-	-	-	-	-	-		-	-	E	DO	C 1	ВА
Reset value	0	0 0	0	) 0	0	0	0	0	0 (	0 (	0 0	0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0 0	0 (	0 0
ID RW Field Value ID	Val	ue		D	Descr	iptic	n																					
A RW READY_SAMPLE	0 1			I	Enab Disat Enab	ole sh	orto	cut.	shoi	rtcu	ıt bet	wee	n RE	EAD	Y ev	/ent	and	d SA	MP	LE t	ask:	•						
B RW READY_STOP	0 1			I	Enab Disat Enab	ole sh	orto	cut.	shoi	rtcu	ıt bet	wee	n RE	EAD	Y ev	/ent	and	d ST	ΌP	tasł	κ.							
C RW DOWN_STOP	0 1			I	Enab Disat Enab	ole sh	orto	cut.	shoi	rtcu	ıt bet	wee	n D	OWI	N ev	vent	and	d ST	ΌP	tasl	κ.							
D RW UP_STOP	0 1			I	Enab Disat Enab	ole sh	orto	cut.	shoi	rtcu	ıt bet	wee	n Ul	P ev	ent	and	I ST	OP	task	•								
E RW CROSS_STOP	0 1			I	Enab Disat Enab	ole sh	orto	cut.	shoi	rtcu	ıt bet	wee	n Cl	ROS	S ev	/ent	and	d ST	ΌP	task	κ.							



#### 31.4.2 ENABLE

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	A /
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID RW Field Value ID	Val	lue			De	escri	ipti	on																							

#### 31.4.3 RESULT

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	10
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- A
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID RW Field	Value ID	Va	lue			De	escri	ipti	on																							
AR						R	esul	t of	last	con	npa	re. (	Deci	sior	n po	int	SAN	ИPL	E ta	sk.												

## 31.4.4 PSEL

Bit number	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6 !	5 4	43	3 2	: 1	0
ID (Field ID)	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-				A	A	AA
Reset value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0	0	) ()	0	0 (
ID RW Field Value ID	Va	lue			De	scr	ipti	on																							
A RW					A	nalo	og p	in s	elec	t.																					
AINO	0				A	IN0	sele	ecte	d as	ana	alog	, inp	ut.																		
AIN1	1				A	IN1	sele	ecte	d as	ana	alog	, inp	ut.																		
AIN2	2				A	IN2	sele	ecte	d as	ana	alog	, inp	ut.																		
AIN3	3				A	IN3	sele	ecte	d as	ana	log	ı inp	ut.																		
AIN4	4				A	IN4	sele	ecte	d as	ana	alog	, inp	ut.																		
AIN5	5				A	IN5	sele	ecte	d as	ana	log	ı inp	ut.																		
AIN6	6				A	IN6	sele	ecte	d as	ana	alog	, inp	ut.																		
AIN7	7				A	IN7	sele	ecte	d as	ana	alog	ı inp	ut.																		



#### 31.4.5 REFSEL

Bit number	31 30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8 3	7 6	5 5	4	3	2	10
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-			-	-	-	A	A A
Reset value	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	) (	) (	0	0	0	0 0
ID RW Field Value ID	Value			Des	scri	ptic	on																						
A RW				Re	fere	ence	e sel	ect.																					
AVDD_1_8	0			AV	DD	* 1,	/8 se	elec	ted	as	refe	reno	ce.																
AVDD_2_8	1			AV	DD	* 2,	/8 se	elec	ted	as	refe	ren	ce.																
AVDD_3_8	2			AV	DD	* 3,	/8 se	elec	ted	as	refe	reno	ce.																
AVDD_4_8	3			AV	DD	* 4,	/8 se	elec	ted	as	refe	rene	ce.																
AVDD_5_8	4			AV	DD	* 5,	/8 se	elec	ted	as	refe	reno	ce.																
AVDD_6_8	5			AV	DD	* 6,	/8 se	elec	ted	as	refe	rene	ce.																
AVDD_7_8	6			AV	DD	* 7,	/8 se	elec	ted	as	refe	rene	ce.																
AREF	7			Ext	terr	nal a	nal	og r	efei	ren	ce s	elec	ted	•															

#### 31.4.6 EXTREFSEL

Bit number		31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1 (
ID (Field ID)		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- 4
Reset value		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 0
ID RW Field	Value ID	Va	lue			De	scri	pti	on																							
A RW						Ex	terr	nal a	anal	oa r	refer	enc	ce se	elec	t.																	

### 31.4.7 ANADETECT

Bit number		31	30	29	28	27	26	25	24	23	22	21 2	) 19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2 <sup>.</sup>	10
ID (Field ID)		-	-	-	-	-	-	-	-	-	-		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	- /	A A
Reset value		0	0	0	0	0	0	0	0	0	0	0 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0 (	0 0
ID RW Field	Value ID	Va	lue			De	scri	ptic	on																						



# Appendix A: SoftDevice architecture

*Figure 73* is a block diagram of the nRF51 series software architecture including the standard ARM<sup>®</sup> CMSIS interface for nRF51 hardware, profile and application code, application specific peripheral drivers, and a firmware module identified as a SoftDevice.

A SoftDevice is precompiled and linked binary software implementing a wireless protocol. While it is software, application developers have minimal compile-time dependence on its features. The unique hardware and software supported framework, in which it executes, provides run-time isolation and determinism in its behavior. These characteristics make the interface comparable to a hardware peripheral abstraction with a functional, programmatic interface.

The SoftDevice Application Program Interface (API) is available to applications as a high-level programming language interface, for example, a C header file.

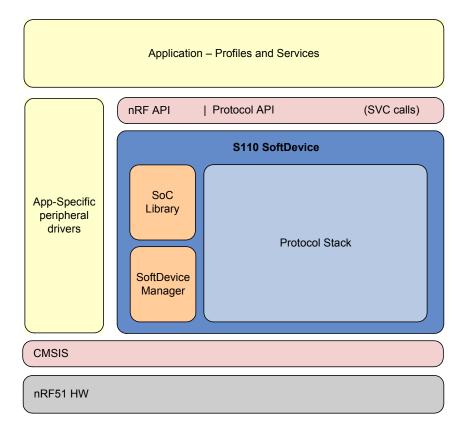


Figure 73 Software architecture block diagram

A SoftDevice consists of three main components:

- 1. SoC Library API for shared hardware resource management (application coexistence).
- 2. SoftDevice manager SoftDevice management API (enabling/disabling the SoftDevice, etc.).
- 3. Protocol stack Implementation of protocol stack and API.

When the SoftDevice is disabled, only the SoftDevice Manager API is available for the application. For more information about enabling/disabling the SoftDevice, see the Softdevice enable and disable section on page 194.



## SoC library

The SoC library provides functions for accessing shared hardware resources. The features of this library will vary between implementations of SoftDevices so detailed descriptions of the SoC library API are made available with the Software Development Kits (SDK) specific to each SoftDevice. The following is a summary of common components in the library.

Component	Description
NVIC	Wrapper functions for the CMSIS NVIC functions provided by ARM <sup>®</sup> . <b>Note:</b> To ensure reliable usage of the SoftDevice you must use the wrapper functions when the SoftDevice is enabled.
MUTEX	Disabling interrupts shall not be done while the SoftDevice is enabled. Mutex functions have been implemented to provide safe regions.
RAND	Random number generator - hardware sharing between SoftDevice and application.
POWER	Power management - Functions for power management.
CLOCK	Clock management – Functions for managing clock sources.
PPI	Safe PPI access to dedicated Application PPI channels.
PWR_MNG	Power management support (not a full implementation) for the application.

#### **SoftDevice Manager**

The SoftDevice Manager (SDM) API implements functions for controlling the state of the SoftDevice enabled/disabled. When enabled, the SDM configures low frequency clock (LFCLK) source, interrupt management and the embedded protocol stack.

Detailed documentation of the SDM API is made available with the Software Development Kits (SDK) specific to each SoftDevice.

### **Protocol stack**

The major component in each SoftDevice is a wireless protocol stack providing abstract control of the RF transceiver features for wireless applications. For example, fully qualified *Bluetooth* low energy and ANT<sup>™</sup> protocols layers may be implemented in a SoftDevice to provide application developers with an out-of-thebox solution for applications using standard 2.4 GHz protocols.

### **Application Program Interface (API)**

In addition, to a Protocol API enabling wireless applications, there is a nRF API that supports both the SoftDevice manager and the SoC library. The nRF API is consistent across SoftDevices in the nRF51 range of ANT<sup>™</sup> and *Bluetooth* products for code compatibility.

The SoftDevice API is implemented using thread-safe Supervisor Calls (SVC). All application interaction with the stack and libraries is asynchronous and event driven. From the application this looks like regular functions, but no compiling or linking is required. All SVC interface functions will be provided through header files for the SDM, SoC Library, and protocol(s).

SVC calls are conceptually software triggered interrupts with a procedure call standard for parameter passing and return values. Each API call generates an interrupt allowing single-thread API context and SoftDevice function locations to be independent from the application perspective at compile-time. SoftDevice API functions can only be called from lower interrupt priority when compared to the SVC priority. See the Exception (interrupt) management with a SoftDevice section *on page 190*.

RAM



#### Memory isolation and run-time protection

SoftDevice program and data memory are sandboxed<sup>3</sup> to prevent SoftDevice program corruption by the application ensuring robust and predictable performance.

Program memory and RAM are divided into two regions using registers. Region 0 is occupied by the SoftDevice while Region 1 is available to the application.

**Program Memory** 

Code regions are defined when programming a SoftDevice by setting a register defining program code length. RAM regions are defined at run-time when the SoftDevice is enabled. See *Figure 74* for an overview of regions.

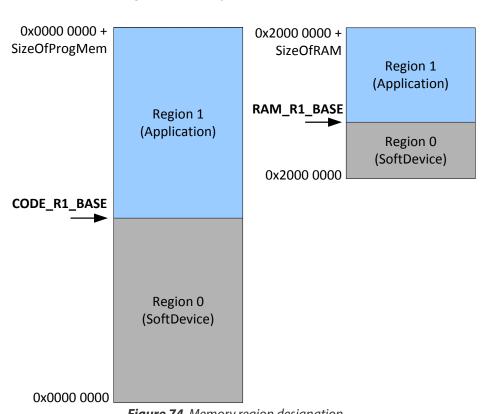


Figure 74 Memory region designation

The SoftDevice uses a fixed amount of flash (program) memory and a variable amount of RAM depending on its state. The flash and RAM usage is specified by size (kilobytes or bytes) and the CODE\_R1\_BASE and RAM\_R1\_BASE addresses which are the usable base addresses of Application code and RAM respectively. Application code must have base address CODE\_R1\_BASE while the Application RAM must be allocated between RAM\_R1\_BASE and the top of RAM, excluding the allocation for the call stack and heap.

Example Application program code address range:

 $CODE_R1_BASE = Program \leq SizeOfProgMem$ 

<sup>3.</sup> A sandbox is a set of access rules for memory imposed on the user.



Example Application RAM address range assuming call stack and heap location as shown in:

#### $RAM_R1\_BASE \le RAM \le (0x2000\ 0000 + SizeOfRAM) - (Call Stack + Heap)$

Sandboxing protects region 0 memory. Region 0 program memory cannot be read, written, or erased<sup>4</sup> at runtime. Region 0 RAM cannot be written to by an application at runtime. Violation of these rules, for example an attempt to write to the protected Region 0 memory, will result in a system Hard Fault as defined in the ARM<sup>®</sup> architecture. There are debugger restrictions applied to these regions which are outlined in *Chapter 8 "Memory Protection Unit (MPU)"* on page 26 that do not affect execution.

When the SoftDevice is disabled the whole of RAM is available to the application. In the context of an enabled SoftDevice however, lower address space of RAM will be "consumed" by the SoftDevice and be marked as write protected.

It is important to note that when the SoftDevice is disabled, the RAM previously used by the application will not be restored. In practice, the application will in many cases want to specify its RAM region from the protected memory length until the end of RAM. This is to make application development easy without having to think about what data to put where.

#### Note:

- The call stack is conventionally located by the initial value of Main Stack Pointer (MSP) at the top address of RAM.
- By default RAM1 block is OFF in System ON-mode. If the MSP initial value defined in the application vector table is in the RAM1 block, the RAM block will be enabled before the application reset vector is executed.
- Do not change the value of MSP dynamically (i.e. never set the MSP register directly).
- RAM located in the SoftDevice's region will be scrambled once the SoftDevice is enabled.
- The RAM scrambled by the SoftDevice will not be recovered on SoftDevice disable.

#### Call stack

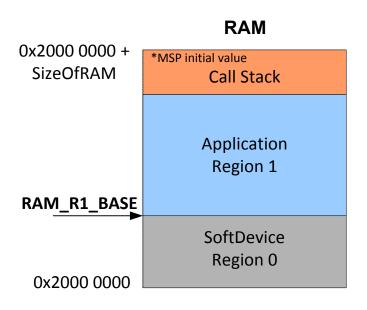
The call stack is defined by the application. The main stack pointer (MSP) gets initialized on reset to the address specified by the application vector table entry 0. The application may, in its reset vector, configure the CPU to use the process stack pointer (PSP) in thread mode. This configuration is optional but may be used by an operating system (OS), for example, to isolate application threads and OS context memory. The application programmer must be aware that the SoftDevice will use the MSP as it is always executed in exception mode.

In configurations without an OS, the main stack grows down and is shared with the nRF51 SoftDevice. The Cortex-M0 has no hardware for detecting stack overflow, and the application is responsible for leaving enough space both for the application itself and the nRF51 SoftDevice stack requirements.

It is customary, but not required, to let the stack run downwards from the upper limit of RAM Region 1.

<sup>4.</sup> An exception is via the Erase All feature which removes all program code from a device.





*Figure 75 Call stack location example* 

With each release of a nRF51 SoftDevice its maximum (worst case) call stack requirement is specified, see the SoftDevice specification for more information. The SoftDevice uses the call stack when LowerStack or UpperStack events occur. These events are asynchronous to the application so the application programmer must reserve call stack for the application in addition to the call stack requirement for the SoftDevice.

#### Heap

At this time there is no heap required by nRF51 SoftDevices. The application is free to allocate and use a heap without disrupting the function of a SoftDevice.

### Peripheral run-time protection

To prevent the application from accidentally disrupting the protocol stack in any way, the application sandbox also protects SoftDevice peripherals. As with program and data memory protection, an attempt to perform a write to a protected peripheral will result in a Hard Fault. Protected peripheral registers are readable by the application, but a write will cause a Hard Fault. Note that peripherals are only protected while the SoftDevice is enabled, otherwise they are available to the application. See the SoftDevice specification for an overview of the peripherals that are restricted by the SoftDevice.

# Exception (interrupt) management with a SoftDevice

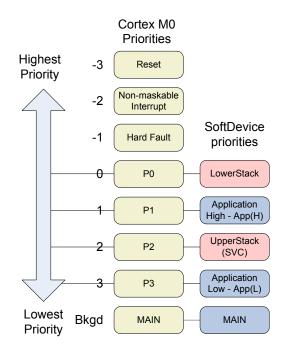
To implement Service Call (SVC) APIs and ensure that embedded protocol real-time requirements are met independent of application processing, the SoftDevice implements an exception model for execution as shown in *Figure 76* on page 191. Care must be taken when selecting the correct interrupt priority for application events according to the guidelines that follow. The NVIC API to the SoC Library supports safe configuration of interrupt priority from the application.

The Cortex-M0 processor has four configurable interrupt priorities ranging from 0 to 3 (with 0 being highest priority). On reset, all interrupts are configured with the highest priority (0).



The highest priority (LowerStack) is reserved by the SoftDevice to service real-time protocol timing requirements and thus must remain unused by the application programmer. The SoftDevice also reserves priority 2 (UpperStack (SVC) priority). This priority is used by higher level, deferrable, SoftDevice tasks and the API functions executed as SVC interrupts (see Interface section *on page 187*).

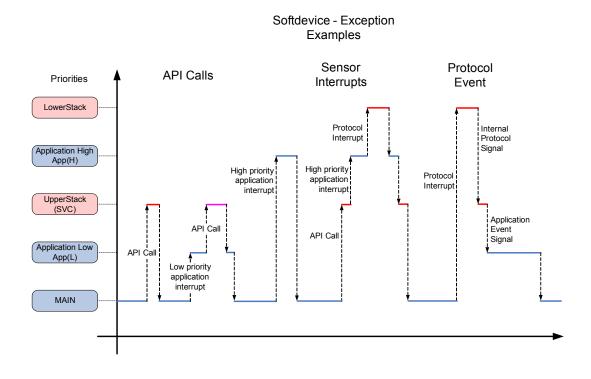
The application provides two configurable priorities, App(H) and App(L), in addition to the background level - main.

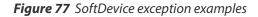


*Figure 76 Exception model* 

As seen from the figure, App(H) is located between the two priorities reserved by the SoftDevice. This enables a low-latency application interrupt in order to support fast sensor interfaces. The App(H) will only experience latency from interrupts in the LowerStack priority, while App(L) can experience latency from LowerStack, App(H) and UpperStack context interrupts.





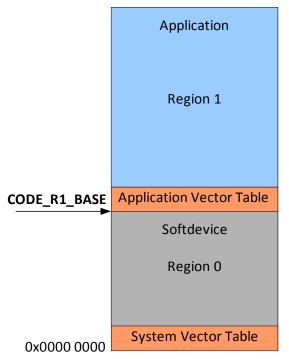


#### Interrupt forwarding to the application

At the lowest level, the SoftDevice Manager receives all interrupts regardless of enabled state. When the SoftDevice is enabled, some interrupt numbers are reserved for use by the protocol stack implemented in the SoftDevice and any handler defined by the application will not receive these interrupts. The reserved interrupts directly correspond to the hardware resource usage of the SoftDevice which can be found in the corresponding SoftDevice Specification. For example, if a SoftDevice (or embedded protocol stack) requires the exclusive use of a peripheral "TIMERO", that peripheral's interrupt handler can be implemented in the application, but will not be executed while the SoftDevice is enabled.

All interrupts corresponding to hardware peripherals not used by the SoftDevice are forwarded directly to the application defined interrupt handler. For the SoftDevice Manager to locate the application interrupt vectors, the application must define its interrupt vector table at the bottom of code Region 1 (see *Figure 78* on page 193). In a majority of toolchains, the base address of the application code is positioned after the top address of the SoftDevice. Then, the code can be developed as a standard ARM<sup>®</sup> Cortex<sup>TM</sup>-M0 application project with the compiler tool creating the interrupt vector table as normal.





#### **Program Memory**

*Figure 78 System and application interrupt vector tables* 

SVC interrupt is handled by SoftDevice manager and the SVC number inspected. If equal or greater than 0x10, the interrupt is processed by the SoftDevice. Values below 0x10 cause the SVC to be forwarded to the application. This allows the application to make use of a range of SVC numbers for its own purpose, for example, for an RTOS.

**Note:** While the Cortex<sup>™</sup>-M0 allows each interrupt to be assigned to an IRQ level 0 to 3, the priorities of the interrupts reserved by the SoftDevice cannot be changed. This includes the SVC interrupt. Handlers running at Application High level have neither access to SoftDevice functions nor to application specific SVCs or RTOS functions running at Application Low level.

If the SoftDevice is not enabled, all interrupts are immediately forwarded to the application specified handler. The exception to this is that SVC interrupts with an SVC number above or equal to 0x10 are not forwarded.

#### **Events - SoftDevice to application**

Software triggered interrupts in reserved IRQ slots are used to signal events from SoftDevice to application. For details on this technique and how to implement handling of these events, refer to the Software Development Kit (SDK) for your device.



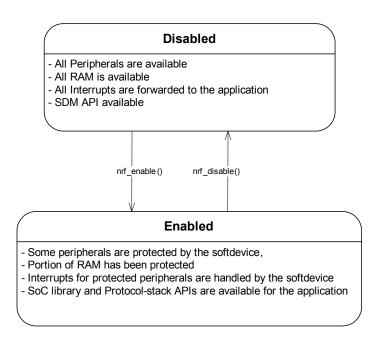
### SoftDevice enable and disable

Before enabling the SoftDevice, you cannot use any capabilities of the SoftDevice. This extends to the use of the SoC library and protocol stack functions. All of the chip's resources are freely available to the application, with some exceptions:

- SVC numbers 0x10 to 0xFF are reserved.
- SoftDevice program memory is reserved.

Once the SoftDevice has been enabled, more restrictions apply:

- Some RAM will be reserved.
- · Some peripherals will be reserved.
- Some of the peripherals that are reserved will have a SoC library interface.
- Interrupts will not arrive in the application for reserved peripherals.
- The reserved peripherals are reset upon SoftDevice disable.
- nrf\_nvic\_ functions must be used instead of CMSIS NVIC\_ functions for safe use of the SoftDevice.
- Maximum interrupt latency will be determined by the SoftDevice .



#### **Power management**

While the SoftDevice is disabled, the application must implement power management at the highest level. After a SoftDevice is enabled, the POWER peripheral will be protected. This means that all interactions with the POWER peripheral must happen through the SoC Library Power API. This API provides an interface for turning on/off peripherals and checking the power status of peripherals that are not protected by the SoftDevice. The application will also have the ability to set the other registers in the peripheral and put the chip in System OFF.



# **Error handling**

All SoftDevice API functions return an error code on success and failure.

Hard Faults are triggered if an application attempts to access memory contrary to the sandbox rules or peripheral configurations at runtime.

An assertion mechanism through a registered callback can indicate fatal failures in the SoftDevice to the application.